A Galois Connection Calculus for Abstract Interpretation (Auxiliary Materials)

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Abstract We introduce a Galois connection calculus for language independent specification of abstract interpretations used in programming language semantics, formal verification, and static analysis. This Galois connection calculus and its type system are typed by abstract interpretation.

Categories and Subject Descriptors D.2.4 [*Software/Program Verification*]

General Terms Algorithms, Languages, Reliability, Security, Theory, Verification.

Keywords Abstract Interpretation, Galois connection, Static Analysis, Verification.

1. Galois connections in Abstract Interpretation In *Abstract interpretation* [3, 4, 6, 7] concrete properties (for example (*e.g.*) of computations) are related to abstract properties (*e.g.* types). The abstract properties are always *sound* approximations of the concrete properties (abstract proofs/static analyzes are always correct in the concrete) and are sometimes *complete* (proofs/analyzes of abstract properties can all be done in the abstract only). *E.g.* types are sound but incomplete [2] while abstract semantics are usually complete [9]. The *concrete domain* $\langle \mathcal{C},\ \sqsubseteq\rangle$ and *abstract domain* $\langle \mathcal{A},\ \preccurlyeq\rangle$ of properties are posets (partial orders being interpreted as implication). When concrete properties all have a $\preccurlyeq$-most precise abstraction, the correspondence is a *Galois connection* (*GC*) $\langle \mathcal{C},\ \sqsubseteq\rangle \underset{\alpha}{\overset{\gamma}{\leftrightarrows}} \langle \mathcal{A},\ \preccurlyeq\rangle$ with *abstraction* $\alpha \in \mathcal{C} \mapsto \mathcal{A}$ and *concretization* $\gamma \in \mathcal{A} \mapsto \mathcal{C}$ satisfying $\forall P \in \mathcal{C} : \forall Q \in \mathcal{A} : \alpha(x) \preccurlyeq y \Leftrightarrow x \sqsubseteq \gamma(y)$ ($\Rightarrow$ expresses soundness and $\Leftarrow$ best abstraction). Each adjoint α/γ uniquely determines the other γ/α. A *Galois retraction* (or *insertion*) has α onto, so γ is one-to-one, and $\alpha \circ \gamma$ is the identity. A *Galois isomorphism* $\langle \mathcal{C},\ \subseteq\rangle \underset{\alpha}{\overset{\gamma}{\leftrightarrows}} \langle \mathcal{A},\ \preccurlyeq\rangle$ has both α and γ one-to-one onto.

2. *GC* calculus syntax The *GC* calculus can describes complex *GCs* between posets representing concrete and abstract properties of the semantics of programs.

2.1 Syntax of program variables

$\mathtt{x}, \ldots \in \mathbb{X}$

2.2 Syntax of labels

$\ell, \ldots \in \mathbb{L}$

2.3 Syntax of ur-elements

The ur-elements are objects (concrete or abstract) which are not a set, but that may be elements of sets.

$e \in \mathbb{E}$
$e ::= \mathsf{true} \mid 1 \mid \infty \mid \mathtt{x} \mid \ell \mid -e \mid \ldots$

2.4 Syntax of sets

$s \in \mathbb{S}$
$s ::= \mathbb{B} \mid \mathbb{Z} \mid \mathbb{X} \mid \mathbb{L} \mid \{e\} \mid [e, e] \mid \mathbb{I}(s, o) \mid s^\infty \mid s \cup s \mid s \mapsto s \mid s \times s \mid \wp(s) \mid \ldots$

Sets are the sets of Booleans, integers, variables, labels, singletons, intervals, the set of all intervals of a poset $\langle s,\ o\rangle$, sets of finite or infinite sequences, unions of sets, sets of functions, of pairs, power sets, *etc.*

2.5 Syntax of partial orders

$o \in \mathbb{O}$
$o ::= \Rightarrow \mid \Leftrightarrow \mid \leq \mid \subseteq \mid \trianglelefteq \mid = \mid o^{-1} \mid o_1 \times o_2 \mid \dot{o} \mid \ddot{o} \mid \ldots$

Orders are implication, equivalence on Booleans, the natural order of integers, inclusion, interval inclusion, equality, the inverse of an order, the pairwise/component by component order, the pointwise order, the double pointwise order, *etc.*

2.6 Syntax of posets

$p \in \mathbb{P}$
$p ::= \langle s,\ o\rangle$

A poset is a set equipped with a partial order.

2.7 Syntax of *GCs*

$g \in \mathbb{G}$
$g ::= \mathbb{1}[p] \mid \top[p, e] \mid \mathbb{I}[p, e, e] \mid \frown[s, s] \mid \cup[s] \mid \neg[s] \mid \infty[s] \mid \rightsquigarrow[s, s] \mid \mapsto[s, s] \mid \times[s, s] \mid \ldots \mid \mathsf{R}[g] \mid s \rightarrowtriangle g \mid g \fatsemi g \mid g \circledast g \mid g \Longmapsto g \mid \ldots$

Basic *GCs* include the identity abstraction $\mathbb{1}[p]$, the top/most abstract abstraction $\top[p, e]$, the interval abstraction $\mathbb{I}[p, e, e]$, the right-image abstraction $\frown[s, s]$, the join abstraction $\cup[s]$, the negation abstraction $\neg[s]$, the sequence abstraction $\infty[s]$, the transformer abstraction $\rightsquigarrow[s, s]$, the function abstraction $\mapsto[s, s]$, the cartesian abstraction $\times[s, s]$, *etc.*.

The Galois connectors include the reduction $\mathsf{R}[g]$, the pointwise extension $s \rightarrowtriangle g$, the composition $g \fatsemi g$, the componentwise composition $g \circledast g$, the function extension $g \Longmapsto g$ at higher-order, *etc.*

3. *GC* calculus semantics The *GC* calculus semantics defines which *GC* is defined by the expressions of the *GC* calculus. Since some *GC* expressions may be ill-defined, the semantics may return a static error Ω (expected to be detectable by sound type systems) or dynamic errors ω (expected no to be detectable by some sound type systems).

* Work supported in part by CMACS, NSF Expedition in Computing award 0926166.

POPL '14, January 22–24, 2014, San Diego, CA, USA.

ACM 978-1-4503-2544-8/14/01.
http://dx.doi.org/10.1145/2535838.2537850

3.1 Semantics of ur-elements

The semantics $\mathcal{S}[\![e]\!]$ of an ur-element e is its value. $\mathcal{O}[\![e]\!]$ is the poset to which this value $\mathcal{S}[\![e]\!]$ belongs.

- $\mathcal{S}[\![\mathsf{x}]\!] \triangleq \mathsf{x}$
  $\mathcal{O}[\![\mathsf{x}]\!] \triangleq \langle \mathbb{X},\ = \rangle$
- $\mathcal{S}[\![\ell]\!] \triangleq \ell$
  $\mathcal{O}[\![\ell]\!] \triangleq \langle \mathbb{L},\ = \rangle$
- $\mathcal{S}[\![\mathsf{true}]\!] \triangleq \mathsf{true}$
  $\mathcal{O}[\![\mathsf{true}]\!] \triangleq \langle \mathbb{B},\ \Rightarrow \rangle$
- $\mathcal{S}[\![1]\!] \triangleq 1$
  $\mathcal{O}[\![1]\!] \triangleq \langle \mathbb{Z} \cup \{-\infty, \infty\},\ \leq \rangle$
- $\mathcal{S}[\![\infty]\!] \triangleq \infty$
  $\mathcal{O}[\![\infty]\!] \triangleq \langle \mathbb{Z} \cup \{-\infty, \infty\},\ \leq \rangle$
- $\mathcal{S}[\![-e]\!] \triangleq \big(\mathcal{S}[\![e]\!] \in \mathbb{B}\ ?\ \neg\mathcal{S}[\![e]\!]\ \|\ \mathcal{S}[\![e]\!] \in \mathbb{Z} \cup \{-\infty, \infty\}\ ?\ -\mathcal{S}[\![e]\!] \mathbin{:} \Omega\big)$
  $\mathcal{O}[\![-e]\!] \triangleq \big(\mathcal{S}[\![e]\!] \in \mathbb{B} \vee \mathcal{S}[\![e]\!] \in \mathbb{Z} \cup \{-\infty, \infty\}\ ?\ \mathcal{O}[\![e]\!] \mathbin{:} \Omega\big)$
- Notice that errors on the semantics $\mathcal{S}[\![e]\!]$ of elements $e \in \mathbb{E}$ are all static *i.e.* Ω.

3.2 Errors

Errors can either be static (Ω) or dynamic (ω). Static errors are expected to be all captured by a static analysis *e.g.* by typing rules. An example is the addition of an integer and a list of booleans. Dynamic errors are those for which static analyses may fail *e.g.* arithmetic overflow which is undecidable. The combination $\mathtt{error}(x, y)$ of errors $x \in \{\omega, \Omega\}$ or $y \in \{\omega, \Omega\}$ is defined so as to give priority to static errors (and to ignore non-erroneous cases).

$\mathtt{error}(x,y)$		x		
		ω	Ω	$\notin \{\omega, \Omega\}$
y	ω	ω	Ω	ω
	Ω	Ω	Ω	Ω
	$\notin \{\omega, \Omega\}$	ω	Ω	———

3.3 Semantics of sets

The semantics $\mathcal{S}[\![s]\!]$ of set expressions $s \in \mathbb{S}$ is the set denoted by that expression.

- $\mathcal{S}[\![\mathbb{X}]\!] \triangleq \mathbb{X}$
- $\mathcal{S}[\![\mathbb{L}]\!] \triangleq \mathbb{L}$
- $\mathcal{S}[\![\mathbb{B}]\!] \triangleq \mathbb{B}$
- $\mathcal{S}[\![\mathbb{Z}]\!] \triangleq \mathbb{Z}$
- $\mathcal{S}[\![\{e\}]\!] \triangleq \big(\mathcal{S}[\![e]\!] \neq \Omega\ ?\ \{\mathcal{S}[\![e]\!]\} \mathbin{:} \Omega\big)$
- $\mathcal{S}[\![[e_1, e_2]]\!] \triangleq \big(\mathcal{O}[\![e_1]\!] = \mathcal{O}[\![e_2]\!] = \langle S,\ o\rangle\ ?\ \{v \in S \mid \langle \mathcal{S}[\![e_1]\!],\ v\rangle \in o \wedge \langle v,\ \mathcal{S}[\![e_2]\!]\rangle \in 0\} \mathbin{:} \Omega\big)$
- $\Im(S, \leqslant) \triangleq \{[v_1, v_2]_\leqslant \mid v_1, v_2 \in S\}$, intervals of S for $\leqslant$ where $[v_1, v_2]_\leqslant \triangleq \{v \in S \mid v_1 \leqslant v \wedge v \leqslant v_2\}$ so $\Im(S, \leqslant) \subseteq \wp(S)$ hence $\Im(S, \leqslant) \in \wp(\wp(S))$
- $\mathcal{S}[\![\mathbb{I}(s, o)]\!] \triangleq \big(\mathcal{S}[\![s]\!] \notin \{\omega, \Omega\}\ ?\ \big(\mathcal{S}[\![o]\!] \subseteq \mathcal{S}[\![s]\!] \times \mathcal{S}[\![s]\!]\ ?\ \Im(\mathcal{S}[\![s]\!], \mathcal{S}[\![o]\!]) \mathbin{:} \omega\big) \mathbin{:} \mathcal{S}[\![s]\!]\big)$
- $\mathcal{S}[\![s^\infty]\!] \triangleq \big(\mathcal{S}[\![s]\!] \notin \{\omega, \Omega\}\ ?\ \mathcal{S}[\![s]\!]^\infty \mathbin{:} \mathcal{S}[\![s]\!]\big)$
- $\mathcal{S}[\![s_1 \cup s_2]\!] \triangleq \big(\mathcal{S}[\![s_1]\!] \notin \{\omega, \Omega\} \wedge \mathcal{S}[\![s_2]\!] \notin \{\omega, \Omega\}\ ?\ \mathcal{S}[\![s_1]\!] \cup \mathcal{S}[\![s_2]\!] \mathbin{:} \mathtt{error}(\mathcal{S}[\![s_1]\!], \mathcal{S}[\![s_2]\!])\big)$
  where $X^\infty \triangleq X^+ \cup X^\omega$, $X^n \triangleq [0, n] \mapsto X$, $n \in \mathbb{N}$ ($\mathsf{dom}(\sigma) = [0, n]$ when $\sigma \in X^n$, $n \in \mathbb{N}$), $X^+ \triangleq \bigcup_{n \in \mathbb{N}} X^n$, and $X^\omega \triangleq \mathbb{N} \mapsto X$ ($\mathsf{dom}(\sigma) = \mathbb{N}$ when $\sigma \in X^\omega$)
- $\mathcal{S}[\![s_1 \mapsto s_2]\!] \triangleq \big(\mathcal{S}[\![s_1]\!] \notin \{\omega, \Omega\} \wedge \mathcal{S}[\![s_2]\!] \notin \{\omega, \Omega\}\ ?\ \mathcal{S}[\![s_1]\!] \mapsto \mathcal{S}[\![s_2]\!] \mathbin{:} \mathtt{error}(\mathcal{S}[\![s_1]\!], \mathcal{S}[\![s_2]\!])\big)$
- $\mathcal{S}[\![s_1 \times s_2]\!] \triangleq \big(\mathcal{S}[\![s_1]\!] \notin \{\omega, \Omega\} \wedge \mathcal{S}[\![s_2]\!] \notin \{\omega, \Omega\}\ ?\ \mathcal{S}[\![s_1]\!] \times \mathcal{S}[\![s_2]\!] \mathbin{:} \mathtt{error}(\mathcal{S}[\![s_1]\!], \mathcal{S}[\![s_2]\!])\big)$
- $\mathcal{S}[\![\wp(s)]\!] \triangleq \big(\mathcal{S}[\![s]\!] \notin \{\omega, \Omega\}\ ?\ \wp(\mathcal{S}[\![s]\!]) \mathbin{:} \mathcal{S}[\![s]\!]\big)$

3.4 Semantics of partial orders

The semantics $\mathcal{S}[\![o]\!]$ of partial order expressions $o \in \mathbb{O}$ is the partial order denoted by that expression. The fact that partial orders are always well-defined *i.e.* $\mathcal{S}[\![o]\!] \notin \{\Omega, \omega\}$ follows from the calculus syntax.

- $\mathcal{S}[\![\Rightarrow]\!] \triangleq \Rightarrow \triangleq \{\langle \mathsf{false},\ \mathsf{false}\rangle, \langle \mathsf{false},\ \mathsf{true}\rangle, \langle \mathsf{true},\ \mathsf{true}\rangle\}$ [1]
- $\mathcal{S}[\![\Leftrightarrow]\!] \triangleq \Leftrightarrow \triangleq \{\langle \mathsf{false},\ \mathsf{false}\rangle, \langle \mathsf{true},\ \mathsf{true}\rangle\}$
- $\mathcal{S}[\![\leq]\!] \triangleq \leq \triangleq \{\langle x,\ y\rangle \mid x, y \in \mathbb{Z} \cup \{-\infty, \infty\} \wedge x \leqslant y\}$
- $\mathcal{S}[\![\subseteq]\!] \triangleq \subseteq \triangleq \{\langle X,\ Y\rangle \mid \forall x \in X : x \in Y\}$ [2]
- $\mathcal{S}[\![\Subset]\!] \triangleq \mathcal{S}[\![\subseteq]\!] = \subseteq$
- $\mathcal{S}[\![=]\!] \triangleq = \triangleq \mathcal{S}[\![\subseteq]\!] \cap \mathcal{S}[\![\subseteq^{-1}]\!]$
- $\mathcal{S}[\![o^{-1}]\!] \triangleq (\mathcal{S}[\![o]\!])^{-1}$ where $R^{-1} \triangleq \{\langle y,\ x\rangle \mid \langle x,\ y\rangle \in R\}$
- $\mathcal{S}[\![o_1 \times o_2]\!] \triangleq \mathcal{S}[\![o_1]\!] \times \mathcal{S}[\![o_2]\!]$ where $R_1 \times R_2 \triangleq \{\langle\langle x,\ x'\rangle,\ \langle y,\ y'\rangle\rangle \mid \langle x,\ y\rangle \in R_1 \wedge \langle x',\ y'\rangle \in R_2\}$
- $\mathcal{S}[\![\dot{o}]\!] \triangleq (\mathcal{S}[\![o]\!]\dot{})$ where $\dot{R} \triangleq \{\langle f,\ g\rangle \mid \forall x \in \mathsf{dom}(f) : x \in \mathsf{dom}(g) \wedge \langle f(x),\ g(x)\rangle \in R\}$
- $\mathcal{S}[\![\ddot{o}]\!] \triangleq (\mathcal{S}[\![\dot{o}]\!]\dot{})$

3.5 Semantics of posets

The semantics $\mathcal{S}[\![p]\!]$ of poset expressions $p \in \mathbb{P}$ is the poset denoted by that expression.

- $\mathcal{S}[\![\langle s,\ o\rangle]\!] \triangleq \big(\mathcal{S}[\![s]\!] \notin \{\omega, \Omega\} \wedge \mathcal{S}[\![o]\!] \notin \{\omega, \Omega\}\ ?\ \langle \mathcal{S}[\![s]\!],\ \mathcal{S}[\![o]\!]\rangle \mathbin{:} \mathtt{error}(\mathcal{S}[\![s]\!], \mathcal{S}[\![o]\!])\big)$

3.6 Semantics of *GCs*

The semantics $\mathcal{S}[\![g]\!]$ of *GCs* $g \in \mathbb{G}$ is the *GC* between posets denoted by that expression.

- $\mathcal{S}[\![\mathbb{1}[p]]\!] \triangleq \big(\mathcal{S}[\![p]\!] \notin \{\omega, \Omega\}\ ?\ \mathcal{S}[\![p]\!] \xleftarrow[\boldsymbol{\lambda}\, P \bullet P]{\boldsymbol{\lambda}\, Q \bullet Q} \mathcal{S}[\![p]\!] \mathbin{:} \mathcal{S}[\![p]\!]\big)$
- $\mathcal{S}[\![\top[p, e]]\!] \triangleq \big(\exists S, \leqslant : \mathcal{S}[\![p]\!] = \langle S,\ \leqslant\rangle\ ?\ \big(\mathcal{S}[\![e]\!] \in S \setminus \{\omega\}\ ?\ \big(\forall x \in S : x \leqslant \mathcal{S}[\![e]\!]\ ?\ \mathcal{S}[\![p]\!] \xleftarrow[\boldsymbol{\lambda}\, P \bullet \mathcal{S}[\![e]\!]]{\boldsymbol{\lambda}\, Q \bullet \mathcal{S}[\![e]\!]} \mathcal{S}[\![p]\!] \mathbin{:} \omega\big) \mathbin{:} \omega\big) \mathbin{:} \mathtt{error}(\mathcal{S}[\![p]\!], \mathcal{S}[\![e]\!])\big)$
- $\mathcal{S}[\![\mathbb{I}[\langle s,\ o\rangle, e_1, e_2]]\!] \triangleq \big(\mathcal{S}[\![s]\!] \notin \{\omega, \Omega\} \wedge \mathcal{S}[\![e_1]\!] \notin \{\omega, \Omega\} \wedge \mathcal{S}[\![e_2]\!] \notin \{\omega, \Omega\}\ ?\ \big(\mathcal{S}[\![o]\!] \subseteq (\mathcal{S}[\![s]\!] \cup \{\mathcal{S}[\![e_1]\!], \mathcal{S}[\![e_2]\!]) \times (\mathcal{S}[\![s]\!] \cup \{\mathcal{S}[\![e_1]\!], \mathcal{S}[\![e_2]\!]) \wedge \forall x \in \mathcal{S}[\![s]\!] : \langle \mathcal{S}[\![e_1]\!], x\rangle \in \mathcal{S}[\![o]\!] \wedge \langle x, \mathcal{S}[\![e_2]\!]\rangle \in \mathcal{S}[\![o]\!]\ ?\ \langle \wp(\mathcal{S}[\![s]\!]), \subseteq\rangle \xleftarrow[\alpha^{\mathbb{I}}]{\gamma^{\mathbb{I}}} \langle \Im(\mathcal{S}[\![s]\!] \cup \{\mathcal{S}[\![e_1]\!], \mathcal{S}[\![e_2]\!]\}, \mathcal{S}[\![o]\!]), \subseteq\rangle \mathbin{:} \omega\big) \mathbin{:} \mathtt{error}(\mathcal{S}[\![s]\!], \mathtt{error}(\mathcal{S}[\![e_1]\!], \mathcal{S}[\![e_2]\!]))\big)$ where
  - $\Im(S, \leqslant) \triangleq \{[v_1, v_2] \mid v_1, v_2 \in S\}$, intervals of S for $\leqslant$ where $[v_1, v_2] \triangleq \{v \in S \mid v_1 \leqslant v \wedge v \leqslant v_2\}$

[1] we use the same mathematical symbol for the syntax and semantics of partial orders in the *GC* calculus, which one is meant should be clear from the context.

[2] The use of proper classes can be avoided by assuming that all sets such as X and Y are taken in a universe of sets such as a Grothendieck universe or the Zermelo/von Neumann hierarchy of sets.

- $\alpha^{\mathbb{I}}(X) \triangleq [\min_{\mathcal{S}[\![o]\!]} X, \max_{\mathcal{S}[\![o]\!]} X], \min_{\mathcal{S}[\![o]\!]} \emptyset \triangleq \mathcal{S}[\![e_2]\!], \max_{\mathcal{S}[\![o]\!]} \emptyset \triangleq \mathcal{S}[\![e_1]\!]$, and
- $\gamma^{\mathbb{I}}([a,b]) \triangleq \{x \in \mathcal{S}[\![s]\!] \mid \langle a,\, x\rangle \in \mathcal{S}[\![o]\!] \wedge \langle x,\, b\rangle \in \mathcal{S}[\![o]\!]\}$

- $\mathcal{S}[\![\curvearrowright[s_{\mathbb{L}}, s_{\mathcal{M}}]]\!] \triangleq \big(\mathcal{S}[\![s_{\mathbb{L}}]\!] \notin \{\Omega,\omega\} \wedge \mathcal{S}[\![s_{\mathcal{M}}]\!] \notin \{\Omega,\omega\} \;?\; \langle \wp(\mathcal{S}[\![s_{\mathbb{L}}]\!] \times \mathcal{S}[\![s_{\mathcal{M}}]\!]),\ \subseteq\rangle \underset{\alpha^{\curvearrowright}}{\overset{\gamma^{\curvearrowright}}{\leftrightarrows}} \langle \mathcal{S}[\![s_{\mathbb{L}}]\!] \mapsto \wp(\mathcal{S}[\![s_{\mathcal{M}}]\!]),\ \dot{\subseteq}\rangle \fatsemi \texttt{error}(\mathcal{S}[\![s_{\mathbb{L}}]\!], \mathcal{S}[\![s_{\mathcal{M}}]\!])\big)$ where
  - $\alpha^{\curvearrowright} \triangleq \boldsymbol{\lambda}\, P \bullet \boldsymbol{\lambda}\, \ell \bullet \{m \mid \langle \ell,\, m\rangle \in P\}$
  - $\gamma^{\curvearrowright} \triangleq \boldsymbol{\lambda}\, Q \bullet \{\langle \ell,\, m\rangle \mid m \in Q(\ell)\}$
  - $\dot{\subseteq}$ is the pointwise extension of inclusion $\subseteq$ that is $f \;\dot{\subseteq}\; g \Leftrightarrow \forall x \in \mathsf{dom} f : x \in \mathsf{dom}(g) \wedge f(x) \subseteq g(x)$.

- $\mathcal{S}[\![\cup[s]]\!] \triangleq \big(\mathcal{S}[\![s]\!] \notin \{\omega,\Omega\} \;?\; \langle \wp(\wp(\mathcal{S}[\![s]\!])),\ \subseteq\rangle \underset{\alpha^{\wp}}{\overset{\gamma^{\wp}}{\leftrightarrows}} \langle \wp(\mathcal{S}[\![s]\!]),\ \subseteq\rangle \fatsemi \mathcal{S}[\![s]\!]\big)$ with
  - $\alpha^{\wp} \triangleq \boldsymbol{\lambda}\, P \in \wp(\wp(\mathcal{S}[\![s]\!])) \bullet \bigcup P$
  - $\gamma^{\wp} \triangleq \boldsymbol{\lambda}\, Q \in \wp(\mathcal{S}[\![s]\!]) \bullet \wp(Q))$

- $\mathcal{S}[\![\neg[s]]\!] \triangleq \big(\mathcal{S}[\![s]\!] \notin \{\omega,\Omega\} \;?\; \langle \wp(\mathcal{S}[\![s]\!]),\ \subseteq\rangle \underset{\neg}{\overset{\neg}{\leftrightarrows}} \langle \wp(\mathcal{S}[\![s]\!]),\ \supseteq\rangle \fatsemi \mathcal{S}[\![s]\!]\big)$,

- $\mathcal{S}[\![\infty[s]]\!] \triangleq \big(\mathcal{S}[\![s]\!] \notin \{\omega,\Omega\} \;?\; \langle \wp(\mathcal{S}[\![s]\!]^{\infty}),\ \subseteq\rangle \underset{\alpha^{\infty}}{\overset{\gamma^{\infty}}{\leftrightarrows}} \langle \wp(\mathcal{S}[\![s]\!]),\ \subseteq\rangle \fatsemi \mathcal{S}[\![s]\!]\big)$ with
  - $\alpha^{\infty} \triangleq \boldsymbol{\lambda}\, P \in \wp(\mathcal{S}[\![s]\!]^{\infty}) \bullet \{\sigma_i \in \mathcal{S}[\![s]\!] \mid \sigma \in P \wedge i \in \mathsf{dom}(\sigma)\}$
  - $\gamma^{\infty} \triangleq \boldsymbol{\lambda}\, Q \in \wp(\mathcal{S}[\![s]\!]) \bullet \{\sigma \in \mathcal{S}[\![s]\!]^{\infty} \mid \forall i \in \mathsf{dom}(\sigma) : \sigma_i \in Q\}$

- $\mathcal{S}[\![\rightsquigarrow[s_1, s_2]]\!] \triangleq \big(\mathcal{S}[\![s_1]\!] \notin \{\omega,\Omega\} \wedge \mathcal{S}[\![s_1]\!] \notin \{\omega,\Omega\} \;?\; \langle \wp(\mathcal{S}[\![s_1]\!] \times \mathcal{S}[\![s_2]\!]),\ \subseteq\rangle \underset{\alpha^{\rightsquigarrow}}{\overset{\gamma^{\rightsquigarrow}}{\leftrightarrows}} \langle \wp(\mathcal{S}[\![s_1]\!]) \xrightarrow{\cup} \wp(\mathcal{S}[\![s_2]\!]),\ \dot{\subseteq}\rangle \fatsemi \texttt{error}(\mathcal{S}[\![s_1]\!], \mathcal{S}[\![s_2]\!])\big)$ mapping relations to join-preserving set transformers with
  - $\alpha^{\rightsquigarrow} \triangleq \boldsymbol{\lambda}\, R \in \wp(\mathcal{S}[\![s_1]\!] \times \mathcal{S}[\![s_2]\!]) \bullet \boldsymbol{\lambda}\, X \in \wp(\mathcal{S}[\![s_1]\!]) \bullet \{y \mid \exists x \in X : \langle x,\, y\rangle \in R\}$
  - $\gamma^{\rightsquigarrow} \triangleq \boldsymbol{\lambda}\, g \in \wp(\mathcal{S}[\![s_1]\!]) \xrightarrow{\cup} \wp(\mathcal{S}[\![s_2]\!]) \bullet \{\langle x,\, y\rangle \mid y \in g(\{x\})\}$

- $\mathcal{S}[\![\mapsto[s_1, s_2]]\!] \triangleq \big(\mathcal{S}[\![s_1]\!] \notin \{\omega,\Omega\} \wedge \mathcal{S}[\![s_2]\!] \notin \{\omega,\Omega\} \;?\; \langle \wp(\mathcal{S}[\![s_1]\!] \mapsto \mathcal{S}[\![s_2]\!]),\ \subseteq\rangle \underset{\alpha^{\mapsto}}{\overset{\gamma^{\mapsto}}{\leftrightarrows}} \langle \wp(\mathcal{S}[\![s_1]\!]) \mapsto \wp(\mathcal{S}[\![s_2]\!]),\ \dot{\subseteq}\rangle \fatsemi \texttt{error}(\mathcal{S}[\![s_1]\!], \mathcal{S}[\![s_2]\!])\big)$ with
  - $\alpha^{\mapsto} \triangleq \boldsymbol{\lambda}\, P \in \wp(\mathcal{S}[\![s_1]\!] \mapsto \mathcal{S}[\![s_2]\!]) \bullet \boldsymbol{\lambda}\, X \in \wp(\mathcal{S}[\![s_1]\!]) \bullet \{f(x) \mid f \in P \wedge x \in X\}$
  - $\gamma^{\mapsto} \triangleq \boldsymbol{\lambda}\, g \in \wp(\mathcal{S}[\![s_1]\!]) \mapsto \wp(\mathcal{S}[\![s_2]\!]) \bullet \{f \in \mathcal{S}[\![s_1]\!] \mapsto \mathcal{S}[\![s_2]\!] \mid \forall X \in \wp(\mathcal{S}[\![s_1]\!]) : \forall x \in X : f(x) \in g(X)\}$,

- $\mathcal{S}[\![\times[s_1, s_2]]\!] \triangleq \big(\mathcal{S}[\![s_1]\!] \notin \{\omega,\Omega\} \wedge \mathcal{S}[\![s_2]\!] \notin \{\omega,\Omega\} \;?\; \langle \wp(\mathcal{S}[\![s_1]\!] \mapsto \mathcal{S}[\![s_2]\!]),\ \subseteq\rangle \underset{\alpha^{\times}}{\overset{\gamma^{\times}}{\leftrightarrows}} \langle \mathcal{S}[\![s_1]\!] \mapsto \wp(\mathcal{S}[\![s_2]\!]),\ \dot{\subseteq}\rangle \fatsemi \texttt{error}(\mathcal{S}[\![s_1]\!], \mathcal{S}[\![s_2]\!])\big)$ with
  - $\alpha^{\times}(X) \triangleq \boldsymbol{\lambda}\, i \in \mathcal{S}[\![s_1]\!] \bullet \{x \in \mathcal{S}[\![s_2]\!] \mid \exists f \in \mathcal{S}[\![s_1]\!] \mapsto \mathcal{S}[\![s_2]\!] : f[i \leftarrow x] \in X\}$
  - $\gamma^{\times}(Y) \triangleq \{f \in \mathcal{S}[\![s_1]\!] \mapsto \mathcal{S}[\![s_2]\!] \mid \forall i \in \mathcal{S}[\![s_1]\!] : f(i) \in Y(i)\}$, and
  - $\dot{\subseteq}$ is the the pointwise extension of $\subseteq$ to $\mathcal{S}[\![s_1]\!]$

- $\mathcal{S}[\![\mathsf{R}[g]]\!] \triangleq \big(\mathcal{S}[\![g]\!] = \langle \mathcal{C},\, \sqsubseteq\rangle \underset{\alpha}{\overset{\gamma}{\leftrightarrows}} \langle \mathcal{A},\, \leqslant\rangle \;?\; \langle \mathcal{C},\, \sqsubseteq\rangle \underset{\alpha}{\overset{\gamma}{\leftrightarrows}} \langle \{\alpha(P) \mid P \in \mathcal{C}\},\, \leqslant\rangle \fatsemi \big(\mathcal{S}[\![g]\!] = \omega \;?\; \omega \fatsemi \Omega\big)\big)$

- $\mathcal{S}[\![s \twoheadrightarrow g]\!] \triangleq \big(\mathcal{S}[\![s]\!] = X \notin \{\omega,\Omega\} \wedge \mathcal{S}[\![g]\!] = \langle \mathcal{C},\, \sqsubseteq\rangle \underset{\alpha}{\overset{\gamma}{\leftrightarrows}} \langle \mathcal{A},\, \leqslant\rangle \;?\; \langle X \mapsto \mathcal{C},\, \dot{\sqsubseteq}\rangle \underset{\boldsymbol{\lambda}\, \rho \bullet \alpha \circ \rho}{\overset{\boldsymbol{\lambda}\, \overline{\rho} \bullet \gamma \circ \overline{\rho}}{\leftrightarrows}} \langle X \mapsto \mathcal{A},\, \dot{\leqslant}\rangle \fatsemi \texttt{error}(\mathcal{S}[\![s]\!], \big(\mathcal{S}[\![g]\!] = \omega \;?\; \omega \fatsemi \Omega\big))\big)$ for the pointwise orderings $\dot{\sqsubseteq}$ and $\dot{\leqslant}$.

- $\mathcal{S}[\![g_1 \fatsemi g_2]\!] \triangleq \big(\mathcal{S}[\![g_1]\!] = p_1 \underset{\alpha_1}{\overset{\gamma_1}{\leftrightarrows}} p_2 \wedge \mathcal{S}[\![g_2]\!] = p_3 \underset{\alpha_2}{\overset{\gamma_2}{\leftrightarrows}} p_4 \;?\; \big(p_2 = p_3 \;?\; p_1 \underset{\alpha_2 \circ \alpha_1}{\overset{\gamma_1 \circ \gamma_2}{\leftrightarrows}} p_4 \fatsemi \omega\big) \fatsemi \texttt{error}(\big(\mathcal{S}[\![g_1]\!] = \omega \;?\; \omega \fatsemi \Omega\big)), \big(\mathcal{S}[\![g_2]\!] = \omega \;?\; \omega \fatsemi \Omega\big))\big)$,

- $\mathcal{S}[\![g_1 \circledast g_2]\!] \triangleq \big(\mathcal{S}[\![g_1]\!] = \langle \mathcal{C}_1,\, \subseteq\rangle \underset{\alpha_1}{\overset{\gamma_1}{\leftrightarrows}} \langle \mathcal{A}_1,\, \sqsubseteq\rangle \wedge \mathcal{S}[\![g_2]\!] = \langle \mathcal{C}_2,\, \Subset\rangle \underset{\alpha_2}{\overset{\gamma_2}{\leftrightarrows}} \langle \mathcal{A}_2,\, \preceq\rangle \;?\; \langle \mathcal{C}_1 \times \mathcal{C}_2,\, \subseteq \times \subseteq\rangle \underset{\alpha_1 \times \alpha_2}{\overset{\gamma_1 \times \gamma_2}{\leftrightarrows}} \langle \mathcal{A}_1 \times \mathcal{A}_2,\, \sqsubseteq \times \preceq\rangle \fatsemi \texttt{error}(\big(\mathcal{S}[\![g_1]\!] = \omega \;?\; \omega \fatsemi \Omega\big)), \big(\mathcal{S}[\![g_2]\!] = \omega \;?\; \omega \fatsemi \Omega\big))\big)$ (generalizing to tuples),

- $\mathcal{S}[\![g_1 \Longrightarrow g_2]\!] \triangleq \big(\mathcal{S}[\![g_1]\!] = \langle \mathcal{C}_1,\, \subseteq\rangle \underset{\alpha_1}{\overset{\gamma_1}{\leftrightarrows}} \langle \mathcal{A}_1,\, \sqsubseteq\rangle \wedge \mathcal{S}[\![g_2]\!] = \langle \mathcal{C}_2,\, \preceq\rangle \underset{\alpha_2}{\overset{\gamma_2}{\leftrightarrows}} \langle \mathcal{A}_2,\, \leqslant\rangle \;?\; \langle \mathcal{C}_1 \xrightarrow{\nearrow} \mathcal{C}_2,\, \dot{\preceq}\rangle \underset{\boldsymbol{\lambda}\, f \bullet \alpha_2 \circ f \circ \gamma_1}{\overset{\boldsymbol{\lambda}\, g \bullet \gamma_2 \circ g \circ \alpha_1}{\leftrightarrows}} \langle \mathcal{A}_1 \xrightarrow{\nearrow} \mathcal{A}_2,\, \dot{\leqslant}\rangle \fatsemi \texttt{error}(\big(\mathcal{S}[\![g_1]\!] = \omega \;?\; \omega \fatsemi \Omega\big)), \big(\mathcal{S}[\![g_2]\!] = \omega \;?\; \omega \fatsemi \Omega\big))\big)$ for increasing maps and pointwise orderings $\dot{\sqsubseteq}$ and $\dot{\leqslant}$.

4. The POPL'77 abstraction

We express the abstraction that was used in the interval example of [3, 4, p. 247] in the *GC* calculus.

4.1 POPL'77 semantics and semantic properties

- $\mathbb{L}$ labels
- $\mathbb{X}$ variables
- $\langle \mathcal{V}, \leq\rangle$ (*e.g.* $\langle \mathbb{Z}, \leqslant\rangle$ complete partial order of values (*e.g.* $\langle \mathbb{Z}, \leqslant\rangle$ or $\langle [\mathtt{minint}, \mathtt{maxint}],\ \leqslant\rangle$)
- $\mathcal{M} \triangleq \mathbb{X} \mapsto \mathcal{V}$ memory states
- $\Sigma \triangleq \mathbb{L} \times \mathcal{M} = \mathbb{L} \times (\mathbb{X} \mapsto \mathcal{V})$ states
- Σ^{∞} program runs/executions, *i.e.* finite or infinite sequences of states
- $\mathcal{S} \triangleq \wp(\Sigma^{\infty})$ semantic domain
- $\wp(\mathcal{S})$ semantic properties (*i.e.* the sets of sets of finite or infinite state sequences $\wp(\wp((\mathbb{L} \times (\mathbb{X} \mapsto \mathbb{Z}))^{\infty}))$)

4.2 POPL'77 reachability abstraction

The reachability abstraction collects states appearing along execution traces into local invariants on variables attached to program points designated by labels.

- G^*
$\triangleq \cup[\Sigma^{\infty}] \fatsemi \infty[\Sigma] \fatsemi \curvearrowright[\mathbb{L}, \mathcal{M}]$
$= \langle \wp(\wp(\Sigma^{\infty})), \subseteq\rangle \underset{\alpha^{\wp}}{\overset{\gamma^{\wp}}{\leftrightarrows}} \langle \wp(\Sigma^{\infty}), \subseteq\rangle \fatsemi \langle \wp(\Sigma^{\infty}), \subseteq\rangle \underset{\alpha^{\infty}}{\overset{\gamma^{\infty}}{\leftrightarrows}} \langle \wp(\Sigma), \subseteq\rangle \fatsemi \curvearrowright[\mathbb{L}, \mathcal{M}]$
$= \langle \wp(\wp(\Sigma^{\infty})), \subseteq\rangle \underset{\alpha^{\infty} \circ \alpha^{\wp}}{\overset{\gamma^{\wp} \circ \gamma^{\infty}}{\leftrightarrows}} \langle \wp(\Sigma), \subseteq\rangle \fatsemi \curvearrowright[\mathbb{L}, \mathcal{M}]$
$= \langle \wp(\wp(\Sigma^{\infty})), \subseteq\rangle \underset{\alpha^{\infty} \circ \alpha^{\wp}}{\overset{\gamma^{\wp} \circ \gamma^{\infty}}{\leftrightarrows}} \langle \wp(\mathbb{L} \times \mathcal{M}), \subseteq\rangle \fatsemi \langle \wp(\mathbb{L} \times \mathcal{M}), \subseteq\rangle \underset{\alpha^{\curvearrowright}}{\overset{\gamma^{\curvearrowright}}{\leftrightarrows}} \langle \mathbb{L} \mapsto \wp(\mathcal{M}),\ \dot{\subseteq}\rangle$
$= \langle \wp(\wp(\Sigma^{\infty})), \subseteq\rangle \underset{\alpha^{\curvearrowright} \circ \alpha^{\infty} \circ \alpha^{\wp}}{\overset{\gamma^{\wp} \circ \gamma^{\infty} \circ \gamma^{\curvearrowright}}{\leftrightarrows}} \langle \mathbb{L} \mapsto \wp(\mathcal{M}),\ \dot{\subseteq}\rangle$
$= \langle \wp(\wp((\mathbb{L} \times (\mathbb{X} \mapsto \mathcal{V}))^{\infty})), \subseteq\rangle \underset{\alpha^{\curvearrowright} \circ \alpha^{\infty} \circ \alpha^{\wp}}{\overset{\gamma^{\wp} \circ \gamma^{\infty} \circ \gamma^{\curvearrowright}}{\leftrightarrows}} \langle \mathbb{L} \mapsto \wp(\mathbb{X} \mapsto \mathcal{V}),\ \dot{\subseteq}\rangle$

4.3 POPL'77 interval cartesian abstraction

- $G^{\Im}$

$\triangleq \mathbb{L} \rightarrowtriangle (\times[\mathbb{X}, \mathcal{V}] \fatsemi (\mathbb{X} \rightarrowtriangle \mathbb{I}[\langle \mathcal{V}, \leq\rangle, -\infty, \infty]))$

$= \mathbb{L} \rightarrowtriangle (\times[\mathbb{X}, \mathcal{V}] \fatsemi (\mathbb{X} \rightarrowtriangle \langle \wp(\mathcal{V}), \subseteq\rangle \underset{\alpha^{\mathbb{I}}}{\overset{\gamma^{\mathbb{I}}}{\leftrightarrows}} \langle \mathbb{I}(\mathcal{V} \cup \{-\infty, \infty\}, \leq), \dot{\in}\rangle))$

$= \mathbb{L} \rightarrowtriangle (\times[\mathbb{X}, \mathcal{V}] \fatsemi \langle \mathbb{X} \mapsto \wp(\mathcal{V}), \dot{\subseteq}\rangle \underset{\boldsymbol{\lambda}\rho\bullet\alpha^{\mathbb{I}}\circ\rho}{\overset{\boldsymbol{\lambda}\overline{\rho}\bullet\gamma^{\mathbb{I}}\circ\overline{\rho}}{\leftrightarrows}} \langle \mathbb{X} \mapsto \mathbb{I}(\mathcal{V} \cup \{-\infty, \infty\}, \leq), \dot{\in}\rangle)$

$= \mathbb{L} \rightarrowtriangle (\langle \wp(\mathbb{X} \mapsto \mathcal{V}), \subseteq\rangle \underset{\alpha^{\times}}{\overset{\gamma^{\times}}{\leftrightarrows}} \langle \mathbb{X} \mapsto \wp(\mathcal{V}), \dot{\subseteq}\rangle \fatsemi \langle \mathbb{X} \mapsto \wp(\mathcal{V}), \dot{\subseteq}\rangle \underset{\boldsymbol{\lambda}\rho\bullet\alpha^{\mathbb{I}}\circ\rho}{\overset{\boldsymbol{\lambda}\overline{\rho}\bullet\gamma^{\mathbb{I}}\circ\overline{\rho}}{\leftrightarrows}} \langle \mathbb{X} \mapsto \mathbb{I}(\mathcal{V} \cup \{-\infty, \infty\}, \leq), \dot{\in}\rangle)$

$= \mathbb{L} \rightarrowtriangle (\langle \wp(\mathbb{X} \mapsto \mathcal{V}), \subseteq\rangle \underset{\boldsymbol{\lambda}P\bullet\alpha^{\mathbb{I}}\circ(\alpha^{\times}(P))}{\overset{\boldsymbol{\lambda}\overline{\rho}\bullet\gamma^{\times}(\gamma^{\mathbb{I}}\circ\overline{\rho})}{\leftrightarrows}} \langle \mathbb{X} \mapsto \mathbb{I}(\mathcal{V} \cup \{-\infty, \infty\}, \leq), \dot{\in}\rangle)$

since

- $(\boldsymbol{\lambda}\rho\bullet\alpha^{\mathbb{I}}\circ\rho)\circ\alpha^{\times}$
- $= \boldsymbol{\lambda}P\bullet(\boldsymbol{\lambda}\rho\bullet\alpha^{\mathbb{I}}\circ\rho)(\alpha^{\times}(P))$
- $= \boldsymbol{\lambda}P\bullet\alpha^{\mathbb{I}}\circ(\alpha^{\times}(P))$
- $= \boldsymbol{\lambda}P\bullet\boldsymbol{\lambda}\mathtt{x}\bullet\alpha^{\mathbb{I}}\circ(\alpha^{\times}(P))(\mathtt{x})$
- $= \boldsymbol{\lambda}P\bullet\boldsymbol{\lambda}\mathtt{x}\bullet\alpha^{\mathbb{I}}((\alpha^{\times}(P))(\mathtt{x}))$

and

- $\gamma^{\times}\circ(\boldsymbol{\lambda}\overline{\rho}\bullet\gamma^{\mathbb{I}}\circ\overline{\rho})$
- $= \boldsymbol{\lambda}\overline{\rho}\bullet\gamma^{\times}((\boldsymbol{\lambda}\overline{\rho}\bullet\gamma^{\mathbb{I}}\circ\overline{\rho})(\overline{\rho}))$
- $= \boldsymbol{\lambda}\overline{\rho}\bullet\gamma^{\times}(\gamma^{\mathbb{I}}\circ\overline{\rho})$

so that

- $G^{\Im}$
- $= \langle \mathbb{L} \mapsto \wp(\mathbb{X} \mapsto \mathcal{V}), \dot{\subseteq}\rangle \underset{\boldsymbol{\lambda}\rho\bullet\boldsymbol{\lambda}\ell\bullet\boldsymbol{\lambda}\mathtt{x}\bullet\alpha^{\mathbb{I}}(\alpha^{\times}(\rho(\ell))(\mathtt{x}))}{\overset{\boldsymbol{\lambda}\overline{\rho}\bullet\boldsymbol{\lambda}\ell\bullet\gamma^{\times}(\boldsymbol{\lambda}\mathtt{x}\bullet(\gamma^{\mathbb{I}}(\overline{\rho}(\ell)(\mathtt{x}))))}{\leftrightarrows}} \langle \mathbb{L} \mapsto \mathbb{X} \mapsto \mathbb{I}(\mathcal{V} \cup \{-\infty, \infty\}, \leq), \ddot{\in}\rangle$

since

- $\boldsymbol{\lambda}\rho\bullet(\boldsymbol{\lambda}P\bullet\alpha^{\mathbb{I}}\circ(\alpha^{\times}(P)))\circ\rho$
- $= \boldsymbol{\lambda}\rho\bullet\boldsymbol{\lambda}\ell\bullet((\boldsymbol{\lambda}P\bullet\alpha^{\mathbb{I}}\circ(\alpha^{\times}(P)))\circ\rho)(\ell)$
- $= \boldsymbol{\lambda}\rho\bullet\boldsymbol{\lambda}\ell\bullet(\boldsymbol{\lambda}P\bullet\alpha^{\mathbb{I}}\circ(\alpha^{\times}(P)))(\rho(\ell))$
- $= \boldsymbol{\lambda}\rho\bullet\boldsymbol{\lambda}\ell\bullet\alpha^{\mathbb{I}}\circ(\alpha^{\times}(\rho(\ell)))$
- $= \boldsymbol{\lambda}\rho\bullet\boldsymbol{\lambda}\ell\bullet\boldsymbol{\lambda}\mathtt{x}\bullet(\alpha^{\mathbb{I}}\circ(\alpha^{\times}(\rho(\ell))))(\mathtt{x})$
- $= \boldsymbol{\lambda}\rho\bullet\boldsymbol{\lambda}\ell\bullet\boldsymbol{\lambda}\mathtt{x}\bullet\alpha^{\mathbb{I}}(\alpha^{\times}(\rho(\ell))(\mathtt{x}))$

and

- $\boldsymbol{\lambda}\overline{\rho}\bullet(\boldsymbol{\lambda}\overline{\rho}'\bullet\gamma^{\times}(\gamma^{\mathbb{I}}\circ\overline{\rho}'))\circ\overline{\rho}$
- $= \boldsymbol{\lambda}\overline{\rho}\bullet\boldsymbol{\lambda}\ell\bullet((\boldsymbol{\lambda}\overline{\rho}'\bullet\gamma^{\times}(\gamma^{\mathbb{I}}\circ\overline{\rho}'))\circ\overline{\rho})(\ell)$
- $= \boldsymbol{\lambda}\overline{\rho}\bullet\boldsymbol{\lambda}\ell\bullet(\boldsymbol{\lambda}\overline{\rho}'\bullet\gamma^{\times}(\gamma^{\mathbb{I}}\circ\overline{\rho}'))(\overline{\rho}(\ell))$
- $= \boldsymbol{\lambda}\overline{\rho}\bullet\boldsymbol{\lambda}\ell\bullet(\gamma^{\times}(\gamma^{\mathbb{I}}\circ(\overline{\rho}(\ell))))$
- $= \boldsymbol{\lambda}\overline{\rho}\bullet\boldsymbol{\lambda}\ell\bullet(\gamma^{\times}(\boldsymbol{\lambda}\mathtt{x}\bullet(\gamma^{\mathbb{I}}\circ(\overline{\rho}(\ell)))(\mathtt{x}))$
- $= \boldsymbol{\lambda}\overline{\rho}\bullet\boldsymbol{\lambda}\ell\bullet\gamma^{\times}(\boldsymbol{\lambda}\mathtt{x}\bullet(\gamma^{\mathbb{I}}(\overline{\rho}(\ell)(\mathtt{x}))))$

4.4 POPL'77 reduced interval cartesian reachability abstraction

- $G^{\Im^*}$

$\triangleq \mathsf{R}[G^* \fatsemi \mathbb{G}^{\Im}]$

$= \mathsf{R}[\langle \wp(\wp(\Sigma^{\infty})), \subseteq\rangle \underset{\alpha^{\frown}\circ\alpha^{\infty}\circ\alpha^{\wp}}{\overset{\gamma^{\wp}\circ\gamma^{\infty}\circ\gamma^{\frown}}{\leftrightarrows}} \langle \mathbb{L} \mapsto \wp(\mathbb{X} \mapsto \mathcal{V}), \dot{\subseteq}\rangle \fatsemi \langle \mathbb{L} \mapsto \wp(\mathbb{X} \mapsto \mathcal{V}), \dot{\subseteq}\rangle \underset{\boldsymbol{\lambda}\rho\bullet\boldsymbol{\lambda}\ell\bullet\boldsymbol{\lambda}\mathtt{x}\bullet\alpha^{\mathbb{I}}(\alpha^{\times}(\rho(\ell))(\mathtt{x}))}{\overset{\boldsymbol{\lambda}\overline{\rho}\bullet\boldsymbol{\lambda}\ell\bullet\gamma^{\times}(\boldsymbol{\lambda}\mathtt{x}\bullet(\gamma^{\mathbb{I}}(\overline{\rho}(\ell)(\mathtt{x}))))}{\leftrightarrows}} \langle \mathbb{L} \mapsto \mathbb{X} \mapsto \mathbb{I}(\mathcal{V} \cup \{-\infty, \infty\}, \leq), \ddot{\in}\rangle]$

$= \mathsf{R}[\langle \wp(\wp(\Sigma^{\infty})), \subseteq\rangle \underset{\boldsymbol{\lambda}P\bullet\boldsymbol{\lambda}\ell\bullet\boldsymbol{\lambda}\mathtt{x}\bullet\alpha^{\mathbb{I}}(\alpha^{\times}((\alpha^{\frown}(\alpha^{\infty}(\alpha^{\wp}(P))))(\ell))(\mathtt{x}))}{\overset{\boldsymbol{\lambda}\overline{\rho}\bullet(\gamma^{\wp}\circ\gamma^{\infty}\circ\gamma^{\frown})(\boldsymbol{\lambda}\ell\bullet\gamma^{\times}(\boldsymbol{\lambda}\mathtt{x}\bullet(\gamma^{\mathbb{I}}(\overline{\rho}(\ell)(\mathtt{x})))))}{\leftrightarrows}} \langle \mathbb{L} \mapsto \mathbb{X} \mapsto \mathbb{I}(\mathcal{V} \cup \{-\infty, \infty\}, \leq), \ddot{\in}\rangle]$

since

- $(\boldsymbol{\lambda}\rho\bullet\boldsymbol{\lambda}\ell\bullet\boldsymbol{\lambda}\mathtt{x}\bullet\alpha^{\mathbb{I}}(\alpha^{\times}(\rho(\ell))(\mathtt{x})))\circ(\alpha^{\frown}\circ\alpha^{\infty}\circ\alpha^{\wp})$
- $= \boldsymbol{\lambda}P\bullet(\boldsymbol{\lambda}\rho\bullet\boldsymbol{\lambda}\ell\bullet\boldsymbol{\lambda}\mathtt{x}\bullet\alpha^{\mathbb{I}}(\alpha^{\times}(\rho(\ell))(\mathtt{x})))((\alpha^{\frown}\circ\alpha^{\infty}\circ\alpha^{\wp})(P))$
- $= \boldsymbol{\lambda}P\bullet\boldsymbol{\lambda}\ell\bullet\boldsymbol{\lambda}\mathtt{x}\bullet\alpha^{\mathbb{I}}(\alpha^{\times}(((\alpha^{\frown}\circ\alpha^{\infty}\circ\alpha^{\wp})(P))(\ell))(\mathtt{x}))$
- $= \boldsymbol{\lambda}P\bullet\boldsymbol{\lambda}\ell\bullet\boldsymbol{\lambda}\mathtt{x}\bullet\alpha^{\mathbb{I}}(\alpha^{\times}(((\alpha^{\frown}\circ\alpha^{\infty}\circ\alpha^{\wp})(P))(\ell))(\mathtt{x}))$
- $= \boldsymbol{\lambda}P\bullet\boldsymbol{\lambda}\ell\bullet\boldsymbol{\lambda}\mathtt{x}\bullet\alpha^{\mathbb{I}}(\alpha^{\times}((\alpha^{\frown}(\alpha^{\infty}(\alpha^{\wp}(P))))(\ell))(\mathtt{x}))$

and

- $(\gamma^{\wp}\circ\gamma^{\infty}\circ\gamma^{\frown})\circ(\boldsymbol{\lambda}\overline{\rho}'\bullet\boldsymbol{\lambda}\ell\bullet\gamma^{\times}(\boldsymbol{\lambda}\mathtt{x}\bullet(\gamma^{\mathbb{I}}(\overline{\rho}'(\ell)(\mathtt{x})))))$
- $= \boldsymbol{\lambda}\overline{\rho}\bullet(\gamma^{\wp}\circ\gamma^{\infty}\circ\gamma^{\frown})((\boldsymbol{\lambda}\overline{\rho}'\bullet\boldsymbol{\lambda}\ell\bullet\gamma^{\times}(\boldsymbol{\lambda}\mathtt{x}\bullet(\gamma^{\mathbb{I}}(\overline{\rho}'(\ell)(\mathtt{x})))))(\overline{\rho}))$
- $= \boldsymbol{\lambda}\overline{\rho}\bullet(\gamma^{\wp}\circ\gamma^{\infty}\circ\gamma^{\frown})(\boldsymbol{\lambda}\ell\bullet\gamma^{\times}(\boldsymbol{\lambda}\mathtt{x}\bullet(\gamma^{\mathbb{I}}(\overline{\rho}(\ell)(\mathtt{x})))))$

so that

- $G^{\Im^*}$
- $= \langle \wp(\wp(\Sigma^{\infty})), \subseteq\rangle \underset{\alpha}{\overset{\gamma}{\leftrightarrows}} \langle \mathbb{L} \mapsto \mathbb{X} \mapsto \mathbb{I}(\mathcal{V} \cup \{-\infty, \infty\}, \leq), \ddot{\in}\rangle$

where

- $\alpha \triangleq \boldsymbol{\lambda}P\bullet\boldsymbol{\lambda}\ell\bullet\mathsf{smash}(\boldsymbol{\lambda}\mathtt{x}\bullet\alpha^{\mathbb{I}}(\alpha^{\times}((\alpha^{\frown}(\alpha^{\infty}(\alpha^{\wp}(P))))(\ell))(\mathtt{x})))$
- $\gamma \triangleq \boldsymbol{\lambda}\overline{\rho}\bullet(\gamma^{\wp}\circ\gamma^{\infty}\circ\gamma^{\frown})(\boldsymbol{\lambda}\ell\bullet\gamma^{\times}(\boldsymbol{\lambda}\mathtt{x}\bullet(\gamma^{\mathbb{I}}(\overline{\rho}(\ell)(\mathtt{x})))))$ and
- the reduction $\mathsf{smash}(\boldsymbol{\lambda}\mathtt{x}\in\mathbb{X}\bullet[a_{\mathtt{x}}, b_{\mathtt{x}}])$ is $\boldsymbol{\lambda}\mathtt{x}\in\mathbb{X}\bullet[\infty, -\infty]$ when some $[a_{\mathtt{x}}, b_{\mathtt{x}}]$ is the empty interval with $b_{\mathtt{x}} < a_{\mathtt{x}}$ and $\boldsymbol{\lambda}\mathtt{x}\in\mathbb{X}\bullet[a_{\mathtt{x}}, b_{\mathtt{x}}]$ otherwise.

In conclusion, the *static* (or *collecting*) *semantics* is the *reachability abstraction* of program properties in $\wp(\mathcal{S})$ that is $G^* \triangleq \cup[\Sigma^{\infty}] \fatsemi \infty[\Sigma] \fatsemi \frown[\mathbb{L}, \mathcal{M}]$ with abstract domain $\mathbb{L} \mapsto \wp(\mathbb{X} \mapsto \mathcal{V})$. The *reduced interval cartesian reachbility abstraction* is $\mathsf{R}[G^* \fatsemi (\mathbb{L} \rightarrowtriangle (\times[\mathbb{X}, \mathcal{V}] \fatsemi (\mathbb{X} \rightarrowtriangle \mathbb{I}[\langle \mathcal{V}, \leq\rangle, -\infty, \infty])))]$ that is the abstraction $\langle \wp(\wp((\mathbb{L} \times (\mathbb{X} \mapsto \mathcal{V}))^{\infty})), \subseteq\rangle \underset{\alpha}{\overset{\gamma}{\leftrightarrows}} \langle \mathbb{L} \mapsto \mathbb{X} \mapsto \mathbb{I}(\mathcal{V} \cup \{-\infty, \infty\}), \ddot{\in}\rangle$ where

$$\alpha(P) \triangleq \boldsymbol{\lambda}\ell\bullet\mathsf{smash}(\boldsymbol{\lambda}\mathtt{x}\bullet\alpha^{\mathbb{I}}(\alpha^{\times}((\alpha^{\frown}(\alpha^{\infty}(\alpha^{\wp}(P))))(\ell))(\mathtt{x})))$$

and the reduction $\mathsf{smash}(\boldsymbol{\lambda}\mathtt{x}\in\mathbb{X}\bullet[a_{\mathtt{x}}, b_{\mathtt{x}}])$ is $\boldsymbol{\lambda}\mathtt{x}\in\mathbb{X}\bullet[\infty, -\infty]$ to the empty interval $[\infty, -\infty]$ when some $[a_{\mathtt{x}}, b_{\mathtt{x}}]$ is $[\infty, -\infty]$ and $\boldsymbol{\lambda}\mathtt{x}\in\mathbb{X}\bullet[a_{\mathtt{x}}, b_{\mathtt{x}}]$ otherwise.

5. Syntax of types

The intuition for types is that a *GC* expression has type T if and only if the semantics of this expression belongs to the set described by type T. For example true has type `bool` which describes the set $\mathbb{B}$ of Booleans since true $\in \mathbb{B}$.

5.1 Syntax of ur-element types

$\mathsf{E} \in \mathfrak{E}$

$\mathsf{E} ::= \mathtt{var} \mid \mathtt{lab} \mid \mathtt{bool} \mid \mathtt{int} \mid \mathtt{err}$

5.2 Syntax of set types

$\mathsf{S} \in \mathfrak{S}$

$\mathsf{S} ::= \mathtt{P}\ \mathsf{E} \mid \mathtt{P}\ \mathsf{S} \mid \mathtt{seq}\ \mathsf{S} \mid \mathsf{S} \nrightarrow \mathsf{S} \mid \mathsf{S} * \mathsf{S} \mid \mathtt{err}$

5.3 Syntax of partial order types

$\mathsf{O} \in \mathfrak{O}$

$\mathsf{O} ::= \Rightarrow \mid \Leftrightarrow \mid \leq \mid \subseteq \mid = \mid \mathsf{O}^{-1} \mid \mathsf{O} \star \mathsf{O} \mid \dot{\mathsf{O}} \mid \ldots \mid \mathtt{err}$

5.4 Syntax of poset types

$\mathsf{P} \in \mathfrak{P}$

$\mathsf{P} ::= \mathsf{S} \circledast \mathsf{O} \mid \mathtt{err}$

5.5 Syntax of *GC* types

$\mathsf{T} \in \mathfrak{T}$

$\mathsf{T} ::= \mathsf{P} \leftrightharpoons \mathsf{P} \mid \mathsf{S} \nrightarrow \mathsf{T} \mid \mathtt{err}$

6. Semantics of types A error type $\mathtt{err}$ denotes any possible value, including the static error (Ω *e.g.* captured by type systems) and the dynamic error (ω *e.g.* possibly captured by static analyzes more powerful than type systems). The error type $\mathtt{err}$ allows to conclude absolutely nothing on a value of that type, including this these value exists or is erroneous.

- $\gamma^{\mathfrak{U}}(\mathtt{err}) \triangleq \mathcal{S}(\mathbb{U}) \cup \{\omega, \Omega\}$ for all $\langle \mathbb{U}, \mathfrak{U} \rangle \in \{\langle \mathbb{E}, \mathfrak{E} \rangle, \langle \mathbb{S}, \mathfrak{S} \rangle, \langle \mathbb{O}, \mathfrak{O} \rangle, \langle \mathbb{P}, \mathfrak{P} \rangle, \langle \mathbb{G}, \mathfrak{T} \rangle\}$ where $\mathcal{S}(\mathbb{U}) \triangleq \{\mathcal{S}[\![u]\!] \mid u \in \mathbb{U}\}$.

6.1 Semantics of ur-element types

- $\gamma^{\mathfrak{E}}(\mathtt{bool}) \triangleq \mathbb{B}$
- $\gamma^{\mathfrak{E}}(\mathtt{int}) \triangleq \mathbb{Z} \cup \{-\infty, \infty\}$
- $\gamma^{\mathfrak{E}}(\mathtt{var}) \triangleq \mathbb{X}$
- $\gamma^{\mathfrak{E}}(\mathtt{lab}) \triangleq \mathbb{L}$
- Observe that $\gamma^{\mathfrak{E}}$ is injective ($\mathsf{E} \neq \mathsf{E}' \Rightarrow \gamma^{\mathfrak{E}}(\mathsf{E}) \neq \gamma^{\mathfrak{E}}(\mathsf{E}')$).

6.2 Semantics of set types

- $\gamma^{\mathfrak{S}}(\mathtt{P}\ \mathsf{E}) \triangleq \wp(\gamma^{\mathfrak{E}}(\mathsf{E}))$
- $\gamma^{\mathfrak{S}}(\mathtt{P}\ \mathsf{S}) \triangleq \wp(\gamma^{\mathfrak{S}}(\mathsf{S}))$
  It follows that $\{\wp(X) \mid X \in \gamma^{\mathfrak{S}}(\mathsf{S})\} \subseteq \gamma^{\mathfrak{S}}(\mathtt{P}\ \mathsf{S})$.
- $\gamma^{\mathfrak{S}}(\mathtt{seq}\ \mathsf{S}) \triangleq \{X^{\infty} \mid X \in \gamma^{\mathfrak{S}}(\mathsf{S})\}$ where X^{∞} is the set of all finite or infinite sequences of elements of X
- $\gamma^{\mathfrak{S}}(\mathsf{S}_1 \nrightarrow \mathsf{S}_2) \triangleq \{X \mapsto Y \mid X \in \gamma^{\mathfrak{S}}(\mathsf{S}_1) \wedge Y \in \gamma^{\mathfrak{S}}(\mathsf{S}_2)\}$
- $\gamma^{\mathfrak{S}}(\mathsf{S}_1 * \mathsf{S}_2) \triangleq \{X \times Y \mid X \in \gamma^{\mathfrak{S}}(\mathsf{S}_1) \wedge Y \in \gamma^{\mathfrak{S}}(\mathsf{S}_2)\}$
- $\forall \mathsf{S} \in \mathfrak{S} : \gamma^{\mathfrak{S}}(\mathsf{S})$ is $\subseteq$-downward closed meaning that $X \subseteq Y \in \gamma^{\mathfrak{S}}(\mathsf{S}) \Rightarrow X \in \gamma^{\mathfrak{S}}(\mathsf{S})$. The proof is by structural induction on S.
- Observe that $\gamma^{\mathfrak{S}}$ is injective (the proof is by structural induction where $\gamma^{\mathfrak{E}}$ is injective for the base case).

6.3 Semantics of partial order types

- $\gamma^{\mathfrak{O}}(\mathsf{O}) \triangleq \{\mathcal{S}[\![\mathsf{O}]\!]\}, \quad \mathsf{O} \in \{\Rightarrow, \Leftrightarrow, \leq, \subseteq, =\}$
- $\gamma^{\mathfrak{O}}(\mathsf{O}^{-1}) \triangleq \{R^{-1} \mid R \in \gamma^{\mathfrak{O}}(\mathsf{O})\}$
- $\gamma^{\mathfrak{O}}(\mathsf{O}_1 \star \mathsf{O}_2) \triangleq \gamma^{\mathfrak{O}}(\mathsf{O}_1) \times \gamma^{\mathfrak{O}}(\mathsf{O}_2) \triangleq \{R_1 \times R_2 \mid R_1 \in \gamma^{\mathfrak{O}}(\mathsf{O}_1) \wedge R_2 \in \gamma^{\mathfrak{O}}(\mathsf{O}_2)\}$
- $\gamma^{\mathfrak{O}}(\dot{\mathsf{O}}) \triangleq \{\dot{\leqslant} \mid \leqslant \in \gamma^{\mathfrak{O}}(\mathsf{O})\}$

Observe, by structural induction on O, that $\gamma^{\mathfrak{O}}(\mathsf{O})$ is always a singleton and $\gamma^{\mathfrak{O}}$ is injective.

6.4 Semantics of poset types

- $\gamma^{\mathfrak{P}}(\mathsf{S} \circledast \mathsf{O}) \triangleq \gamma^{\mathfrak{S}}(\mathsf{S}) \times \gamma^{\mathfrak{O}}(\mathsf{O})$
- $\gamma^{\mathfrak{P}}$ is injective.

6.5 Semantics of *GC* types

- $\gamma^{\mathfrak{T}}(\mathsf{P} \leftrightharpoons \mathsf{P}') \triangleq \{P \underset{\alpha}{\overset{\gamma}{\leftrightarrows}} P' \mid P \in \gamma^{\mathfrak{P}}(\mathsf{P}) \wedge P' \in \gamma^{\mathfrak{P}}(\mathsf{P}')\}$ [3].
- $\gamma^{\mathfrak{T}}(\mathsf{S} \nrightarrow \mathsf{T}) \triangleq \{\langle X \mapsto \mathcal{C}, \dot{\sqsubseteq} \rangle \underset{\alpha'}{\overset{\gamma'}{\leftrightarrows}} \langle X \mapsto \mathcal{A}, \dot{\leqslant} \rangle \mid X \in \gamma^{\mathfrak{S}}(\mathsf{S}) \wedge \langle \mathcal{C}, \sqsubseteq \rangle \underset{\alpha}{\overset{\gamma}{\leftrightarrows}} \langle \mathcal{A}, \leqslant \rangle \in \gamma^{\mathfrak{T}}(\mathsf{T})\}$
- $\gamma^{\mathfrak{T}}$ is injective.

7. Type membership Type membership $\mathbb{C}$ is the abstraction of set membership $\in$.

- $\mathsf{E}\ \mathbb{C}\ \mathsf{S} \Rightarrow \forall e \in \gamma^{\mathfrak{E}}(\mathsf{E}) : \exists s \in \gamma^{\mathfrak{S}}(\mathsf{S}) : e \in s$
- $\mathsf{S}_1\ \mathbb{C}\ \mathsf{S}_2 \Rightarrow \forall s_1 \in \gamma^{\mathfrak{S}}(\mathsf{S}_1) : \exists s_2 \in \gamma^{\mathfrak{S}}(\mathsf{S}_2) : s_1 \in s_2$

so that

- $\mathsf{E}\ \mathbb{C}\ \mathtt{P}\ \mathsf{E}$
  For example $\mathtt{bool}\ \mathbb{C}\ \mathtt{P}\ \mathtt{bool}$ since $\forall b \in \gamma^{\mathfrak{E}}(\mathtt{bool}) = \mathbb{B}$: $\exists B \in \{\emptyset, \{\mathsf{true}\}, \{\mathsf{false}\}, \mathbb{B}\} = \wp(\mathbb{B}) = \wp(\gamma^{\mathfrak{E}}(\mathtt{bool})) = \gamma^{\mathfrak{S}}(\mathtt{P}\ \mathtt{bool})$: $b \in B$.
- $\mathsf{S}\ \mathbb{C}\ \mathtt{P}\ \mathsf{S}$

Observe that if $\mathsf{E} \notin \mathsf{S}$ (*e.g.* $\mathtt{bool} \notin \mathtt{P}\ \mathtt{int}$) then $\exists e \in \gamma^{\mathfrak{E}}(\mathsf{E}) : \forall s \in \gamma^{\mathfrak{S}}(\mathsf{S}) : e \notin s$ so e of type S belongs to none of the sets s of type S (*e.g.* true belongs to no set of integers).

8. Type preorder

- $\mathsf{T}_1 \trianglelefteqslant \mathsf{T}_2 \triangleq \gamma^{\mathfrak{T}}(\mathsf{T}_1) \subseteq \gamma^{\mathfrak{T}}(\mathsf{T}_2)$ and similarly for $\mathfrak{S}$, $\mathfrak{O}$, $\mathfrak{P}$, specifically $\mathsf{E}_1 \trianglelefteqslant \mathsf{E}_2 \triangleq \mathsf{E}_1 = \mathsf{E}_2$ for $\mathfrak{E}$.

The situation would be different if *e.g.* basic types where enriched by intervals *e.g.* as in the PASCAL language, in which case $\trianglelefteqslant$ would be interval inclusion.

The relation $\trianglelefteqslant$ is reflexive $\mathsf{T} \trianglelefteqslant \mathsf{T}$, transitive $\mathsf{T}_1 \trianglelefteqslant \mathsf{T}_2 \wedge \mathsf{T}_2 \trianglelefteqslant \mathsf{T}_3 \Rightarrow \mathsf{T}_1 \trianglelefteqslant \mathsf{T}_3$ and $\mathtt{err}$ is the supremum (top element). To ensure the existence of a best abstraction of the empty set (*i.e.* false property), it may be useful to add a type infimum (bottom element) $\varnothing$ such that $\gamma^{\mathfrak{E}}(\varnothing) \triangleq \emptyset$. It follows from the definition of $\trianglelefteqslant$ that

- $\mathsf{E} \trianglelefteqslant \mathsf{E}' \Rightarrow \mathtt{P}\ \mathsf{E} \trianglelefteqslant \mathtt{P}\ \mathsf{E}'$
- $\mathsf{S} \trianglelefteqslant \mathsf{S}' \Rightarrow \mathtt{P}\ \mathsf{S} \trianglelefteqslant \mathtt{P}\ \mathsf{S}'$
- $\mathsf{S} \trianglelefteqslant \mathsf{S}' \Rightarrow \mathtt{seq}\ \mathsf{S} \trianglelefteqslant \mathtt{seq}\ \mathsf{S}'$
- $\mathsf{S}_1 \trianglelefteqslant \mathsf{S}'_1 \wedge \mathsf{S}_2 \trianglelefteqslant \mathsf{S}'_2 \Rightarrow \mathsf{S}_1 \nrightarrow \mathsf{S}_2 \trianglelefteqslant \mathsf{S}'_1 \nrightarrow \mathsf{S}'_2$
- $\mathsf{S}_1 \trianglelefteqslant \mathsf{S}'_1 \wedge \mathsf{S}_2 \trianglelefteqslant \mathsf{S}'_2 \Rightarrow \mathsf{S}_1 * \mathsf{S}_2 \trianglelefteqslant \mathsf{S}'_1 * \mathsf{S}'_2$
- $\Leftrightarrow \trianglelefteqslant \Rightarrow, = \trianglelefteqslant \leq, = \trianglelefteqslant \subseteq, = \trianglelefteqslant \Subset$
- $\mathsf{O} \trianglelefteqslant \mathsf{O}' \Rightarrow \mathsf{O}^{-1} \trianglelefteqslant \mathsf{O}'^{-1}$
- $\mathsf{O}_1 \trianglelefteqslant \mathsf{O}'_1 \wedge \mathsf{O}_2 \trianglelefteqslant \mathsf{O}'_2 \Rightarrow \mathsf{O}_1 \star \mathsf{O}_2 \trianglelefteqslant \mathsf{O}'_1 \star \mathsf{O}'_2$
- $\mathsf{O} \trianglelefteqslant \mathsf{O}' \Rightarrow \dot{\mathsf{O}} \trianglelefteqslant \dot{\mathsf{O}}'$
- $\mathsf{S} \trianglelefteqslant \mathsf{S}' \wedge \mathsf{O} \trianglelefteqslant \mathsf{O}' \Rightarrow \mathsf{S} \circledast \mathsf{O} \trianglelefteqslant \mathsf{S}' \circledast \mathsf{O}'$
- $\mathsf{P}_1 \trianglelefteqslant \mathsf{P}'_1 \wedge \mathsf{P}_2 \trianglelefteqslant \mathsf{P}'_2 \Rightarrow \mathsf{P}_1 \leftrightharpoons \mathsf{P}_2 \trianglelefteqslant \mathsf{P}'_1 \leftrightharpoons \mathsf{P}'_2$
- $\mathsf{S} \trianglelefteqslant \mathsf{S}' \wedge \mathsf{T} \trianglelefteqslant \mathsf{T}' \Rightarrow \mathsf{S} \nrightarrow \mathsf{T} \trianglelefteqslant \mathsf{S}' \nrightarrow \mathsf{T}'$

9. Type equivalence

- $\mathsf{T}_1 \cong \mathsf{T}_2 \triangleq \mathsf{T}_1 \trianglelefteqslant \mathsf{T}_2 \wedge \mathsf{T}_2 \trianglelefteqslant \mathsf{T}_1 = \gamma^{\mathfrak{S}}(\mathsf{T}_1) = \gamma^{\mathfrak{S}}(\mathsf{T}_2)$

It follows that we have the rule $\mathsf{X} \trianglelefteqslant \mathsf{Y} \wedge \mathsf{Y} \trianglelefteqslant \mathsf{X} \Rightarrow \mathsf{X} \cong \mathsf{Y}$ and so $(\mathsf{O}^{-1})^{-1} \cong \mathsf{O}$, $\ddot{\mathsf{O}} \cong \dot{\mathsf{O}}$, *etc.*

[3] By the *GC* notation $\langle C, \preceq \rangle \underset{\alpha}{\overset{\gamma}{\leftrightarrows}} \langle A, \sqsubseteq \rangle, \forall x \in C : \forall y \in A : \alpha(x) \sqsubseteq y \Leftrightarrow x \preceq \gamma(y)$.

10. Type inference algorithm We provide the type inference algorithm and prove its soundness. The classical theory of types, which proceeds by subject reduction, that is by induction on the computation steps of an operational semantics, is not helpful since the *GC* calculus does not describe computations. However Abstract Interpretation is not bound to a specific kind of language, semantics (*e.g.* denotational semantics in [2]), and form of induction (*e.g.* no subject reduction in [2]). So soundness is proved by structural induction on the syntax of types.

The type inference algorithm is

- $\mathscr{C}[\![\mathsf{true}]\!] = \mathscr{C}[\![\mathsf{false}]\!] \triangleq \mathtt{bool}$
- $\mathscr{C}[\![0]\!] = \mathscr{C}[\![1]\!] = \ldots = \mathscr{C}[\![\infty]\!] \triangleq \mathtt{int}$,
- $\mathscr{C}[\![\mathsf{x}]\!] = \mathscr{C}[\![\mathsf{y}]\!] = \ldots \triangleq \mathtt{var}$
- $\mathscr{C}[\![\ell]\!] = \ldots \triangleq \mathtt{lab}$
- $\mathscr{C}[\![-e]\!] \triangleq \big(\mathscr{C}[\![e]\!] = \mathtt{bool} \vee \mathscr{C}[\![e]\!] = \mathtt{int} \mathbin{?} \mathscr{C}[\![e]\!] \mathbin{:} \mathtt{err}\big)$
- $\mathcal{S}[\![\mathbb{B}]\!] \triangleq \mathbf{P}\ \mathtt{bool}$
- $\mathcal{S}[\![\mathbb{Z}]\!] \triangleq \mathbf{P}\ \mathtt{int}$
- $\mathcal{S}[\![\mathbb{X}]\!] \triangleq \mathbf{P}\ \mathtt{var}$
- $\mathcal{S}[\![\mathbb{L}]\!] \triangleq \mathbf{P}\ \mathtt{lab}$
- $\mathcal{S}[\![\{e\}]\!] \triangleq \big(\mathscr{C}[\![e]\!] \neq \mathtt{err} \mathbin{?} \mathbf{P}\ \mathscr{C}[\![e]\!] \mathbin{:} \mathtt{err}\big)$
- $\mathcal{S}[\![[e_1, e_2]]\!] \triangleq \big(\mathscr{C}[\![e_1]\!] \cong \mathscr{C}[\![e_2]\!] \neq \mathtt{err} \mathbin{?} \mathbf{P}\ \mathscr{C}[\![e_1]\!] \mathbin{:} \mathtt{err}\big)$
- $\mathcal{S}[\![\mathbb{I}(s, o)]\!] \triangleq \big(\mathcal{S}[\![s]\!] \neq \mathtt{err} \mathbin{?} \mathbf{P}\ \mathcal{S}[\![s]\!] \mathbin{:} \mathtt{err}\big)$
- $\mathcal{S}[\![s^\infty]\!] \triangleq \big(\mathcal{S}[\![s]\!] \neq \mathtt{err} \mathbin{?} \mathtt{seq}\ \mathcal{S}[\![s]\!] \mathbin{:} \mathtt{err}\big)$
- $\mathcal{S}[\![s_1 \cup s_2]\!] \triangleq \big(\mathtt{err} \neq \mathcal{S}[\![s_1]\!] \cong \mathcal{S}[\![s_2]\!] \neq \mathtt{err} \mathbin{?} \mathcal{S}[\![s_1]\!] \mathbin{:} \mathtt{err}\big)$ (note the approximation that s_1 and s_2 must be of the same type as is the case for alternatives of conditionals in functional languages)
- $\mathcal{S}[\![s_1 \mapsto s_2]\!] \triangleq \big(\mathcal{S}[\![s_1]\!] \neq \mathtt{err} \wedge \mathcal{S}[\![s_2]\!] \neq \mathtt{err} \mathbin{?} \mathcal{S}[\![s_1]\!] \mathbin{\ast\!\!\rightarrow} \mathcal{S}[\![s_2]\!] \mathbin{:} \mathtt{err}\big)$
- $\mathcal{S}[\![s_1 \times s_2]\!] \triangleq \big(\mathcal{S}[\![s_1]\!] \neq \mathtt{err} \wedge \mathcal{S}[\![s_2]\!] \neq \mathtt{err} \mathbin{?} \mathcal{S}[\![s_1]\!] * \mathcal{S}[\![s_2]\!] \mathbin{:} \mathtt{err}\big)$
- $\mathcal{S}[\![\wp(s)]\!] \triangleq \big(\mathcal{S}[\![s]\!] \neq \mathtt{err} \mathbin{?} \mathbf{P}\ \mathcal{S}[\![s]\!] \mathbin{:} \mathtt{err}\big)$
- $\mathcal{O}[\![o]\!] \triangleq o, \quad o \in \{\Rightarrow, \Leftrightarrow, \leq, \subseteq, =\}$
- $\mathcal{O}[\![\in]\!] \triangleq \subseteq$
- $\mathcal{O}[\![o^{-1}]\!] \triangleq (\mathcal{O}[\![o]\!])^{-1}$
- $\mathcal{O}[\![o_1 \times o_2]\!] \triangleq \mathcal{O}[\![o_1]\!] * \mathcal{O}[\![o_2]\!]$
- $\mathcal{O}[\![\dot{o}]\!] \triangleq (\dot{\mathcal{O}}[\![o]\!])$
- $\mathcal{O}[\![\ddot{o}]\!] \triangleq (\ddot{\mathcal{O}}[\![o]\!]))$
- $\mathcal{P}[\![\langle s,\ o\rangle]\!] \triangleq \big(\mathcal{S}[\![s]\!] \neq \mathtt{err} \wedge \mathcal{O}[\![o]\!] \neq \mathtt{err} \mathbin{?} \mathcal{S}[\![s]\!] \circledast \mathcal{O}[\![o]\!] \mathbin{:} \mathtt{err}\big)$
- $\mathcal{T}[\![\mathbb{1}[p]]\!] \triangleq \big(\mathcal{P}[\![p]\!] \neq \mathtt{err} \mathbin{?} \mathcal{P}[\![p]\!] \rightleftharpoons \mathcal{P}[\![p]\!] \mathbin{:} \mathtt{err}\big)$
- $\mathcal{T}[\![\top[p, e]]\!] \triangleq \big(\mathscr{C}[\![e]\!] \neq \mathtt{err} \wedge \exists \mathsf{S} \in \mathfrak{S}, \mathsf{O} \in \mathfrak{O} : \mathcal{P}[\![p]\!] = \mathsf{S} \circledast \mathsf{O} \wedge \mathscr{C}[\![e]\!] \Subset \mathsf{S} \mathbin{?} \mathcal{P}[\![p]\!] \rightleftharpoons \mathcal{P}[\![p]\!] \mathbin{:} \mathtt{err}\big)$
- $\mathcal{T}[\![\mathbb{I}[\langle s,\ o\rangle, b, t]]\!] \triangleq \big(\mathtt{err} \neq \mathscr{C}[\![b]\!] \Subset \mathcal{S}[\![s]\!] \neq \mathtt{err} \wedge \mathtt{err} \neq \mathscr{C}[\![t]\!] \Subset \mathcal{S}[\![s]\!] \mathbin{?} (\mathbf{P}\ \mathcal{S}[\![s]\!] \circledast \subseteq) \rightleftharpoons (\mathbf{P}\ \mathcal{S}[\![s]\!] \circledast \subseteq) \mathbin{:} \mathtt{err}\big)$
- $\mathcal{T}[\![\frown[s_{\mathbb{L}}, s_{\mathcal{M}}]]\!] \triangleq \big(\mathcal{S}[\![s_{\mathbb{L}}]\!] \neq \mathtt{err} \wedge \mathcal{S}[\![s_{\mathcal{M}}]\!] \neq \mathtt{err} \mathbin{?} \mathbf{P}\ (\mathcal{S}[\![s_{\mathbb{L}}]\!] * \mathcal{S}[\![s_{\mathcal{M}}]\!]) \circledast \subseteq \rightleftharpoons \mathcal{S}[\![s_{\mathbb{L}}]\!] \mathbin{\ast\!\!\rightarrow} \mathbf{P}\ \mathcal{S}[\![s_{\mathcal{M}}]\!] \circledast \dot{\subseteq} \mathbin{:} \mathtt{err}\big)$
- $\mathcal{T}[\![\cup[s]]\!] \triangleq \big(\mathcal{S}[\![s]\!] \neq \mathtt{err} \mathbin{?} \mathbf{P}\ (\mathbf{P}\ \mathcal{S}[\![s]\!]) \circledast \subseteq \rightleftharpoons \mathbf{P}\ \mathcal{S}[\![s]\!] \circledast \subseteq \mathbin{:} \mathtt{err}\big)$
- $\mathcal{T}[\![\neg[s]]\!] \triangleq \big(\mathcal{S}[\![s]\!] \neq \mathtt{err} \mathbin{?} \mathbf{P}\ \mathcal{S}[\![s]\!] \circledast \subseteq \rightleftharpoons \mathbf{P}\ \mathcal{S}[\![s]\!] \circledast \subseteq^{-1} \mathbin{:} \mathtt{err}\big)$
- $\mathcal{T}[\![\infty[s]]\!] \triangleq \big(\mathcal{S}[\![s]\!] \neq \mathtt{err} \mathbin{?} \mathbf{P}\ (\mathtt{seq}\ \mathcal{S}[\![s]\!]) \circledast \subseteq \rightleftharpoons \mathbf{P}\ \mathcal{S}[\![s]\!] \circledast \subseteq \mathbin{:} \mathtt{err}\big)$
- $\mathcal{T}[\![\rightsquigarrow[s_1, s_2]]\!] \triangleq \big(\mathcal{S}[\![s_1]\!] \neq \mathtt{err} \wedge \mathcal{S}[\![s_2]\!] \neq \mathtt{err} \mathbin{?} \mathbf{P}(\mathcal{S}[\![s_1]\!] * \mathcal{S}[\![s_2]\!]) \circledast \subseteq \rightleftharpoons \mathbf{P}\ \mathcal{S}[\![s_1]\!] \mathbin{\ast\!\!\rightarrow} \mathbf{P}\ \mathcal{S}[\![s_2]\!] \circledast \dot{\subseteq} \mathbin{:} \mathtt{err}\big)$
- $\mathcal{T}[\![\mapsto[s_1, s_2]]\!] \triangleq \big(\mathcal{S}[\![s_1]\!] \neq \mathtt{err} \wedge \mathcal{S}[\![s_2]\!] \neq \mathtt{err} \mathbin{?} \mathbf{P}(\mathcal{S}[\![s_1]\!] \mathbin{\ast\!\!\rightarrow} \mathcal{S}[\![s_2]\!]) \circledast \subseteq \rightleftharpoons \mathbf{P}\ \mathcal{S}[\![s_1]\!] \mathbin{\ast\!\!\rightarrow} \mathbf{P}\ \mathcal{S}[\![s_2]\!] \circledast \dot{\subseteq} \mathbin{:} \mathtt{err}\big)$
- $\mathcal{T}[\![\times[s_1, s_2]]\!] \triangleq \big(\mathcal{S}[\![s_1]\!] \neq \mathtt{err} \wedge \mathcal{S}[\![s_2]\!] \neq \mathtt{err} \mathbin{?} \mathbf{P}(\mathcal{S}[\![s_1]\!] \mathbin{\ast\!\!\rightarrow} \mathcal{S}[\![s_2]\!]) \circledast \subseteq \rightleftharpoons \mathcal{S}[\![s_1]\!] \mathbin{\ast\!\!\rightarrow} \mathbf{P}\ \mathcal{S}[\![s_2]\!] \circledast \dot{\subseteq} \mathbin{:} \mathtt{err}\big)$
- $\mathcal{T}[\![\mathsf{R}[g]]\!] \triangleq \mathcal{T}[\![g]\!]$
- $\mathcal{T}[\![s \twoheadrightarrow g]\!] \triangleq \big(\mathcal{S}[\![s]\!] \neq \mathtt{err} \wedge \mathcal{T}[\![g]\!] \neq \mathtt{err} \mathbin{?} \mathcal{S}[\![s]\!] \mathbin{\ast\!\!\mapsto} \mathcal{T}[\![g]\!] \mathbin{:} \mathtt{err}\big)$

  For brevity, we write $\mathcal{T}[\![s \twoheadrightarrow g]\!] \triangleq \mathcal{S}[\![s]\!] \twoheadrightarrow \mathcal{T}[\![g]\!]$ by defining $\mathsf{S} \twoheadrightarrow \mathsf{T} \triangleq \big(\mathsf{S} \neq \mathtt{err} \wedge \mathsf{T} \neq \mathtt{err} \mathbin{?} \mathsf{S} \mathbin{\ast\!\!\mapsto} \mathsf{T} \mathbin{:} \mathtt{err}\big)$
- $\mathcal{T}[\![g_1 \fatsemi g_2]\!] \triangleq \big(\mathcal{T}[\![g_1]\!] = \mathsf{P}_1 \rightleftharpoons \mathsf{P}_2 \wedge \mathcal{T}[\![g_2]\!] = \mathsf{P}_3 \rightleftharpoons \mathsf{P}_4 \wedge \mathsf{P}_2 \cong \mathsf{P}_3 \mathbin{?} \mathsf{P}_1 \rightleftharpoons \mathsf{P}_4 \mathbin{:} \mathtt{err}\big)$

  For brevity, we write $\mathcal{T}[\![g_1 \fatsemi g_2]\!] = \mathcal{T}[\![g_1]\!] \fatsemi \mathcal{T}[\![g_2]\!]$ by defining $\mathsf{T}_1 \fatsemi \mathsf{T}_2 \triangleq \big(\mathsf{T}_1 = \mathsf{P}_1 \rightleftharpoons \mathsf{P}_2 \wedge \mathsf{T}_2 = \mathsf{P}_3 \rightleftharpoons \mathsf{P}_4 \wedge \mathsf{P}_2 \cong \mathsf{P}_3 \mathbin{?} \mathsf{P}_1 \rightleftharpoons \mathsf{P}_4 \mathbin{:} \mathtt{err}\big)$.
- $\mathcal{T}[\![g_1 \between g_2]\!] \triangleq \big(\mathcal{T}[\![g_1]\!] = \mathsf{S}_1 \circledast \mathsf{O}_1 \rightleftharpoons \mathsf{S}_2 \circledast \mathsf{O}_2 \wedge \mathcal{T}[\![g_2]\!] = \mathsf{S}_3 \circledast \mathsf{O}_3 \rightleftharpoons \mathsf{S}_4 \circledast \mathsf{O}_4 \mathbin{?} \mathsf{S}_1 * \mathsf{S}_2 \circledast \mathsf{O}_1 \star \mathsf{O}_2 \rightleftharpoons \mathsf{S}_3 * \mathsf{S}_4 \circledast \mathsf{O}_3 \star \mathsf{O}_4 \mathbin{:} \mathtt{err}\big)$
- $\mathcal{T}[\![g_1 \Longrightarrow g_2]\!] \triangleq \big(\mathcal{T}[\![g_1]\!] = \mathsf{S}_1 \circledast \mathsf{O}_1 \rightleftharpoons \mathsf{S}_2 \circledast \mathsf{O}_2 \wedge \mathcal{T}[\![g_2]\!] = \mathsf{S}_3 \circledast \mathsf{O}_3 \rightleftharpoons \mathsf{S}_4 \circledast \mathsf{O}_4 \mathbin{?} \mathsf{S}_1 \mathbin{\ast\!\!\rightarrow} \mathsf{S}_3 \circledast \dot{\mathsf{O}}_3 \rightleftharpoons \mathsf{S}_2 \mathbin{\ast\!\!\rightarrow} \mathsf{S}_4 \circledast \dot{\mathsf{O}}_4 \mathbin{:} \mathtt{err}\big)$

For example the type of the *GC* expression describing the interval example of [3, 4, p. 247] is

$$
\begin{aligned}
&\mathcal{T}[\![\cup[(\mathbb{L} \times (\mathbb{X} \mapsto \mathbb{Z}))^\infty] \fatsemi \infty[\mathbb{L} \times (\mathbb{X} \mapsto \mathbb{Z})] \fatsemi \frown[\mathbb{L}, \mathbb{X} \mapsto \mathbb{Z}] \fatsemi \mathbb{L} \twoheadrightarrow (\times[\mathbb{X}, \mathbb{Z}] \fatsemi (\mathbb{X} \twoheadrightarrow \mathbb{I}[\langle \mathbb{Z}, \leq\rangle, -\infty, \infty])))]\!] \\
= \ &\mathcal{T}[\![\cup[(\mathbb{L} \times (\mathbb{X} \mapsto \mathbb{Z}))^\infty] \fatsemi \infty[\mathbb{L} \times (\mathbb{X} \mapsto \mathbb{Z})] \fatsemi \frown[\mathbb{L}, \mathbb{X} \mapsto \mathbb{Z}] \fatsemi \mathbb{L} \twoheadrightarrow (\times[\mathbb{X}, \mathbb{Z}]] \fatsemi (\mathbf{P}\ \mathtt{var} \twoheadrightarrow (\mathbf{P}\ \mathbf{P}\ \mathtt{int} \circledast \subseteq) \rightleftharpoons (\mathbf{P}\ \mathbf{P}\ \mathtt{int} \circledast \subseteq)))]\!] \\
= \ &\mathcal{T}[\![\cup[(\mathbb{L} \times (\mathbb{X} \mapsto \mathbb{Z}))^\infty] \fatsemi \infty[\mathbb{L} \times (\mathbb{X} \mapsto \mathbb{Z})] \fatsemi \frown[\mathbb{L}, \mathbb{X} \mapsto \mathbb{Z}]]\!] \fatsemi \mathbf{P}\ \mathtt{lab} \twoheadrightarrow ((\mathbf{P}(\mathbf{P}\ \mathtt{var} \mathbin{\ast\!\!\rightarrow} \mathbf{P}\ \mathtt{int}) \circledast \subseteq) \rightleftharpoons (\mathbf{P}\ \mathtt{var} \mathbin{\ast\!\!\rightarrow} \mathbf{P}\ \mathbf{P}\ \mathtt{int} \circledast \dot{\subseteq}) \fatsemi (\mathbf{P}\ \mathtt{var} \mathbin{\ast\!\!\rightarrow} \mathbf{P}\ \mathbf{P}\ \mathtt{int} \circledast \dot{\subseteq}) \rightleftharpoons (\mathbf{P}\ \mathtt{var} \mathbin{\ast\!\!\rightarrow} \mathbf{P}\ \mathbf{P}\ \mathtt{int} \circledast \dot{\subseteq})) \\
= \ &\mathcal{T}[\![\cup[(\mathbb{L} \times (\mathbb{X} \mapsto \mathbb{Z}))^\infty] \fatsemi \infty[\mathbb{L} \times (\mathbb{X} \mapsto \mathbb{Z})] \fatsemi \frown[\mathbb{L}, \mathbb{X} \mapsto \mathbb{Z}]]\!] \fatsemi \mathbf{P}\ \mathtt{lab} \twoheadrightarrow ((\mathbf{P}(\mathbf{P}\ \mathtt{var} \mathbin{\ast\!\!\rightarrow} \mathbf{P}\ \mathtt{int}) \circledast \subseteq) \rightleftharpoons (\mathbf{P}\ \mathtt{var} \mathbin{\ast\!\!\rightarrow} \mathbf{P}\ \mathbf{P}\ \mathtt{int} \circledast \dot{\subseteq})) \\
= \ &\mathcal{T}[\![\cup[(\mathbb{L} \times (\mathbb{X} \mapsto \mathbb{Z}))^\infty]]\!] \fatsemi \mathbf{P}\ (\mathbf{P}\ \mathtt{lab} * (\mathbf{P}\ \mathtt{var} \mathbin{\ast\!\!\rightarrow} \mathbf{P}\ \mathtt{int})) \circledast \subseteq \rightleftharpoons \mathbf{P}\ \mathtt{lab} \mathbin{\ast\!\!\rightarrow} \mathbf{P}\ (\mathbf{P}\ \mathtt{var} \mathbin{\ast\!\!\rightarrow} \mathbf{P}\ \mathtt{int}) \circledast \dot{\subseteq} \fatsemi ((\mathbf{P}\ \mathtt{lab} \mathbin{\ast\!\!\rightarrow} \mathbf{P}(\mathbf{P}\ \mathtt{var} \mathbin{\ast\!\!\rightarrow} \mathbf{P}\ \mathtt{int}) \circledast \dot{\subseteq}) \rightleftharpoons (\mathbf{P}\ \mathtt{lab} \mathbin{\ast\!\!\rightarrow} \mathbf{P}\ \mathtt{var} \mathbin{\ast\!\!\rightarrow} \mathbf{P}\ \mathbf{P}\ \mathtt{int} \circledast \ddot{\subseteq})) \\
= \ &\mathcal{T}[\![\cup[(\mathbb{L} \times (\mathbb{X} \mapsto \mathbb{Z}))^\infty]]\!] \fatsemi \mathbf{P}\ (\mathbf{P}\ (\mathtt{seq}\ (\mathbf{P}\ \mathtt{lab} * (\mathbf{P}\ \mathtt{var} \mathbin{\ast\!\!\rightarrow} \mathbf{P}\ \mathtt{int})))) \circledast \subseteq \rightleftharpoons \mathbf{P}\ (\mathbf{P}\ \mathtt{lab} * (\mathbf{P}\ \mathtt{var} \mathbin{\ast\!\!\rightarrow} \mathbf{P}\ \mathtt{int})) \circledast \subseteq \fatsemi \mathbf{P}\ (\mathbf{P}\ \mathtt{lab} * (\mathbf{P}\ \mathtt{var} \mathbin{\ast\!\!\rightarrow} \mathbf{P}\ \mathtt{int})) \circledast \subseteq \rightleftharpoons (\mathbf{P}\ \mathtt{lab} \mathbin{\ast\!\!\rightarrow} \mathbf{P}\ \mathtt{var} \mathbin{\ast\!\!\rightarrow} \mathbf{P}\ \mathbf{P}\ \mathtt{int} \circledast \ddot{\subseteq}) \\
= \ &\mathbf{P}\ (\mathbf{P}\ (\mathtt{seq}\ (\mathbf{P}\ \mathtt{lab} * (\mathbf{P}\ \mathtt{var} \mathbin{\ast\!\!\rightarrow} \mathbf{P}\ \mathtt{int})))) \circledast \subseteq \rightleftharpoons \mathbf{P}\ (\mathbf{P}\ (\mathtt{seq}\ (\mathbf{P}\ \mathtt{lab} * (\mathbf{P}\ \mathtt{var} \mathbin{\ast\!\!\rightarrow} \mathbf{P}\ \mathtt{int}))) \circledast \subseteq \fatsemi \mathbf{P}\ (\mathtt{seq}\ (\mathbf{P}\ \mathtt{lab} * (\mathbf{P}\ \mathtt{var} \mathbin{\ast\!\!\rightarrow} \mathbf{P}\ \mathtt{int}))) \circledast \subseteq \rightleftharpoons (\mathbf{P}\ \mathtt{lab} \mathbin{\ast\!\!\rightarrow} \mathbf{P}\ \mathtt{var} \mathbin{\ast\!\!\rightarrow} \mathbf{P}\ \mathbf{P}\ \mathtt{int} \circledast \ddot{\subseteq}) \\
= \ &\mathbf{P}\ (\mathbf{P}\ (\mathtt{seq}\ (\mathbf{P}\ \mathtt{lab} * (\mathbf{P}\ \mathtt{var} \mathbin{\ast\!\!\rightarrow} \mathbf{P}\ \mathtt{int})))) \circledast \subseteq \rightleftharpoons (\mathbf{P}\ \mathtt{lab} \mathbin{\ast\!\!\rightarrow} \mathbf{P}\ \mathtt{var} \mathbin{\ast\!\!\rightarrow} \mathbf{P}\ \mathbf{P}\ \mathtt{int} \circledast \ddot{\subseteq})
\end{aligned}
$$

Intuitively, sets of sets of sequences of states are abstracted by a map of labels to variables to intervals *i.e.* sets of integers (the difference as abstracted away by the type abstraction).

10.1 Soundness of the type inference algorithm

$\forall \mathsf{T} \in \mathfrak{T} \setminus \{\mathtt{err}\} : \Omega \notin \gamma^{\mathfrak{T}}(\mathsf{T})$ so the soundness of type inference is $\big(\mathcal{T}[\![g]\!] \neq \mathtt{err} \mathbin{?} \mathcal{S}[\![g]\!] \in \gamma^{\mathfrak{T}}(\mathcal{T}[\![g]\!]) \cup \{\omega\}\big)$ *i.e.* "typepable

expressions cannot go wrong (*i.e.* return Ω but may return a runtime error ω)". There is no soundness requirement to prove in case of static type error $\mathcal{T}[\![g]\!] = \texttt{err}$ since untypble expressions are rejected. There is also nothing to prove when the semantics returns a dynamic error ω since types say nothing on possible dynamic errors (*e.g.* the type `int` does no prevent overflows).

For example,

$$
\begin{aligned}
& \mathcal{S}[\![\mathbb{I}(\mathbb{Z}, \leq)]\!] \\
\triangleq\ & \big(\mathcal{S}[\![\mathbb{Z}]\!] = S \notin \{\omega, \Omega\} \ ? \ \big(\mathcal{S}[\![\leq]\!] \subseteq S \times S \ ? \ \{\{v \in S \mid \langle v_1,\ v\rangle \in \mathcal{S}[\![\leq]\!] \wedge \langle v,\ v_2\rangle \in \mathcal{S}[\![\leq]\!]\} \mid v_1, v_2 \in S\} : \omega\big) : \mathcal{S}[\![\mathbb{Z}]\!]\big) \\
=\ & \big(\mathcal{S}[\![\leq]\!] \subseteq \mathbb{Z} \times \mathbb{Z} \ ? \ \{\{v \in \mathbb{Z} \mid \langle v_1,\ v\rangle \in \mathcal{S}[\![\leq]\!] \wedge \langle v,\ v_2\rangle \in \mathcal{S}[\![\leq]\!]\} \mid v_1, v_2 \in \mathbb{Z}\} : \omega\big) \\
=\ & \{\{v \in \mathbb{Z} \mid \langle v_1,\ v\rangle \in \mathcal{S}[\![\leq]\!] \wedge \langle v,\ v_2\rangle \in \mathcal{S}[\![\leq]\!]\} \mid v_1, v_2 \in \mathbb{Z}\} \\
=\ & \{\{v \in \mathbb{Z} \mid v_1 \leq v \langle v,\ \leq \rangle v_2\} \mid v_1, v_2 \in \mathbb{Z}\} \\
=\ & \{[v_1, v_2] \mid v_1, v_2 \in \mathbb{Z}\} \\
\in\ & \{\wp(X) \mid X \in \wp(\mathbb{Z} \cup \{-\infty, \infty\})\} \\
=\ & \{\wp(X) \mid X \in \wp(\gamma^{\mathfrak{e}}(\texttt{int}))\} \\
=\ & \{\wp(X) \mid X \in \gamma^{\mathfrak{S}}(\texttt{P int})\} \\
\subseteq\ & \gamma^{\mathfrak{S}}(\texttt{P P int}) \\
=\ & \gamma^{\mathfrak{S}}(\texttt{P}\ \mathcal{S}[\![\mathbb{Z}]\!]) \\
\subseteq\ & \gamma^{\mathfrak{S}}(\texttt{P}\ \mathcal{S}[\![\mathbb{Z}]\!]) \cup \{\omega\}
\end{aligned}
$$

10.2 Soundness of element type inference

- $\mathscr{E}[\![\mathsf{true}]\!] = \mathscr{E}[\![\mathsf{false}]\!] \triangleq \texttt{bool}$
  $\mathcal{S}[\![\mathsf{true}]\!] = \mathsf{true} \in \mathbb{B} = \gamma^{\mathfrak{e}}(\texttt{bool}) = \gamma^{\mathfrak{e}}(\mathscr{E}[\![\mathsf{true}]\!])$, *etc.*
- $\mathscr{E}[\![0]\!] = \mathscr{E}[\![1]\!] = \ldots = \mathscr{E}[\![\infty]\!] \triangleq \texttt{int}$,
  $\mathcal{S}[\![0]\!] = 0 \in \mathbb{Z} \cup \{-\infty, \infty\} = \gamma^{\mathfrak{e}}(\texttt{int}) = \gamma^{\mathfrak{e}}(\mathscr{E}[\![0]\!])$, *etc.*
- $\mathscr{E}[\![\mathsf{x}]\!] = \mathscr{E}[\![\mathsf{y}]\!] = \ldots \triangleq \texttt{var}$
  $\mathcal{S}[\![\mathsf{x}]\!] = \mathsf{x} \in \mathbb{X} = \gamma^{\mathfrak{e}}(\texttt{var}) = \gamma^{\mathfrak{e}}(\mathscr{E}[\![\mathsf{x}]\!])$, *etc.*
- $\mathscr{E}[\![\ell]\!] = \ldots \triangleq \texttt{lab}$
  $\mathcal{S}[\![\ell]\!] = \ell \in \mathbb{L} = \gamma^{\mathfrak{e}}(\texttt{lab}) = \gamma^{\mathfrak{e}}(\mathscr{E}[\![\ell]\!])$,
- $\mathscr{E}[\![-e]\!] = \big(\mathscr{E}[\![e]\!] = \texttt{bool} \vee \mathscr{E}[\![e]\!] = \texttt{int} \ ? \ \mathscr{E}[\![e]\!] : \texttt{err}\big)$
  - if $\mathscr{E}[\![e]\!] = \texttt{bool}$ then $\mathcal{S}[\![e]\!] \in \gamma^{\mathfrak{e}}(\mathscr{E}[\![e]\!]) = \gamma^{\mathfrak{e}}(\texttt{bool}) = \mathbb{B}$ by structural induction hypothesis so $\mathcal{S}[\![-e]\!] = \neg\mathcal{S}[\![e]\!] \in \mathbb{B} = \gamma^{\mathfrak{e}}(\texttt{bool}) = \gamma^{\mathfrak{e}}(\mathscr{E}[\![-e]\!])$
  - else if $\mathscr{E}[\![e]\!] = \texttt{int}$, *Id.*
  - else $\mathscr{E}[\![e]\!] = \texttt{err}$, no requirement to prove.
- Observe that when $\mathcal{O}[\![e]\!] \neq \Omega$, $\mathcal{O}[\![e]\!]$ has been defined in Sect. 3.1 to be a poset of the form $\mathcal{O}[\![e]\!] = \langle \gamma^{\mathfrak{e}}(\mathscr{E}[\![e]\!]),\ o\rangle$

10.3 Soundness of set type inference

- $\mathcal{S}[\![\mathbb{B}]\!] \triangleq \texttt{P bool}$
  $\mathcal{S}[\![\mathbb{B}]\!] = \mathbb{B} \in \wp(\mathbb{B}) = \wp(\gamma^{\mathfrak{e}}(\texttt{bool})) = \gamma^{\mathfrak{S}}(\texttt{P bool}) = \gamma^{\mathfrak{S}}(\mathcal{S}[\![\mathbb{B}]\!])$ *QED.*

- $\mathcal{S}[\![\mathbb{Z}]\!] \triangleq \texttt{P int}$
  Id.

- $\mathcal{S}[\![\mathbb{X}]\!] \triangleq \texttt{P var}$
  Id.

- $\mathcal{S}[\![\mathbb{L}]\!] \triangleq \texttt{P lab}$
  Id.

- $\mathcal{S}[\![\{e\}]\!] \triangleq \big(\mathscr{E}[\![e]\!] \neq \texttt{err} \ ? \ \texttt{P}\ \mathscr{E}[\![e]\!] : \texttt{err}\big)$

  By structural induction hypothesis, $\mathcal{S}[\![e]\!] \in \gamma^{\mathfrak{e}}(\mathscr{E}[\![e]\!])$ so $\forall \mathsf{E} \in \mathfrak{E} \setminus \{\texttt{err}\} : \Omega \notin \gamma^{\mathfrak{e}}(\mathsf{E})$ and $\mathscr{E}[\![e]\!] \neq \texttt{err}$ imply that $\mathcal{S}[\![e]\!] \neq \Omega$ so $\mathcal{S}[\![\{e\}]\!] = \big(\mathcal{S}[\![e]\!] \neq \Omega \ ? \ \{\mathcal{S}[\![e]\!]\} : \Omega\big) = \{\mathcal{S}[\![e]\!]\} \in \wp(\gamma^{\mathfrak{e}}(\mathscr{E}[\![e]\!])) = \gamma^{\mathfrak{S}}(\texttt{P}\ \mathscr{E}[\![e]\!]) = \gamma^{\mathfrak{S}}(\mathcal{S}[\![\{e\}]\!])$

- $\mathcal{S}[\![[e_1, e_2]]\!] \triangleq \big(\mathscr{E}[\![e_1]\!] \cong \mathscr{E}[\![e_2]\!] \neq \texttt{err} \ ? \ \texttt{P}\ \mathscr{E}[\![e_1]\!] : \texttt{err}\big)$

  No requirement has to be proved when $\mathcal{S}[\![[e_1, e_2]]\!] = \texttt{err}$. Otherwise $\mathscr{E}[\![e_1]\!] \cong \mathscr{E}[\![e_2]\!] \neq \texttt{err}$ and $\mathcal{S}[\![[e_1, e_2]]\!] = \texttt{P}\ \mathscr{E}[\![e_1]\!]$, in which case we must prove that $\mathcal{S}[\![[e_1, e_2]]\!] \in \gamma^{\mathfrak{S}}(\mathcal{S}[\![[e_1, e_2]]\!])$.

  $\mathscr{E}[\![e_1]\!] \cong \mathscr{E}[\![e_2]\!]$ implies $\gamma^{\mathfrak{e}}(\mathscr{E}[\![e_1]\!]) = \gamma^{\mathfrak{e}}(\mathscr{E}[\![e_2]\!])$ (and reciprocally). But $\gamma^{\mathfrak{e}}$ is injective so $\mathscr{E}[\![e_1]\!] = \mathscr{E}[\![e_2]\!] \neq \texttt{err}$. Morover, $\forall \mathsf{E} \in \mathfrak{E} \setminus \{\texttt{err}\} : \Omega \notin \gamma^{\mathfrak{e}}(\mathsf{E})$ and $\mathscr{E}[\![e_i]\!] \neq \texttt{err}$ imply that $\Omega \notin \gamma^{\mathfrak{e}}(\mathscr{E}[\![e_i]\!])$, $i = 1, 2$. By structural induction hypothesis, $\mathcal{S}[\![e_i]\!] \in \gamma^{\mathfrak{e}}(\mathscr{E}[\![e_i]\!])$, $i = 1, 2$, so that $\mathcal{S}[\![e_i]\!] \neq \Omega$.

  - In case $\mathcal{O}[\![e_1]\!] = \mathcal{O}[\![e_2]\!] = \langle S,\ o\rangle$, $\mathcal{O}[\![e_i]\!] \neq \Omega$ so $\mathcal{O}[\![e_i]\!] = \langle \gamma^{\mathfrak{e}}(\mathscr{E}[\![e_i]\!]),\ o_i\rangle$ implies that $S = \gamma^{\mathfrak{e}}(\mathscr{E}[\![e_i]\!])$. It follows that $\mathcal{S}[\![[e_1, e_2]]\!] \triangleq \big(\mathcal{O}[\![e_1]\!] = \mathcal{O}[\![e_2]\!] = \langle S,\ o\rangle \ ? \ \{v \in S \mid \langle \mathcal{S}[\![e_1]\!],\ v\rangle \in o \wedge \langle v,\ \mathcal{S}[\![e_2]\!]\rangle \in 0\} : \Omega\big) = \{v \in \gamma^{\mathfrak{e}}(\mathscr{E}[\![e_1]\!]) \mid \langle \mathcal{S}[\![e_1]\!],\ v\rangle \in o \wedge \langle v,\ \mathcal{S}[\![e_2]\!]\rangle \in 0\} \in \wp(\gamma^{\mathfrak{e}}(\mathscr{E}[\![e_1]\!])) \triangleq \gamma^{\mathfrak{S}}(\texttt{P}\ \mathscr{E}[\![e_1]\!]) = \gamma^{\mathfrak{S}}(\mathcal{S}[\![[e_1, e_2]]\!])$, *QED.*
  - In case $\mathcal{O}[\![e_1]\!] = \langle S_1,\ o_1\rangle \neq \mathcal{O}[\![e_2]\!] = \langle S_2,\ o_2\rangle$, we have $\mathcal{O}[\![e_i]\!] \neq \Omega$ so $\mathcal{O}[\![e_i]\!] = \langle \gamma^{\mathfrak{e}}(\mathscr{E}[\![e_i]\!]),\ o_i\rangle$. Therefore $\mathscr{E}[\![e_1]\!] = \mathscr{E}[\![e_2]\!]$ implies $S_1 = S_2$. It follows that $o_1 \neq o_2$, which is in contradiction with the definition of $\mathcal{O}$ in Sect. 3.1. By reductio ad absurdum, this case is impossible.
  - The last possible case is $\mathcal{O}[\![e_1]\!] = \Omega$ or $\mathcal{O}[\![e_2]\!] = \Omega$. By def. of $\mathcal{O}$ in Sect. 3.1, one of e_1 or e_2 is of the form $e_i = -e_i'$ with $\mathcal{S}[\![e_i']\!] \notin \mathbb{B} \wedge \mathcal{S}[\![e_i']\!] \notin \mathbb{Z} \cup \{-\infty, \infty\}$, which implies that $\mathcal{S}[\![e_i]\!] = \Omega$, a contradiction. By reductio ad absurdum, this case is also impossible.

  In conclusion, $\mathcal{S}[\![[e_1, e_2]]\!] \in \gamma^{\mathfrak{S}}(\mathcal{S}[\![[e_1, e_2]]\!])$. *QED.*

- $\mathcal{S}[\![\mathbb{I}(s, o)]\!] \triangleq \big(\mathcal{S}[\![s]\!] \neq \texttt{err} \ ? \ \texttt{P}\ \mathcal{S}[\![s]\!] : \texttt{err}\big)$

  In case $\mathcal{S}[\![\mathbb{I}(s, o)]\!] \neq \texttt{err}$, $\mathcal{S}[\![s]\!] \neq \texttt{err}$ so, by structural induction hypothesis, $\mathcal{S}[\![s]\!] \in (\gamma^{\mathfrak{S}}(\mathcal{S}[\![s]\!]) \cup \{\omega\}) \setminus \{\Omega\}$.
  - In case, $\mathcal{S}[\![s]\!] = \omega$, we have $\mathcal{S}[\![\mathbb{I}(s, o)]\!] \triangleq \big(\mathcal{S}[\![s]\!] = S \notin \{\omega, \Omega\} \ ? \ \big(\mathcal{S}[\![o]\!] \subseteq S \times S \ ? \ \{\{v \in S \mid \langle v_1,\ v\rangle \in \mathcal{S}[\![o]\!] \wedge \langle v,\ v_2\rangle \in \mathcal{S}[\![o]\!]\} \mid v_1, v_2 \in S\} : \omega\big) : \mathcal{S}[\![s]\!]\big) = \mathcal{S}[\![s]\!] = \omega$ in which case $\mathcal{S}[\![\mathbb{I}(s, o)]\!] \in \{\omega\}$
  - Otherwise, in case $\mathcal{S}[\![s]\!] \in \gamma^{\mathfrak{S}}(\mathcal{S}[\![s]\!])$ and $\mathcal{S}[\![o]\!] \notin S \times S$, $\mathcal{S}[\![\mathbb{I}(s, o)]\!] \triangleq \mathcal{S}[\![\mathbb{I}(s, o)]\!] \triangleq \big(\mathcal{S}[\![s]\!] = S \notin \{\omega, \Omega\} \ ? \ \big(\mathcal{S}[\![o]\!] \subseteq S \times S \ ? \ \{\{v \in S \mid \langle v_1,\ v\rangle \in \mathcal{S}[\![o]\!] \wedge \langle v,\ v_2\rangle \in \mathcal{S}[\![o]\!]\} \mid v_1, v_2 \in S\} : \omega\big) : \mathcal{S}[\![s]\!]\big) = \omega$ so $\mathcal{S}[\![\mathbb{I}(s, o)]\!] \in \{\omega\}$
  - Otherwise $\mathcal{S}[\![s]\!] \in \gamma^{\mathfrak{S}}(\mathcal{S}[\![s]\!])$ and $\mathcal{S}[\![o]\!] \subseteq S \times S$ and so $\mathcal{S}[\![\mathbb{I}(s, o)]\!] \triangleq \big(\mathcal{S}[\![s]\!] = S \notin \{\omega, \Omega\} \ ? \ \big(\mathcal{S}[\![o]\!] \subseteq S \times S \ ? \ \{\{v \in S \mid \langle v_1,\ v\rangle \in \mathcal{S}[\![o]\!] \wedge \langle v,\ v_2\rangle \in \mathcal{S}[\![o]\!]\} \mid v_1, v_2 \in S\} : \omega\big) : \mathcal{S}[\![s]\!]\big) = \{\{v \in \mathcal{S}[\![s]\!] \mid \langle v_1,\ v\rangle \in \mathcal{S}[\![o]\!] \wedge \langle v,\ v_2\rangle \in \mathcal{S}[\![o]\!]\} \mid v_1, v_2 \in \mathcal{S}[\![s]\!]\} \in \{\wp(X) \mid X \in \gamma^{\mathfrak{S}}(\mathcal{S}[\![s]\!])\} \subseteq \gamma^{\mathfrak{S}}(\texttt{P}\ \mathcal{S}[\![s]\!])$
  - In all cases, we have $\mathcal{S}[\![\mathbb{I}(s, o)]\!] \in \gamma^{\mathfrak{S}}(\texttt{P}\ \mathcal{S}[\![s]\!]) \cup \{\omega\}$. *QED.*

- $\mathcal{S}[\![s^\infty]\!] \triangleq \big(\mathcal{S}[\![s]\!] \neq \texttt{err} \ ? \ \texttt{seq}\ \mathcal{S}[\![s]\!] : \texttt{err}\big)$

  In case $\mathcal{S}[\![s^\infty]\!] \neq \texttt{err}$, $\mathcal{S}[\![s]\!] \neq \texttt{err}$ so, by structural induction hypothesis, $\mathcal{S}[\![s]\!] \in (\gamma^{\mathfrak{S}}(\mathcal{S}[\![s]\!]) \cup \{\omega\}) \setminus \{\Omega\}$.

- The case $\mathcal{S}[\![s]\!] = \Omega$ is impossible since $\mathcal{S}[\![s^\infty]\!] \neq \texttt{err}$
- In case $\mathcal{S}[\![s]\!] = \omega$, $\mathcal{S}[\![s^\infty]\!] \triangleq \left(\mathcal{S}[\![s]\!] \notin \{\omega, \Omega\} \;?\; \mathcal{S}[\![s]\!]^\infty \;:\; \mathcal{S}[\![s]\!]\right) = \omega \in \{\omega\}$
- Otherwise, $\mathcal{S}[\![s]\!] \notin \{\omega, \Omega\}$ so $\mathcal{S}[\![s]\!] \in \gamma^{\mathfrak{S}}(\mathcal{S}[\![s]\!])$ and $\mathcal{S}[\![s^\infty]\!] \triangleq \left(\mathcal{S}[\![s]\!] \notin \{\omega, \Omega\} \;?\; \mathcal{S}[\![s]\!]^\infty \;:\; \mathcal{S}[\![s]\!]\right) = \mathcal{S}[\![s]\!]^\infty \in \{X^\infty \mid X \in \gamma^{\mathfrak{S}}(\mathcal{S}[\![s]\!])\} \triangleq \gamma^{\mathfrak{S}}(\texttt{seq}\, \mathcal{S}[\![s]\!]) = \gamma^{\mathfrak{S}}(\mathcal{S}[\![s^\infty]\!])$.
We conclude that $\mathcal{S}[\![s^\infty]\!] \neq \Omega$ implies that $\mathcal{S}[\![s^\infty]\!] \in \gamma^{\mathfrak{S}}(\mathcal{S}[\![s^\infty]\!]) \cup \{\omega\}$. *QED.*

- $\mathcal{S}[\![s_1 \cup s_2]\!] \triangleq \left(\texttt{err} \neq \mathcal{S}[\![s_1]\!] \cong \mathcal{S}[\![s_2]\!] \neq \texttt{err} \;?\; \mathcal{S}[\![s_1]\!] \;:\; \texttt{err}\right)$

  In case $\mathcal{S}[\![s_1 \cup s_2]\!] \neq \texttt{err}$, we have $\mathcal{S}[\![s_1]\!] \neq \texttt{err}$ and $\mathcal{S}[\![s_2]\!] \neq \texttt{err}$ so $\mathcal{S}[\![s_1 \cup s_2]\!] = \mathcal{S}[\![s_1]\!]$ and, by structural induction hypothesis, $\mathcal{S}[\![s_1]\!] \in (\gamma^{\mathfrak{S}}(\mathcal{S}[\![s_1]\!]) \cup \{\omega\}) \setminus \{\Omega\} = (\gamma^{\mathfrak{S}}(\mathcal{S}[\![s_2]\!]) \cup \{\omega\}) \setminus \{\Omega\} \ni \mathcal{S}[\![s_2]\!]$.
  - In case $\mathcal{S}[\![s_1]\!] \in \{\omega, \Omega\} \vee \mathcal{S}[\![s_2]\!] \in \{\omega, \Omega\}$, we have $\mathcal{S}[\![s_1]\!] = \omega \vee \mathcal{S}[\![s_2]\!] = \omega$ so $\mathcal{S}[\![s_1 \cup s_2]\!] \triangleq \left(\mathcal{S}[\![s_1]\!] \notin \{\omega, \Omega\} \wedge \mathcal{S}[\![s_2]\!] \notin \{\omega, \Omega\} \;?\; \mathcal{S}[\![s_1]\!] \cup \mathcal{S}[\![s_2]\!] \;:\; \texttt{error}(\mathcal{S}[\![s_1]\!], \mathcal{S}[\![s_2]\!])\right) = \texttt{error}(\mathcal{S}[\![s_1]\!], \mathcal{S}[\![s_2]\!]) = \omega \in \{\omega\}$
  - Otherwise, $\mathcal{S}[\![s_1]\!] \neq \omega \wedge \mathcal{S}[\![s_2]\!] \neq \omega$ so $\mathcal{S}[\![s_1 \cup s_2]\!] \triangleq \left(\mathcal{S}[\![s_1]\!] \notin \{\omega, \Omega\} \wedge \mathcal{S}[\![s_2]\!] \notin \{\omega, \Omega\} \;?\; \mathcal{S}[\![s_1]\!] \cup \mathcal{S}[\![s_2]\!] \;:\; \texttt{error}(\mathcal{S}[\![s_1]\!], \mathcal{S}[\![s_2]\!])\right) = \mathcal{S}[\![s_1]\!] \cup \mathcal{S}[\![s_2]\!] \in \gamma^{\mathfrak{S}}(\mathcal{S}[\![s_1]\!]) \cup \gamma^{\mathfrak{S}}(\mathcal{S}[\![s_2]\!]) = \gamma^{\mathfrak{S}}(\mathcal{S}[\![s_1]\!]) = \gamma^{\mathfrak{S}}(\mathcal{S}[\![s_1 \cup s_2]\!])$
  - In both cases, $\mathcal{S}[\![s_1 \cup s_2]\!] \in \gamma^{\mathfrak{S}}(\mathcal{S}[\![s_1 \cup s_2]\!]) \cup \{\omega\}$. *QED.*

- $\mathcal{S}[\![s_1 \mapsto s_2]\!] \triangleq \left(\mathcal{S}[\![s_1]\!] \neq \texttt{err} \wedge \mathcal{S}[\![s_2]\!] \neq \texttt{err} \;?\; \mathcal{S}[\![s_1]\!] \mathrel{*\!\!\to} \mathcal{S}[\![s_2]\!] \;:\; \texttt{err}\right)$

  In case $\mathcal{S}[\![s_1 \mapsto s_2]\!] \neq \texttt{err}$, we have $\mathcal{S}[\![s_1]\!] \neq \texttt{err}$ and $\mathcal{S}[\![s_2]\!] \neq \texttt{err}$ so $\mathcal{S}[\![s_1 \mapsto s_2]\!] = \mathcal{S}[\![s_1]\!] \mathrel{*\!\!\to} \mathcal{S}[\![s_2]\!]$ and, by structural induction hypothesis, $\mathcal{S}[\![s_1]\!] \in (\gamma^{\mathfrak{S}}(\mathcal{S}[\![s_1]\!]) \cup \{\omega\}) \setminus \{\Omega\}$ and $\mathcal{S}[\![s_2]\!] \in (\gamma^{\mathfrak{S}}(\mathcal{S}[\![s_2]\!]) \cup \{\omega\}) \setminus \{\Omega\}$.
  - If $\mathcal{S}[\![s_1]\!] \in \{\omega, \Omega\}$ or $\mathcal{S}[\![s_2]\!] \in \{\omega, \Omega\}$ then $\mathcal{S}[\![s_1]\!] = \omega$ or $\mathcal{S}[\![s_2]\!] = \omega$ in which case $\mathcal{S}[\![s_1 \mapsto s_2]\!] \triangleq \left(\mathcal{S}[\![s_1]\!] \notin \{\omega, \Omega\} \wedge \mathcal{S}[\![s_2]\!] \notin \{\omega, \Omega\} \;?\; \mathcal{S}[\![s_1]\!] \mapsto \mathcal{S}[\![s_2]\!] \;:\; \texttt{error}(\mathcal{S}[\![s_1]\!], \mathcal{S}[\![s_2]\!])\right) = \texttt{error}(\mathcal{S}[\![s_1]\!], \mathcal{S}[\![s_2]\!]) = \omega \in \{\omega\}$
  - Otherwise $\mathcal{S}[\![s_1]\!] \in \gamma^{\mathfrak{S}}(\mathcal{S}[\![s_1]\!])$, $\mathcal{S}[\![s_2]\!] \in \gamma^{\mathfrak{S}}(\mathcal{S}[\![s_2]\!])$, and $\mathcal{S}[\![s_1 \mapsto s_2]\!] \triangleq \left(\mathcal{S}[\![s_1]\!] \notin \{\omega, \Omega\} \wedge \mathcal{S}[\![s_2]\!] \notin \{\omega, \Omega\} \;?\; \mathcal{S}[\![s_1]\!] \mapsto \mathcal{S}[\![s_2]\!] \;:\; \texttt{error}(\mathcal{S}[\![s_1]\!], \mathcal{S}[\![s_2]\!])\right) = \mathcal{S}[\![s_1]\!] \mapsto \mathcal{S}[\![s_2]\!] \in \{X \mapsto Y \mid X \in \gamma^{\mathfrak{S}}(\mathcal{S}[\![s_1]\!]) \wedge Y \in \gamma^{\mathfrak{S}}(\mathcal{S}[\![s_2]\!])\} \triangleq \gamma^{\mathfrak{S}}(\mathcal{S}[\![s_1]\!] \mathrel{*\!\!\to} \mathcal{S}[\![s_2]\!]) \triangleq \gamma^{\mathfrak{S}}(\mathcal{S}[\![s_1 \mapsto s_2]\!])$.
  - Grouping both cases together, we have $\mathcal{S}[\![s_1 \mapsto s_2]\!] \in \gamma^{\mathfrak{S}}(\mathcal{S}[\![s_1 \mapsto s_2]\!]) \cup \{\omega\}$.

- $\mathcal{S}[\![s_1 \times s_2]\!] \triangleq \left(\mathcal{S}[\![s_1]\!] \neq \texttt{err} \wedge \mathcal{S}[\![s_2]\!] \neq \texttt{err} \;?\; \mathcal{S}[\![s_1]\!] * \mathcal{S}[\![s_2]\!] \;:\; \texttt{err}\right)$

  In case $\mathcal{S}[\![s_1 \times s_2]\!] \neq \texttt{err}$, we have $\mathcal{S}[\![s_1]\!] \neq \texttt{err}$ and $\mathcal{S}[\![s_2]\!] \neq \texttt{err}$ so $\mathcal{S}[\![s_1 \times s_2]\!] = \mathcal{S}[\![s_1]\!] * \mathcal{S}[\![s_2]\!]$ and, by structural induction hypothesis, $\mathcal{S}[\![s_1]\!] \in (\gamma^{\mathfrak{S}}(\mathcal{S}[\![s_1]\!]) \cup \{\omega\}) \setminus \{\Omega\}$ and $\mathcal{S}[\![s_2]\!] \in (\gamma^{\mathfrak{S}}(\mathcal{S}[\![s_2]\!]) \cup \{\omega\}) \setminus \{\Omega\}$.
  - If $\mathcal{S}[\![s_1]\!] \in \{\omega, \Omega\}$ or $\mathcal{S}[\![s_2]\!] \in \{\omega, \Omega\}$ then $\mathcal{S}[\![s_1]\!] = \omega$ or $\mathcal{S}[\![s_2]\!] = \omega$ in which case $\mathcal{S}[\![s_1 \times s_2]\!] \triangleq \mathcal{S}[\![s_1 \times s_2]\!] \triangleq \left(\mathcal{S}[\![s_1]\!] \notin \{\omega, \Omega\} \wedge \mathcal{S}[\![s_2]\!] \notin \{\omega, \Omega\} \;?\; \mathcal{S}[\![s_1]\!] \times \mathcal{S}[\![s_2]\!] \;:\; \texttt{error}(\mathcal{S}[\![s_1]\!], \mathcal{S}[\![s_2]\!])\right) = \texttt{error}(\mathcal{S}[\![s_1]\!], \mathcal{S}[\![s_2]\!]) = \omega \in \{\omega\}$
  - Otherwise $\mathcal{S}[\![s_1]\!] \in \gamma^{\mathfrak{S}}(\mathcal{S}[\![s_1]\!])$, $\mathcal{S}[\![s_2]\!] \in \gamma^{\mathfrak{S}}(\mathcal{S}[\![s_2]\!])$, and $\mathcal{S}[\![s_1 \times s_2]\!] \triangleq \mathcal{S}[\![s_1 \times s_2]\!] \triangleq \left(\mathcal{S}[\![s_1]\!] \notin \{\omega, \Omega\} \wedge \mathcal{S}[\![s_2]\!] \notin \{\omega, \Omega\} \;?\; \mathcal{S}[\![s_1]\!] \times \mathcal{S}[\![s_2]\!] \;:\; \texttt{error}(\mathcal{S}[\![s_1]\!], \mathcal{S}[\![s_2]\!])\right) = \mathcal{S}[\![s_1]\!] \times \mathcal{S}[\![s_2]\!] \in \{X \times Y \mid X \in \gamma^{\mathfrak{S}}(\mathcal{S}[\![s_1]\!]) \wedge Y \in \gamma^{\mathfrak{S}}(\mathcal{S}[\![s_2]\!])\} \triangleq \gamma^{\mathfrak{S}}(\mathcal{S}[\![s_1]\!] * \mathcal{S}[\![s_2]\!]) \triangleq \gamma^{\mathfrak{S}}(\mathcal{S}[\![s_1 \times s_2]\!])$.
  - Grouping both cases together, we have $\mathcal{S}[\![s_1 \times s_2]\!] \in \gamma^{\mathfrak{S}}(\mathcal{S}[\![s_1 \times s_2]\!]) \cup \{\omega\}$.

- $\mathcal{S}[\![\wp(s)]\!] \triangleq \left(\mathcal{S}[\![s]\!] \neq \texttt{err} \;?\; \texttt{P}\; \mathcal{S}[\![s]\!] \;:\; \texttt{err}\right)$

  In case $\mathcal{S}[\![\wp(s)]\!] \neq \texttt{err}$, we have $\mathcal{S}[\![s]\!] \neq \texttt{err}$, $\mathcal{S}[\![\wp(s)]\!] = \texttt{P}\; \mathcal{S}[\![s]\!]$ and, by structural induction hypothesis, $\mathcal{S}[\![s]\!] \in (\gamma^{\mathfrak{S}}(\mathcal{S}[\![s]\!]) \cup \{\omega\}) \setminus \{\Omega\}$.
  - If $\mathcal{S}[\![s]\!] \in \{\omega, \Omega\}$ then $\mathcal{S}[\![s]\!] = \omega$ in which case $\mathcal{S}[\![\wp(s)]\!] \triangleq \left(\mathcal{S}[\![s]\!] \notin \{\omega, \Omega\} \;?\; \wp(\mathcal{S}[\![s]\!]) \;:\; \mathcal{S}[\![s]\!]\right) = \mathcal{S}[\![s]\!] = \omega \in \{\omega\}$
  - Otherwise $\mathcal{S}[\![s]\!] \in \gamma^{\mathfrak{S}}(\mathcal{S}[\![s]\!])$ and $\mathcal{S}[\![\wp(s)]\!] \triangleq \left(\mathcal{S}[\![s]\!] \notin \{\omega, \Omega\} \;?\; \wp(\mathcal{S}[\![s]\!]) \;:\; \mathcal{S}[\![s]\!]\right) = \wp(\mathcal{S}[\![s]\!]) \in \{\wp(X) \mid X \in \gamma^{\mathfrak{S}}(\mathcal{S}[\![s]\!])\} \subseteq \gamma^{\mathfrak{S}}(\texttt{P}\; \mathcal{S}[\![s]\!]) \triangleq \gamma^{\mathfrak{S}}(\mathcal{S}[\![\wp(s)]\!])$.
  - In both cases, we have $\mathcal{S}[\![\wp(s)]\!] \in \gamma^{\mathfrak{S}}(\mathcal{S}[\![\wp(s)]\!]) \cup \{\omega\}$.
- $\mathcal{S}[\![x]\!] \triangleq \texttt{err}$, otherwise (nothing to prove).

10.4 Partial order type inference

- $\mathcal{O}[\![o]\!] \triangleq o, \quad o \in \{\Rightarrow, \Leftrightarrow, \leq, \subseteq, =\}$
- $\mathcal{O}[\![\Subset]\!] \triangleq \subseteq$
- $\mathcal{O}[\![o^{-1}]\!] \triangleq (\mathcal{O}[\![o]\!])^{-1}$
- $\mathcal{O}[\![o_1 \times o_2]\!] \triangleq \mathcal{O}[\![o_1]\!] * \mathcal{O}[\![o_2]\!]$
- $\mathcal{O}[\![\dot{o}]\!] \triangleq (\dot{\mathcal{O}[\![o]\!]})$
- $\mathcal{O}[\![\ddot{o}]\!] \triangleq (\dddot{\mathcal{O}[\![o]\!]})$

  Since $\mathcal{O}[\![o]\!] \neq \texttt{err}$, soundness is $\mathcal{S}[\![o]\!] \in \gamma^{\mathfrak{O}}(\mathcal{O}[\![o]\!])$ (which implies $\mathcal{S}[\![o]\!] \in \gamma^{\mathfrak{O}}(\mathcal{O}[\![o]\!]) \cup \{\omega\}$). By the definition of $\gamma^{\mathfrak{O}}$ in Sec. 6.3, we have $\gamma^{\mathfrak{O}}(\mathsf{O}) \triangleq \{\mathcal{S}[\![\mathsf{O}]\!]\}$ for all $\mathsf{O} \in \mathfrak{O}$ (up to the isomorphism between partial order types and the semantics of partial orders, up to the use of $*$ for the type of pairs $\times$ and $\subseteq$ for interval inclusion $\Subset$) and therefore $\mathcal{S}[\![o]\!] \in \{\mathcal{S}[\![o]\!]\} = \gamma^{\mathfrak{O}}(o) = \gamma^{\mathfrak{O}}(\mathcal{O}[\![o]\!])$.
- $\mathcal{O}[\![x]\!] \triangleq \Omega$, otherwise (nothing to prove).

10.5 Poset type inference

- $\mathcal{P}[\![\langle s,\ o\rangle]\!] \triangleq \left(\mathcal{S}[\![s]\!] \neq \texttt{err} \wedge \mathcal{O}[\![o]\!] \neq \texttt{err} \;?\; \mathcal{S}[\![s]\!] \circledast \mathcal{O}[\![o]\!] \;:\; \texttt{err}\right)$

  If $\mathcal{P}[\![\langle s,\ o\rangle]\!] \neq \texttt{err}$ then $\mathcal{S}[\![s]\!] \neq \texttt{err}$ and $\mathcal{O}[\![o]\!] \neq \texttt{err}$ so by, structural induction hypothesis, $\mathcal{S}[\![s]\!] \in \gamma^{\mathfrak{S}}(\mathcal{S}[\![s]\!]) \cup \{\omega\}$ and $\mathcal{S}[\![o]\!] \in \gamma^{\mathfrak{O}}(\mathcal{O}[\![o]\!]) \cup \{\omega\}$. There is nothing to prove in case of a dynamic error ω so, in the remaining case, $\mathcal{S}[\![\langle s,\ o\rangle]\!] \triangleq \left(\mathcal{S}[\![s]\!] \notin \{\omega, \Omega\} \wedge \mathcal{S}[\![o]\!] \notin \{\omega, \Omega\} \;?\; \langle \mathcal{S}[\![s]\!], \mathcal{S}[\![o]\!]\rangle \;:\; \texttt{error}(\mathcal{S}[\![s]\!], \mathcal{S}[\![o]\!])\right) = \langle \mathcal{S}[\![s]\!], \mathcal{S}[\![o]\!]\rangle \in \gamma^{\mathfrak{S}}(\mathcal{S}[\![s]\!]) \times \gamma^{\mathfrak{O}}(\mathcal{O}[\![o]\!]) \triangleq \gamma^{\mathfrak{P}}(\mathcal{S}[\![s]\!] \circledast \mathcal{O}[\![o]\!]) \triangleq \gamma^{\mathfrak{P}}(\mathcal{P}[\![\langle s,\ o\rangle]\!])$

10.6 Galois connection type inference

- $\mathcal{T}[\![\mathbb{1}[p]]\!] \triangleq \left(\mathcal{P}[\![p]\!] \neq \texttt{err} \;?\; \mathcal{P}[\![p]\!] \leftrightharpoons \mathcal{P}[\![p]\!] \;:\; \texttt{err}\right)$

  The soundness proof must be done in the case when $\mathcal{T}[\![\mathbb{1}[p]]\!] \neq \texttt{err}$ which implies $\mathcal{P}[\![p]\!] \neq \texttt{err}$ and so $\mathcal{T}[\![\mathbb{1}[p]]\!] = \mathcal{P}[\![p]\!] \leftrightharpoons \mathcal{P}[\![p]\!]$. By structural induction hypothesis, $\mathcal{S}[\![p]\!] \in \gamma^{\mathfrak{P}}(\mathcal{P}[\![p]\!]) \cup \{\omega\}$.

- If $\mathcal{S}[\![p]\!] = \omega$ then $\mathcal{S}[\![\mathbb{1}[p]]\!] = \left(\mathcal{S}[\![p]\!] \notin \{\omega, \Omega\} \mathbin{?} \mathcal{S}[\![p]\!] \xleftarrow[\lambda P \bullet P]{\lambda Q \bullet Q} \mathcal{S}[\![p]\!] \mathbin{:} \mathcal{S}[\![p]\!]\right) = \mathcal{S}[\![p]\!] = \omega \in \{\omega\}$
- Otherwise $\mathcal{S}[\![p]\!] \notin \{\omega, \Omega\}$. By the syntactic definition of p in Sec. 2.6, p is of the form $p = \langle s,\ o\rangle$. s and o are syntactic components of p, itself a syntactic component of $\mathbb{1}[p]$, so, by structural induction hypothesis, $\mathcal{S}[\![s]\!] \in \gamma^{\mathfrak{S}}(\mathcal{S}[\![s]\!]) \cup \{\omega\}$ and $\mathcal{S}[\![o]\!] \in \gamma^{\mathfrak{O}}(\mathcal{O}[\![o]\!]) \cup \{\omega\}$. But $\mathcal{S}[\![p]\!] = \mathcal{S}[\![\langle s,\ o\rangle]\!] \triangleq \left(\mathcal{S}[\![s]\!] \notin \{\omega, \Omega\} \wedge \mathcal{S}[\![o]\!] \notin \{\omega, \Omega\} \mathbin{?} \langle \mathcal{S}[\![s]\!],\ \mathcal{S}[\![o]\!]\rangle \mathbin{:} \texttt{error}(\mathcal{S}[\![s]\!], \mathcal{S}[\![o]\!])\right)$ so $\mathcal{S}[\![s]\!] \in \{\omega\}$ or $\mathcal{S}[\![o]\!] \in \{\omega\}$ would yield a contradiction so $\mathcal{S}[\![s]\!] \in \gamma^{\mathfrak{S}}(\mathcal{S}[\![s]\!])$ and $\mathcal{S}[\![o]\!] \in \gamma^{\mathfrak{O}}(\mathcal{O}[\![o]\!])$. It follows that

  $\mathcal{S}[\![\mathbb{1}[p]]\!]$
  $\triangleq \left(\mathcal{S}[\![p]\!] \notin \{\omega, \Omega\} \mathbin{?} \mathcal{S}[\![p]\!] \xleftarrow[\lambda P \bullet P]{\lambda Q \bullet Q} \mathcal{S}[\![p]\!] \mathbin{:} \mathcal{S}[\![p]\!]\right)$
  $= \mathcal{S}[\![p]\!] \xleftarrow[\lambda P \bullet P]{\lambda Q \bullet Q} \mathcal{S}[\![p]\!]$ ⟅case $\mathcal{S}[\![p]\!] \notin \{\omega, \Omega\}$⟆
  $= \mathcal{S}[\![\langle s,\ o\rangle]\!] \xleftarrow[\lambda P \bullet P]{\lambda Q \bullet Q} \mathcal{S}[\![\langle s,\ o\rangle]\!]$ ⟅$p = \langle s,\ o\rangle$⟆
  $\triangleq \left(\mathcal{S}[\![s]\!] \notin \{\omega, \Omega\} \wedge \mathcal{S}[\![o]\!] \notin \{\omega, \Omega\} \mathbin{?} \langle \mathcal{S}[\![s]\!],\ \mathcal{S}[\![o]\!]\rangle \mathbin{:} \texttt{error}(\mathcal{S}[\![s]\!], \mathcal{S}[\![o]\!])\right) \xleftarrow[\lambda P \bullet P]{\lambda Q \bullet Q} \left(\mathcal{S}[\![s]\!] \notin \{\omega, \Omega\} \wedge \mathcal{S}[\![o]\!] \notin \{\omega, \Omega\} \mathbin{?} \langle \mathcal{S}[\![s]\!],\ \mathcal{S}[\![o]\!]\rangle \mathbin{:} \texttt{error}(\mathcal{S}[\![s]\!], \mathcal{S}[\![o]\!])\right)$
  $= \langle \mathcal{S}[\![s]\!],\ \mathcal{S}[\![o]\!]\rangle \xleftarrow[\lambda P \bullet P]{\lambda Q \bullet Q} \langle \mathcal{S}[\![s]\!],\ \mathcal{S}[\![o]\!]\rangle$
  $\in \{\langle C, \preceq\rangle \xleftrightarrow[\alpha]{\gamma} \langle A, \sqsubseteq\rangle \mid C \in \gamma^{\mathfrak{S}}(\mathcal{S}[\![s]\!]) \wedge \preceq \in \gamma^{\mathfrak{O}}(\mathcal{O}[\![o]\!]) \wedge A \in \gamma^{\mathfrak{S}}(\mathcal{S}[\![s]\!]) \wedge \sqsubseteq \in \gamma^{\mathfrak{O}}(\mathcal{O}[\![o]\!])\}$
  $\triangleq \gamma^{\mathfrak{T}}(\mathcal{S}[\![s]\!] \circledast \mathcal{O}[\![o]\!] \rightleftharpoons \mathcal{S}[\![s]\!] \circledast \mathcal{O}[\![o]\!])$
  $= \gamma^{\mathfrak{T}}(\mathcal{P}[\![\langle s,\ o\rangle]\!] \rightleftharpoons \mathcal{P}[\![\langle s,\ o\rangle]\!])$
  $= \gamma^{\mathfrak{T}}(\mathcal{P}[\![p]\!] \rightleftharpoons \mathcal{P}[\![p]\!])$ ⟅$p = \langle s,\ o\rangle$⟆
  $\subseteq \gamma^{\mathfrak{T}}(\mathcal{P}[\![p]\!] \rightleftharpoons \mathcal{P}[\![p]\!]) \cup \{\omega\}$
  $= \gamma^{\mathfrak{T}}(\mathcal{T}[\![\mathbb{1}[p]]\!]) \cup \{\omega\}$ *QED.*

- $\mathcal{T}[\![\top[p, e]]\!] \triangleq \left(\mathcal{C}[\![e]\!] \neq \texttt{err} \wedge \exists \mathsf{S} \in \mathfrak{S}, \mathsf{O} \in \mathfrak{O} : \mathcal{P}[\![p]\!] = \mathsf{S} \circledast \mathsf{O} \wedge \mathcal{C}[\![e]\!] \Subset \mathsf{S} \mathbin{?} \mathcal{P}[\![p]\!] \rightleftharpoons \mathcal{P}[\![p]\!] \mathbin{:} \texttt{err}\right)$

  The soundness proof must be done in the case when $\mathcal{T}[\![\top[p, e]]\!] \neq \texttt{err}$ which implies $\mathcal{C}[\![e]\!] \neq \texttt{err} \wedge \exists \mathsf{S} \in \mathfrak{S}, \mathsf{O} \in \mathfrak{O} : \mathcal{P}[\![p]\!] = \mathsf{S} \circledast \mathsf{O} \wedge \mathcal{C}[\![e]\!] \Subset \mathsf{S}$.
  Since $p = \langle s,\ o\rangle$ and $\mathcal{P}[\![\langle s,\ o\rangle]\!] \triangleq \left(\mathcal{S}[\![s]\!] \neq \texttt{err} \wedge \mathcal{O}[\![o]\!] \neq \texttt{err} \mathbin{?} \mathcal{S}[\![s]\!] \circledast \mathcal{O}[\![o]\!] \mathbin{:} \texttt{err}\right)$ we have $\mathcal{C}[\![e]\!] \neq \texttt{err}$, $\mathsf{S} = \mathcal{S}[\![s]\!] \neq \texttt{err}$, $\mathsf{O} = \mathcal{O}[\![o]\!] \neq \texttt{err}$, $\mathcal{C}[\![e]\!] \Subset \mathsf{S} = \mathcal{S}[\![s]\!]$, and $\mathcal{T}[\![\top[p, e]]\!] = \mathcal{P}[\![p]\!] \rightleftharpoons \mathcal{P}[\![p]\!]$.

  By structural induction hypothesis, $\mathcal{S}[\![e]\!] \in \gamma^{\mathfrak{E}}(\mathcal{C}[\![e]\!]) \not\ni \Omega$, $\mathcal{S}[\![s]\!] \in (\gamma^{\mathfrak{S}}(\mathcal{S}[\![s]\!]) \cup \{\omega\}) \not\ni \Omega$, $\mathcal{S}[\![o]\!] \in (\gamma^{\mathfrak{O}}(\mathcal{O}[\![o]\!]) \cup \{\omega\}) \not\ni \Omega$, and $\mathcal{S}[\![p]\!] \in (\gamma^{\mathfrak{P}}(\mathcal{P}[\![p]\!]) \cup \{\omega\}) \not\ni \Omega$

  - If $\mathcal{S}[\![s]\!] = \omega$, $\mathcal{S}[\![o]\!] = \omega$, or $\mathcal{S}[\![p]\!] = \omega$ then $\mathcal{S}[\![\top[p, e]]\!] = \left(\exists S, \leqslant : \mathcal{S}[\![p]\!] = \langle S,\ \leqslant\rangle \mathbin{?} \left(\mathcal{S}[\![e]\!] \in S \setminus \{\omega\} \mathbin{?} \left(\forall x \in S : x \leqslant \mathcal{S}[\![e]\!] \mathbin{?} \mathcal{S}[\![p]\!] \xleftarrow[\lambda P \bullet \mathcal{S}[\![e]\!]]{\lambda Q \bullet \mathcal{S}[\![e]\!]} \mathcal{S}[\![p]\!] \mathbin{:} \omega\right) \mathbin{:} \omega\right) \mathbin{:} \texttt{error}(\mathcal{S}[\![p]\!], \mathcal{S}[\![e]\!])\right) = \omega \in \{\omega\}$
  - Otherwise $\mathcal{S}[\![p]\!] \notin \{\omega, \Omega\}$ so $\mathcal{S}[\![p]\!] \in \gamma^{\mathfrak{P}}(\mathcal{P}[\![p]\!])$. We have $\mathcal{S}[\![p]\!] = \mathcal{S}[\![\langle s,\ o\rangle]\!] \triangleq \left(\mathcal{S}[\![s]\!] \notin \{\omega, \Omega\} \wedge \mathcal{S}[\![o]\!] \notin \{\omega, \Omega\} \mathbin{?} \langle \mathcal{S}[\![s]\!],\ \mathcal{S}[\![o]\!]\rangle \mathbin{:} \texttt{error}(\mathcal{S}[\![s]\!], \mathcal{S}[\![o]\!])\right)$ so $\mathcal{S}[\![s]\!] \in \{\omega, \Omega\}$ or $\mathcal{S}[\![o]\!] \in \{\omega, \Omega\}$ would imply the contradiction $\mathcal{S}[\![p]\!] \in \{\omega, \Omega\}$ so $\mathcal{S}[\![s]\!] \in \gamma^{\mathfrak{S}}(\mathcal{S}[\![s]\!])$, $\mathcal{S}[\![o]\!] \in \gamma^{\mathfrak{O}}(\mathcal{O}[\![o]\!])$, and $\mathcal{S}[\![p]\!] = \langle \mathcal{S}[\![s]\!],\ \mathcal{S}[\![o]\!]\rangle \in \gamma^{\mathfrak{P}}(\mathcal{P}[\![p]\!])$.

    It follows that $\mathcal{S}[\![e]\!] \in S \setminus \{\omega\} = \mathcal{S}[\![s]\!]$ and therefore

    $\mathcal{S}[\![\top[p, e]]\!]$
    $\triangleq \left(\exists S, \leqslant : \mathcal{S}[\![p]\!] = \langle S,\ \leqslant\rangle \mathbin{?} \left(\mathcal{S}[\![e]\!] \in S \setminus \{\omega\} \mathbin{?} \left(\forall x \in S : x \leqslant \mathcal{S}[\![e]\!] \mathbin{?} \mathcal{S}[\![p]\!] \xleftarrow[\lambda P \bullet \mathcal{S}[\![e]\!]]{\lambda Q \bullet \mathcal{S}[\![e]\!]} \mathcal{S}[\![p]\!] \mathbin{:} \omega\right) \mathbin{:} \omega\right) \mathbin{:} \texttt{error}(\mathcal{S}[\![p]\!], \mathcal{S}[\![e]\!])\right)$
    $= \left(\forall x \in \mathcal{S}[\![s]\!] : x \leqslant \mathcal{S}[\![e]\!] \mathbin{?} \mathcal{S}[\![p]\!] \xleftarrow[\lambda P \bullet \mathcal{S}[\![e]\!]]{\lambda Q \bullet \mathcal{S}[\![e]\!]} \mathcal{S}[\![p]\!] \mathbin{:} \omega\right)$ ⟅since $\mathcal{S}[\![p]\!] = \langle \mathcal{S}[\![s]\!],\ \mathcal{S}[\![o]\!]\rangle$ and $\mathcal{S}[\![e]\!] \in \mathcal{S}[\![s]\!]$⟆
    $\in \{\mathcal{S}[\![p]\!] \xleftarrow[\lambda P \bullet \mathcal{S}[\![e]\!]]{\lambda Q \bullet \mathcal{S}[\![e]\!]} \mathcal{S}[\![p]\!]\} \cup \{\omega\}$
    $\subseteq \{P \xleftrightarrow[\alpha]{\gamma} P' \mid P \in \gamma^{\mathfrak{P}}(\mathcal{P}[\![p]\!]) \wedge P' \in \gamma^{\mathfrak{P}}(\mathcal{P}[\![p]\!])\} \cup \{\omega\}$ ⟅since $\mathcal{S}[\![p]\!] \in \gamma^{\mathfrak{P}}(\mathcal{P}[\![p]\!])$⟆
    $= \gamma^{\mathfrak{T}}(\mathcal{P}[\![p]\!] \rightleftharpoons \mathcal{P}[\![p]\!]) \cup \{\omega\}$
    ⟅by definition $\gamma^{\mathfrak{T}}(\mathsf{P} \rightleftharpoons \mathsf{P}') \triangleq \{P \xleftrightarrow[\alpha]{\gamma} P' \mid P \in \gamma^{\mathfrak{P}}(\mathsf{P}) \wedge P' \in \gamma^{\mathfrak{P}}(\mathsf{P}')\}$⟆
    $= \gamma^{\mathfrak{T}}(\mathcal{T}[\![\top[p, e]]\!]) \cup \{\omega\}$ *QED.*

- $\mathcal{T}[\![\mathbb{I}[\langle s,\ o\rangle, b, t]]\!] \triangleq \left(\texttt{err} \neq \mathcal{C}[\![b]\!] \Subset \mathcal{S}[\![s]\!] \neq \texttt{err} \wedge \texttt{err} \neq \mathcal{C}[\![t]\!] \Subset \mathcal{S}[\![s]\!] \mathbin{?} (\mathbf{P}\ \mathcal{S}[\![s]\!] \circledast \subseteq) \rightleftharpoons (\mathbf{P}\ \mathcal{S}[\![s]\!] \circledast \subseteq) \mathbin{:} \texttt{err}\right)$

  For example, $\mathbb{I}[\langle \mathbb{Z},\ \leq\rangle, -\infty, \infty]$ denotes the *GC* $\langle \wp(\mathbb{Z}),\ \subseteq\rangle \xleftrightarrow[\alpha^{\mathbb{I}}]{\gamma^{\mathbb{I}}} \langle \Im(\mathbb{Z} \cup \{-\infty, \infty\}, \leqslant),\ \subseteq\rangle$ where $\wp(\mathbb{Z}) \in \wp(\wp(\mathbb{Z} \cup \{-\infty, \infty\}))$ so $\wp(\mathbb{Z})$ has type $\mathbf{P}\ \mathbb{Z}$ and $\Im(\mathbb{Z} \cup \{-\infty, \infty\}, \leqslant) \subseteq \wp(\mathbb{Z} \cup \{-\infty, \infty\})$ so $\Im(\mathbb{Z} \cup \{-\infty, \infty\}, \leqslant) \in \wp(\wp(\mathbb{Z} \cup \{-\infty, \infty\}))$ hence $\Im(\mathbb{Z} \cup \{-\infty, \infty\}, \leqslant)$ has type $\mathbf{P}\ \mathbb{Z}$. It follows that $\mathcal{T}[\![\mathbb{I}[\langle \mathbb{Z},\ \leq\rangle, -\infty, \infty]]\!] = (\mathbf{P}\ \mathbb{Z} \circledast \subseteq) \rightleftharpoons (\mathbf{P}\ \mathbb{Z} \circledast \subseteq)$.

  For another example, $\mathbb{I}[\langle \wp(\mathbb{Z}),\ \subseteq\rangle, \emptyset, \mathbb{Z}]$ denotes the *GC* $\langle \wp(\wp(\mathbb{Z})),\ \subseteq\rangle \xleftrightarrow[\alpha^{\mathbb{I}}]{\gamma^{\mathbb{I}}} \langle \Im(\wp(\mathbb{Z}) \cup \{\emptyset, \mathbb{Z}\}, \subseteq),\ \subseteq\rangle = \langle \Im(\wp(\mathbb{Z}), \subseteq),\ \subseteq\rangle$ where $\wp(\wp(\mathbb{Z})) \in \wp(\wp(\wp(\mathbb{Z})))$ and $\Im(\wp(\mathbb{Z}), \subseteq) \subseteq \wp(\wp(\mathbb{Z}))$ such that $\Im(\wp(\mathbb{Z}), \subseteq) \in \wp(\wp(\wp(\mathbb{Z})))$ have type $\mathbf{P}\ \mathbf{P}\ \mathbb{Z}$. It follows that $\mathcal{T}[\![\mathbb{I}[\langle \wp(\mathbb{Z}),\ \subseteq\rangle, \emptyset, \mathbb{Z}]]\!] = (\mathbf{P}\ \mathbf{P}\ \mathbb{Z} \circledast \subseteq) \rightleftharpoons (\mathbf{P}\ \mathbf{P}\ \mathbb{Z} \circledast \subseteq)$

  The soundness proof must be done in the case when $\mathcal{T}[\![\mathbb{I}[\langle s,\ o\rangle, b, t]]\!] \neq \texttt{err}$ which implies $\texttt{err} \neq \mathcal{C}[\![b]\!] \Subset \mathcal{S}[\![s]\!] \neq \texttt{err}$ so $\mathcal{S}[\![b]\!] \in \gamma^{\mathfrak{S}}(\mathcal{S}[\![s]\!]) \cup \{\omega\}$, $\texttt{err} \neq \mathcal{C}[\![t]\!] \Subset \mathcal{S}[\![s]\!]$ so $\mathcal{S}[\![t]\!] \in \gamma^{\mathfrak{S}}(\mathcal{S}[\![s]\!]) \cup \{\omega\}$, and $\mathcal{T}[\![\mathbb{I}[\langle s,\ o\rangle, b, t]]\!] \triangleq (\mathbf{P}\ \mathcal{S}[\![s]\!] \circledast \subseteq) \rightleftharpoons (\mathbf{P}\ \mathcal{S}[\![s]\!] \circledast \subseteq)$.

  - If $\mathcal{S}[\![\mathbb{I}[\langle s,\ o\rangle, b, t]]\!] = \omega$ then obviously $\mathcal{S}[\![\mathbb{I}[\langle s,\ o\rangle, b, t]]\!] \in \gamma^{\mathfrak{T}}(\mathcal{T}[\![\mathbb{I}[\langle s,\ o\rangle, b, t]]\!]) \cup \{\omega\}$.
  - Otherwise $\mathcal{S}[\![\mathbb{I}[\langle s,\ o\rangle, b, t]]\!] \triangleq \left(\mathcal{S}[\![s]\!] \notin \{\omega, \Omega\} \wedge \mathcal{S}[\![b]\!] \notin \{\omega, \Omega\} \wedge \mathcal{S}[\![t]\!] \notin \{\omega, \Omega\} \mathbin{?} \left(\mathcal{S}[\![o]\!] \subseteq (\mathcal{S}[\![s]\!] \cup \{\mathcal{S}[\![b]\!], \mathcal{S}[\![t]\!]\} \times (\mathcal{S}[\![s]\!] \cup \{\mathcal{S}[\![b]\!], \mathcal{S}[\![t]\!]\}) \wedge \forall x \in \mathcal{S}[\![s]\!] : \langle \mathcal{S}[\![b]\!],\ x\rangle \in \mathcal{S}[\![o]\!] \wedge \langle x,\ \mathcal{S}[\![t]\!]\rangle \in \mathcal{S}[\![o]\!] \mathbin{?} \langle \wp(\mathcal{S}[\![s]\!]),\ \subseteq\rangle \xleftrightarrow[\alpha^{\mathbb{I}}]{\gamma^{\mathbb{I}}} \langle \Im(\mathcal{S}[\![s]\!] \cup \{\mathcal{S}[\![b]\!], \mathcal{S}[\![t]\!]\}, \mathcal{S}[\![o]\!]), \subseteq\rangle \mathbin{:} \omega\right) \mathbin{:} \texttt{error}(\mathcal{S}[\![s]\!], \texttt{error}(\mathcal{S}[\![b]\!], \mathcal{S}[\![t]\!]))\right) \neq \omega$ and so necessarily $\mathcal{S}[\![s]\!] \notin \{\omega, \Omega\}$, $\mathcal{S}[\![b]\!] \notin \{\omega, \Omega\}$, $\mathcal{S}[\![t]\!] \notin \{\omega, \Omega\}$, $\mathcal{S}[\![o]\!] \subseteq (\mathcal{S}[\![s]\!] \cup \{\mathcal{S}[\![b]\!], \mathcal{S}[\![t]\!]\}) \times (\mathcal{S}[\![s]\!] \cup \{\mathcal{S}[\![b]\!], \mathcal{S}[\![t]\!]\})$, $\forall x \in \mathcal{S}[\![s]\!] : \langle \mathcal{S}[\![b]\!],\ x\rangle \in \mathcal{S}[\![o]\!] \wedge \langle x,\ \mathcal{S}[\![t]\!]\rangle \in \mathcal{S}[\![o]\!]$, and $\mathcal{S}[\![\mathbb{I}[\langle s,\ o\rangle, b, t]]\!] \triangleq \langle \wp(\mathcal{S}[\![s]\!]),\ \subseteq\rangle \xleftrightarrow[\alpha^{\mathbb{I}}]{\gamma^{\mathbb{I}}} \langle \Im(\mathcal{S}[\![s]\!] \cup \{\mathcal{S}[\![b]\!], \mathcal{S}[\![t]\!]\}, \mathcal{S}[\![o]\!]),\ \subseteq\rangle$. By structural induction, it follows that $\mathcal{S}[\![s]\!] \in \gamma^{\mathfrak{S}}(\mathcal{S}[\![s]\!])$ so

    $\langle \wp(\mathcal{S}[\![s]\!]),\ \subseteq\rangle$
    $\in \gamma^{\mathfrak{S}}(\mathbf{P}\ \mathcal{S}[\![s]\!]) \times \{\subseteq\}$
    ⟅since $\mathcal{S}[\![s]\!] \in \gamma^{\mathfrak{S}}(\mathcal{S}[\![s]\!])$ so $\wp(\mathcal{S}[\![s]\!]) \in \{\wp(X) \mid X \in \gamma^{\mathfrak{S}}(\mathcal{S}[\![s]\!])\} \subseteq \gamma^{\mathfrak{S}}(\mathbf{P}\ \mathcal{S}[\![s]\!])$ and $\subseteq\ \in \{\subseteq\}$⟆
    $= \gamma^{\mathfrak{S}}(\mathbf{P}\ \mathcal{S}[\![s]\!]) \times \gamma^{\mathfrak{O}}(\subseteq)$ ⟅since $\gamma^{\mathfrak{O}}(\subseteq) \triangleq \{\mathcal{S}[\![\subseteq]\!]\} = \{\subseteq\}$⟆
    $= \gamma^{\mathfrak{P}}(\mathbf{P}\ \mathcal{S}[\![s]\!] \circledast \subseteq)$ ⟅since $\gamma^{\mathfrak{P}}(\mathsf{S} \circledast \mathsf{O}) \triangleq \gamma^{\mathfrak{S}}(\mathsf{S}) \times \gamma^{\mathfrak{O}}(\mathsf{O})$⟆

Moreover $\mathscr{E}[\![b]\!], \mathscr{E}[\![t]\!] \Subset \mathcal{S}[\![s]\!]$ implies $\mathcal{S}[\![b]\!], \mathcal{S}[\![t]\!] \in \gamma^{\mathfrak{S}}(\mathcal{S}[\![s]\!])$ and therefore $\mathcal{S}[\![s]\!] \cup \{\mathcal{S}[\![b]\!], \mathcal{S}[\![t]\!]\} \in \gamma^{\mathfrak{S}}(\mathcal{S}[\![s]\!])$ so

$$
\begin{array}{ll}
& \langle \Im(\mathcal{S}[\![s]\!] \cup \{\mathcal{S}[\![b]\!], \mathcal{S}[\![t]\!]\}, \mathcal{S}[\![o]\!]),\ \subseteq \rangle \\
\in & \{\Im(X, \mathcal{S}[\![o]\!]) \mid X \in \gamma^{\mathfrak{S}}(\mathcal{S}[\![s]\!])\} \times \{\subseteq\} \\
& \qquad \text{since } \mathcal{S}[\![s]\!] \cup \{\mathcal{S}[\![b]\!], \mathcal{S}[\![t]\!]\} \in \gamma^{\mathfrak{S}}(\mathcal{S}[\![s]\!]) \\
\subseteq & \{\wp(X) \mid X \in \gamma^{\mathfrak{S}}(\mathcal{S}[\![s]\!])\} \times \{\subseteq\} \\
& \qquad \text{since } \Im(S, \leqslant) \subseteq \wp(S) \\
\subseteq & \wp(\gamma^{\mathfrak{S}}(\mathcal{S}[\![s]\!])) \times \{\subseteq\} \\
& \qquad \text{since } \{\wp(X) \mid X \in \gamma^{\mathfrak{S}}(\mathsf{S})\} \subseteq \gamma^{\mathfrak{S}}(\mathtt{P}\ \mathsf{S}) \\
= & \gamma^{\mathfrak{S}}(\mathtt{P}\ \mathcal{S}[\![s]\!]) \times \{\subseteq\} \qquad \text{since } \gamma^{\mathfrak{S}}(\mathtt{P}\ \mathsf{S}) \triangleq \wp(\gamma^{\mathfrak{S}}(\mathsf{S})) \\
= & \gamma^{\mathfrak{S}}(\mathtt{P}\ \mathcal{S}[\![s]\!]) \times \gamma^{\mathfrak{O}}(\subseteq) \quad \text{since } \gamma^{\mathfrak{O}}(\subseteq) \triangleq \{\mathcal{S}[\![\subseteq]\!]\} = \{\subseteq\} \\
= & \gamma^{\mathfrak{P}}(\mathtt{P}\ \mathcal{S}[\![s]\!] \circledast \subseteq) \quad \text{since } \gamma^{\mathfrak{P}}(\mathsf{S} \circledast \mathsf{O}) \triangleq \gamma^{\mathfrak{S}}(\mathsf{S}) \times \gamma^{\mathfrak{O}}(\mathsf{O})
\end{array}
$$

We conclude that

$$
\begin{array}{ll}
& \mathcal{S}[\![\mathbb{I}[\langle s,\ o\rangle, b, t]]\!] \\
\triangleq & \langle \wp(\mathcal{S}[\![s]\!]),\ \subseteq \rangle \underset{\alpha^{\mathbb{I}}}{\overset{\gamma^{\mathbb{I}}}{\leftrightarrows}} \langle \Im(\mathcal{S}[\![s]\!] \cup \{\mathcal{S}[\![b]\!], \mathcal{S}[\![t]\!]\}, \mathcal{S}[\![o]\!]),\ \subseteq \rangle \\
\in & \{P \underset{\alpha}{\overset{\gamma}{\leftrightarrows}} P' \mid P \in \gamma^{\mathfrak{P}}(\mathtt{P}\ \mathcal{S}[\![s]\!] \circledast \subseteq) \wedge P' \in \gamma^{\mathfrak{P}}(\mathtt{P}\ \mathcal{S}[\![s]\!] \circledast \subseteq)\} \\
& \qquad \text{since } \langle \wp(\mathcal{S}[\![s]\!]), \subseteq\rangle \in \gamma^{\mathfrak{P}}(\mathtt{P}\ \mathcal{S}[\![s]\!] \circledast \subseteq) \text{ and } \langle \Im(\mathcal{S}[\![s]\!] \cup \{\mathcal{S}[\![b]\!], \mathcal{S}[\![t]\!]\}, \mathcal{S}[\![o]\!]),\ \subseteq\rangle \in \gamma^{\mathfrak{P}}(\mathtt{P}\ \mathcal{S}[\![s]\!] \circledast \subseteq) \\
\subseteq & \gamma^{\mathfrak{T}}((\mathtt{P}\ \mathcal{S}[\![s]\!] \circledast \subseteq) \leftrightharpoons (\mathtt{P}\ \mathcal{S}[\![s]\!] \circledast \subseteq)) \cup \{\omega\} \\
& \qquad \text{by definition } \gamma^{\mathfrak{T}}(\mathsf{P} \leftrightharpoons \mathsf{P}') \triangleq \{P \underset{\alpha}{\overset{\gamma}{\leftrightarrows}} P' \mid P \in \gamma^{\mathfrak{P}}(\mathsf{P}) \wedge P' \in \gamma^{\mathfrak{P}}(\mathsf{P}')\} \\
= & \gamma^{\mathfrak{T}}(\mathcal{T}[\![\mathbb{I}[\langle s,\ o\rangle, b, t]]\!]) \cup \{\omega\} \qquad \textit{QED.}
\end{array}
$$

- $\mathcal{T}[\![\frown[s_{\mathbb{L}}, s_{\mathcal{M}}]]\!] \triangleq \big(\mathcal{S}[\![s_{\mathbb{L}}]\!] \neq \mathtt{err} \wedge \mathcal{S}[\![s_{\mathcal{M}}]\!] \neq \mathtt{err}\ ?\ \mathtt{P}\ (\mathcal{S}[\![s_{\mathbb{L}}]\!] * \mathcal{S}[\![s_{\mathcal{M}}]\!]) \circledast \subseteq \leftrightharpoons \mathcal{S}[\![s_{\mathbb{L}}]\!] \twoheadrightarrow \mathtt{P}\ \mathcal{S}[\![s_{\mathcal{M}}]\!] \circledast \dot{\subseteq}\ :\ \mathtt{err}\big)$

Assuming $\mathcal{T}[\![\frown[s_{\mathbb{L}}, s_{\mathcal{M}}]]\!] \neq \mathtt{err}$ hence $\mathcal{S}[\![s_{\mathbb{L}}]\!] \neq \mathtt{err}$ and $\mathcal{S}[\![s_{\mathcal{M}}]\!] \neq \mathtt{err}$ so, by structural induction hypothesis, that $\mathcal{S}[\![s_{\mathbb{L}}]\!] \in \gamma^{\mathfrak{S}}(\mathcal{S}[\![s_{\mathbb{L}}]\!]) \cup \{\omega\}$ and $\mathcal{S}[\![s_{\mathcal{M}}]\!] \in \gamma^{\mathfrak{S}}(\mathcal{S}[\![s_{\mathcal{M}}]\!]) \cup \{\omega\}$, we must prove that $\mathcal{S}[\![\frown[s_{\mathbb{L}}, s_{\mathcal{M}}]]\!] \in \gamma^{\mathfrak{T}}(\mathcal{T}[\![\frown[s_{\mathbb{L}}, s_{\mathcal{M}}]]\!]) \cup \{\omega\}$. Observe that in case $\mathcal{S}[\![s_{\mathbb{L}}]\!] \neq \omega$ and $\mathcal{S}[\![s_{\mathcal{M}}]\!] \neq \omega$, we have

$$
\begin{array}{ll}
& \wp(\mathcal{S}[\![s_{\mathbb{L}}]\!] \times \mathcal{S}[\![s_{\mathcal{M}}]\!]) \\
\in & \{\wp(X \times Y) \mid X \in \gamma^{\mathfrak{S}}(\mathcal{S}[\![s_{\mathbb{L}}]\!]) \wedge Y \in \gamma^{\mathfrak{S}}(\mathcal{S}[\![s_{\mathcal{M}}]\!])\} \\
& \qquad \text{since } \mathcal{S}[\![s_{\mathbb{L}}]\!] \in \gamma^{\mathfrak{S}}(\mathcal{S}[\![s_{\mathbb{L}}]\!]) \text{ and } \mathcal{S}[\![s_{\mathcal{M}}]\!] \in \gamma^{\mathfrak{S}}(\mathcal{S}[\![s_{\mathcal{M}}]\!]) \\
= & \{\wp(X \times Y) \mid X \times Y \in \gamma^{\mathfrak{S}}(\mathcal{S}[\![s_{\mathbb{L}}]\!] * \mathcal{S}[\![s_{\mathcal{M}}]\!])\} \\
& \qquad \text{since } \gamma^{\mathfrak{S}}(\mathsf{S}_1 * \mathsf{S}_2) \triangleq \{X \times Y \mid X \in \gamma^{\mathfrak{S}}(\mathsf{S}_1) \wedge Y \in \gamma^{\mathfrak{S}}(\mathsf{S}_2)\} \\
\subseteq & \gamma^{\mathfrak{S}}(\mathtt{P}\ (\mathcal{S}[\![s_{\mathbb{L}}]\!] * \mathcal{S}[\![s_{\mathcal{M}}]\!])) \\
& \qquad \text{since } \{\wp(Z) \mid Z \in \gamma^{\mathfrak{S}}(\mathsf{S})\} \subseteq \gamma^{\mathfrak{S}}(\mathtt{P}\ \mathsf{S}) \quad .
\end{array}
$$

and $\subseteq\ \in \{\subseteq\} = \{\mathcal{S}[\![\subseteq]\!]\} \triangleq \gamma^{\mathfrak{O}}(\subseteq)$ so $\langle \wp(\mathcal{S}[\![s_{\mathbb{L}}]\!] \times \mathcal{S}[\![s_{\mathcal{M}}]\!]), \subseteq\rangle \in \gamma^{\mathfrak{S}}(\mathtt{P}\ (\mathcal{S}[\![s_{\mathbb{L}}]\!] * \mathcal{S}[\![s_{\mathcal{M}}]\!])) \times \gamma^{\mathfrak{O}}(\subseteq) \triangleq \gamma^{\mathfrak{P}}(\mathtt{P}\ (\mathcal{S}[\![s_{\mathbb{L}}]\!] * \mathcal{S}[\![s_{\mathcal{M}}]\!]) \circledast \subseteq)$. Moreover

$$
\begin{array}{ll}
& \mathcal{S}[\![s_{\mathbb{L}}]\!] \mapsto \wp(\mathcal{S}[\![s_{\mathcal{M}}]\!]) \\
\in & \{X \mapsto \wp(Y) \mid X \in \gamma^{\mathfrak{S}}(\mathcal{S}[\![s_{\mathbb{L}}]\!]) \wedge Y \in \gamma^{\mathfrak{S}}(\mathcal{S}[\![s_{\mathcal{M}}]\!])\} \\
& \qquad \text{since } \mathcal{S}[\![s_{\mathbb{L}}]\!] \in \gamma^{\mathfrak{S}}(\mathcal{S}[\![s_{\mathbb{L}}]\!]) \text{ and } \mathcal{S}[\![s_{\mathcal{M}}]\!] \in \gamma^{\mathfrak{S}}(\mathcal{S}[\![s_{\mathcal{M}}]\!]) \\
= & \{X \mapsto Z \mid X \in \gamma^{\mathfrak{S}}(\mathcal{S}[\![s_{\mathbb{L}}]\!]) \wedge Z \in \gamma^{\mathfrak{S}}(\mathtt{P}\ \mathcal{S}[\![s_{\mathcal{M}}]\!])\} \\
& \qquad \text{since } \gamma^{\mathfrak{S}}(\mathtt{P}\ \mathsf{S}) \triangleq \wp(\gamma^{\mathfrak{S}}(\mathsf{S})) \\
= & \gamma^{\mathfrak{S}}(\mathcal{S}[\![s_{\mathbb{L}}]\!] \twoheadrightarrow \mathtt{P}\ \mathcal{S}[\![s_{\mathcal{M}}]\!]) \\
& \qquad \text{since } \gamma^{\mathfrak{S}}(\mathsf{S}_1 \twoheadrightarrow \mathsf{S}_2) \triangleq \{X \mapsto Y \mid X \in \gamma^{\mathfrak{S}}(\mathsf{S}_1) \wedge Y \in \gamma^{\mathfrak{S}}(\mathsf{S}_2)\}
\end{array}
$$

and $\dot{\subseteq} \in \gamma^{\mathfrak{O}}(\dot{\subseteq})$ so $\langle \mathcal{S}[\![s_{\mathbb{L}}]\!] \mapsto \wp(\mathcal{S}[\![s_{\mathcal{M}}]\!]),\ \dot{\subseteq}\rangle \in \gamma^{\mathfrak{S}}(\mathcal{S}[\![s_{\mathbb{L}}]\!] \twoheadrightarrow \mathtt{P}\ \mathcal{S}[\![s_{\mathcal{M}}]\!]) \times \gamma^{\mathfrak{O}}(\dot{\subseteq}) \triangleq \gamma^{\mathfrak{P}}(\mathcal{S}[\![s_{\mathbb{L}}]\!] \twoheadrightarrow \mathtt{P}\ \mathcal{S}[\![s_{\mathcal{M}}]\!] \circledast \dot{\subseteq})$.
It follows that for the non-trivial case $\mathcal{S}[\![\frown[s_{\mathbb{L}}, s_{\mathcal{M}}]]\!] \neq \omega$, we have

$$
\begin{array}{ll}
& \mathcal{S}[\![\frown[s_{\mathbb{L}}, s_{\mathcal{M}}]]\!] \\
= & \big(\mathcal{S}[\![s_{\mathbb{L}}]\!] \notin \{\Omega, \omega\} \wedge \mathcal{S}[\![s_{\mathcal{M}}]\!] \notin \{\Omega, \omega\}\ ?\ \langle \wp(\mathcal{S}[\![s_{\mathbb{L}}]\!] \times \mathcal{S}[\![s_{\mathcal{M}}]\!]),\ \subseteq\rangle \underset{\alpha^{\frown}}{\overset{\gamma^{\frown}}{\leftrightarrows}} \langle \mathcal{S}[\![s_{\mathbb{L}}]\!] \mapsto \wp(\mathcal{S}[\![s_{\mathcal{M}}]\!]),\ \dot{\subseteq}\rangle\ :\ \mathtt{error}(\mathcal{S}[\![s_{\mathbb{L}}]\!], \mathcal{S}[\![s_{\mathcal{M}}]\!])\big) \\
= & \langle \wp(\mathcal{S}[\![s_{\mathbb{L}}]\!] \times \mathcal{S}[\![s_{\mathcal{M}}]\!]),\ \subseteq\rangle \underset{\alpha^{\frown}}{\overset{\gamma^{\frown}}{\leftrightarrows}} \langle \mathcal{S}[\![s_{\mathbb{L}}]\!] \mapsto \wp(\mathcal{S}[\![s_{\mathcal{M}}]\!]),\ \dot{\subseteq}\rangle \\
\in & \{P \underset{\alpha}{\overset{\gamma}{\leftrightarrows}} P' \mid P \in \gamma^{\mathfrak{P}}(\mathtt{P}\ (\mathcal{S}[\![s_{\mathbb{L}}]\!] * \mathcal{S}[\![s_{\mathcal{M}}]\!]) \circledast \subseteq) \wedge P' \in \gamma^{\mathfrak{P}}(\mathcal{S}[\![s_{\mathbb{L}}]\!] \twoheadrightarrow \mathtt{P}\ \mathcal{S}[\![s_{\mathcal{M}}]\!] \circledast \dot{\subseteq})\} \\
& \qquad \text{since } \langle \wp(\mathcal{S}[\![s_{\mathbb{L}}]\!] \times \mathcal{S}[\![s_{\mathcal{M}}]\!]),\ \subseteq\rangle \in \gamma^{\mathfrak{P}}(\mathtt{P}\ (\mathcal{S}[\![s_{\mathbb{L}}]\!] * \mathcal{S}[\![s_{\mathcal{M}}]\!]) \circledast \subseteq) \text{ and } \langle \mathcal{S}[\![s_{\mathbb{L}}]\!] \mapsto \wp(\mathcal{S}[\![s_{\mathcal{M}}]\!]),\ \dot{\subseteq}\rangle \in \gamma^{\mathfrak{P}}(\mathcal{S}[\![s_{\mathbb{L}}]\!] \twoheadrightarrow \mathtt{P}\ \mathcal{S}[\![s_{\mathcal{M}}]\!] \circledast \dot{\subseteq}) \\
= & \gamma^{\mathfrak{T}}(\mathtt{P}\ (\mathcal{S}[\![s_{\mathbb{L}}]\!] * \mathcal{S}[\![s_{\mathcal{M}}]\!]) \circledast \subseteq \leftrightharpoons \mathcal{S}[\![s_{\mathbb{L}}]\!] \twoheadrightarrow \mathtt{P}\ \mathcal{S}[\![s_{\mathcal{M}}]\!] \circledast \dot{\subseteq}) \\
& \qquad \text{since } \gamma^{\mathfrak{T}}(\mathsf{P} \leftrightharpoons \mathsf{P}') \triangleq \{P \underset{\alpha}{\overset{\gamma}{\leftrightarrows}} P' \mid P \in \gamma^{\mathfrak{P}}(\mathsf{P}) \wedge P' \in \gamma^{\mathfrak{P}}(\mathsf{P}')\} \\
= & \gamma^{\mathfrak{T}}(\mathcal{T}[\![\frown[s_{\mathbb{L}}, s_{\mathcal{M}}]]\!]) \qquad \textit{QED.}
\end{array}
$$

- $\mathcal{T}[\![\cup[s]]\!] \triangleq \big(\mathcal{S}[\![s]\!] \neq \mathtt{err}\ ?\ \mathtt{P}\ (\mathtt{P}\ \mathcal{S}[\![s]\!]) \circledast \subseteq \leftrightharpoons \mathtt{P}\ \mathcal{S}[\![s]\!] \circledast \subseteq\ :\ \mathtt{err}\big)$

Assuming $\mathcal{T}[\![\cup[s]]\!] \neq \mathtt{err}$ hence $\mathcal{S}[\![s]\!] \neq \mathtt{err}$ so, by structural induction hypothesis, that $\mathcal{S}[\![s]\!] \in \gamma^{\mathfrak{S}}(\mathcal{S}[\![s]\!]) \cup \{\omega\}$, we must prove that $\mathcal{S}[\![\cup[s]]\!] \in \gamma^{\mathfrak{T}}(\mathcal{T}[\![\cup[s]]\!]) \cup \{\omega\}$. For the non-trivial case $\mathcal{S}[\![\cup[s]]\!] \neq \omega$, we have

$$
\begin{array}{ll}
& \mathcal{S}[\![\cup[s]]\!] \\
= & \big(\mathcal{S}[\![s]\!] \notin \{\omega, \Omega\}\ ?\ \langle \wp(\wp(\mathcal{S}[\![s]\!])),\ \subseteq\rangle \underset{\alpha^{\wp}}{\overset{\gamma^{\wp}}{\leftrightarrows}} \langle \wp(\mathcal{S}[\![s]\!]),\ \subseteq\rangle\ :\ \mathcal{S}[\![s]\!]\big) \\
= & \langle \wp(\wp(\mathcal{S}[\![s]\!])),\ \subseteq\rangle \underset{\alpha^{\wp}}{\overset{\gamma^{\wp}}{\leftrightarrows}} \langle \wp(\mathcal{S}[\![s]\!]),\ \subseteq\rangle \\
\in & \{P \underset{\alpha}{\overset{\gamma}{\leftrightarrows}} P' \mid P \in \gamma^{\mathfrak{P}}(\mathtt{P}\ (\mathtt{P}\ \mathcal{S}[\![s]\!]) \circledast \subseteq) \wedge P' \in \gamma^{\mathfrak{P}}(\mathtt{P}\ \mathcal{S}[\![s]\!] \circledast \subseteq)\} \\
& \qquad \text{since } \langle \wp(\wp(\mathcal{S}[\![s]\!])),\ \subseteq\rangle \in \gamma^{\mathfrak{P}}(\mathtt{P}\ (\mathtt{P}\ \mathcal{S}[\![s]\!]) \circledast \subseteq) \text{ and } \langle \wp(\mathcal{S}[\![s]\!]),\ \subseteq\rangle \in \gamma^{\mathfrak{P}}(\mathtt{P}\ \mathcal{S}[\![s]\!] \circledast \subseteq) \\
= & \gamma^{\mathfrak{T}}((\mathtt{P}\ (\mathtt{P}\ \mathcal{S}[\![s]\!]) \circledast \subseteq) \leftrightharpoons (\mathtt{P}\ \mathcal{S}[\![s]\!] \circledast \subseteq)) \\
& \qquad \text{since } \gamma^{\mathfrak{T}}(\mathsf{P} \leftrightharpoons \mathsf{P}') \triangleq \{P \underset{\alpha}{\overset{\gamma}{\leftrightarrows}} P' \mid P \in \gamma^{\mathfrak{P}}(\mathsf{P}) \wedge P' \in \gamma^{\mathfrak{P}}(\mathsf{P}')\} \\
= & \gamma^{\mathfrak{T}}(\mathcal{T}[\![\cup[s]]\!]) \qquad \textit{QED.}
\end{array}
$$

- $\mathcal{T}[\![\neg[s]]\!] \triangleq \big(\mathcal{S}[\![s]\!] \neq \mathtt{err}\ ?\ \mathtt{P}\ \mathcal{S}[\![s]\!] \circledast \subseteq \leftrightharpoons \mathtt{P}\ \mathcal{S}[\![s]\!] \circledast \subseteq^{-1}\ :\ \mathtt{err}\big)$

In $\mathcal{T}[\![\neg[s]]\!] \neq \mathtt{err}$ so $\mathcal{S}[\![s]\!] \neq \mathtt{err}$, we have $\mathcal{S}[\![s]\!] \in \gamma^{\mathfrak{S}}(\mathcal{S}[\![s]\!]) \cup \{\omega\}$ by structural induction hypothesis. For the non-trivial case $\mathcal{S}[\![\neg[s]]\!] \neq \omega$ so $\mathcal{S}[\![s]\!] \neq \omega$ hence $\wp(\mathcal{S}[\![s]\!]) \in \gamma^{\mathfrak{S}}(\mathtt{P}\ \mathcal{S}[\![s]\!])$, we have

$$
\begin{array}{ll}
& \mathcal{S}[\![\neg[s]]\!] \\
= & \big(\mathcal{S}[\![s]\!] \notin \{\omega, \Omega\}\ ?\ \langle \wp(\mathcal{S}[\![s]\!]),\ \subseteq\rangle \underset{\neg}{\overset{\neg}{\leftrightarrows}} \langle \wp(\mathcal{S}[\![s]\!]),\ \supseteq\rangle\ :\ \mathcal{S}[\![s]\!]\big) \\
= & \langle \wp(\mathcal{S}[\![s]\!]),\ \subseteq\rangle \underset{\neg}{\overset{\neg}{\leftrightarrows}} \langle \wp(\mathcal{S}[\![s]\!]),\ \supseteq\rangle \\
\in & \{P \underset{\alpha}{\overset{\gamma}{\leftrightarrows}} P' \mid P \in \gamma^{\mathfrak{P}}(\mathtt{P}\ \mathcal{S}[\![s]\!] \circledast \subseteq) \wedge P' \in \gamma^{\mathfrak{P}}(\mathtt{P}\ \mathcal{S}[\![s]\!] \circledast \subseteq^{-1})\} \\
& \qquad \text{since } \wp(\mathcal{S}[\![s]\!]) \in \gamma^{\mathfrak{S}}(\mathtt{P}\ \mathcal{S}[\![s]\!]),\ \subseteq\ \in \gamma^{\mathfrak{O}}(\subseteq),\ \supseteq\ \in \gamma^{\mathfrak{O}}((\subseteq)^{-1}), \text{ and } \gamma^{\mathfrak{P}}(\mathsf{S} \circledast \mathsf{O}) \triangleq \gamma^{\mathfrak{S}}(\mathsf{S}) \times \gamma^{\mathfrak{O}}(\mathsf{O}) \\
= & \gamma^{\mathfrak{T}}(\mathtt{P}\ \mathcal{S}[\![s]\!] \circledast \subseteq \leftrightharpoons \mathtt{P}\ \mathcal{S}[\![s]\!] \circledast \subseteq^{-1}) \qquad \text{def. } \gamma^{\mathfrak{T}}
\end{array}
$$

$= \gamma^{\mathfrak{T}}(\{\mathcal{S}[\![s]\!] \neq \mathtt{err} \mathbin{?} \mathtt{P}\ \mathcal{S}[\![s]\!] \circledast \subseteq \rightleftharpoons \mathtt{P}\ \mathcal{S}[\![s]\!] \circledast \subseteq^{-1} \mathbin{:} \mathtt{err}\})$
$\subseteq \gamma^{\mathfrak{T}}(\mathcal{T}[\![\neg[s]]\!]) \cup \{\omega\}$ *QED.*

- $\mathcal{T}[\![\infty[s]]\!] \triangleq \{\mathcal{S}[\![s]\!] \neq \mathtt{err} \mathbin{?} \mathtt{P}\ (\mathtt{seq}\ \mathcal{S}[\![s]\!]) \circledast \subseteq \rightleftharpoons \mathtt{P}\ \mathcal{S}[\![s]\!] \circledast \subseteq \mathbin{:} \mathtt{err}\}$

  In the non-trivial case $\mathcal{T}[\![\infty[s]]\!] \neq \mathtt{err}$ (so $\mathcal{S}[\![s]\!] \neq \mathtt{err}$) and $\mathcal{S}[\![s]\!] \neq \omega$, we have $\mathcal{S}[\![s^\infty]\!] \triangleq \{\mathcal{S}[\![s]\!] \notin \{\omega, \Omega\} \mathbin{?} \mathcal{S}[\![s]\!]^\infty \mathbin{:} \mathcal{S}[\![s]\!]\} = \mathcal{S}[\![s]\!]^\infty$. By structural induction hypothesis, $\mathcal{S}[\![s]\!] \in \gamma^{\mathfrak{S}}(\mathcal{S}[\![s]\!])$ so $\wp(\mathcal{S}[\![s]\!]) \in \gamma^{\mathfrak{S}}(\mathtt{P}\ \mathcal{S}[\![s]\!])$ hence $\langle \wp(\mathcal{S}[\![s]\!]), \subseteq\rangle \in \gamma^{\mathfrak{P}}(\mathtt{P}\ \mathcal{S}[\![s]\!] \circledast \subseteq)$. Similarly $\mathcal{S}[\![s]\!] \in \gamma^{\mathfrak{S}}(\mathcal{S}[\![s]\!])$ implies that $\mathcal{S}[\![s]\!]^\infty \in \{X^\infty \mid X \in \gamma^{\mathfrak{S}}(\mathcal{S}[\![s]\!])\} \triangleq \gamma^{\mathfrak{S}}(\mathtt{seq}\ \mathcal{S}[\![s]\!])$ so $\wp(\mathcal{S}[\![s]\!]^\infty) \in \gamma^{\mathfrak{S}}(\mathtt{P}(\mathtt{seq}\ \mathcal{S}[\![s]\!]))$ hence $\langle \wp(\mathcal{S}[\![s]\!]^\infty), \subseteq\rangle \in \gamma^{\mathfrak{P}}((\mathtt{P}(\mathtt{seq}\ \mathcal{S}[\![s]\!])) \circledast \subseteq)$. We conclude that
  $\mathcal{S}[\![\infty[s]]\!]$
  $\triangleq \{\mathcal{S}[\![s]\!] \notin \{\omega, \Omega\} \mathbin{?} \langle \wp(\mathcal{S}[\![s]\!]^\infty), \subseteq\rangle \overset{\gamma^\infty}{\underset{\alpha^\infty}{\leftrightarrows}} \langle \wp(\mathcal{S}[\![s]\!]), \subseteq\rangle \mathbin{:} \mathcal{S}[\![s]\!]\}$
  $= \langle \wp(\mathcal{S}[\![s]\!]^\infty), \subseteq\rangle \overset{\gamma^\infty}{\underset{\alpha^\infty}{\leftrightarrows}} \langle \wp(\mathcal{S}[\![s]\!]), \subseteq\rangle$
  $\in \{P \overset{\gamma}{\underset{\alpha}{\leftrightarrows}} P' \mid P \in \gamma^{\mathfrak{P}}(\mathtt{P}\ (\mathtt{seq}\ \mathcal{S}[\![s]\!]) \circledast \subseteq) \wedge P' \in \gamma^{\mathfrak{P}}(\mathtt{P}\ \mathcal{S}[\![s]\!] \circledast \subseteq)\}$
  $\triangleq \gamma^{\mathfrak{T}}(\mathtt{P}\ (\mathtt{seq}\ \mathcal{S}[\![s]\!]) \circledast \subseteq \rightleftharpoons \mathtt{P}\ \mathcal{S}[\![s]\!] \circledast \subseteq)$
  $= \gamma^{\mathfrak{T}}(\{\mathcal{S}[\![s]\!] \neq \mathtt{err} \mathbin{?} \mathtt{P}\ (\mathtt{seq}\ \mathcal{S}[\![s]\!]) \circledast \subseteq \rightleftharpoons \mathtt{P}\ \mathcal{S}[\![s]\!] \circledast \subseteq \mathbin{:} \mathtt{err}\})$
  $\triangleq \gamma^{\mathfrak{T}}(\mathcal{T}[\![\infty[s]]\!])$ *QED.*

- $\mathcal{T}[\![\rightsquigarrow[s_1, s_2]]\!] \triangleq \{\mathcal{S}[\![s_1]\!] \neq \mathtt{err} \wedge \mathcal{S}[\![s_2]\!] \neq \mathtt{err} \mathbin{?} \mathtt{P}(\mathcal{S}[\![s_1]\!] * \mathcal{S}[\![s_2]\!]) \circledast \subseteq \rightleftharpoons \mathtt{P}\ \mathcal{S}[\![s_1]\!] \twoheadrightarrow \mathtt{P}\ \mathcal{S}[\![s_2]\!] \circledast \dot{\subseteq} \mathbin{:} \mathtt{err}\}$

  In the non-trivial case $\mathcal{T}[\![\rightsquigarrow[s_1, s_2]]\!] \neq \mathtt{err}$, $\mathcal{S}[\![s_1]\!] \in \gamma^{\mathfrak{S}}(\mathcal{S}[\![s_1]\!])$, and $\mathcal{S}[\![s_2]\!] \in \gamma^{\mathfrak{S}}(\mathcal{S}[\![s_2]\!])$, we have $\mathcal{S}[\![s_1]\!] \times \mathcal{S}[\![s_2]\!] \in \{\langle X, Y\rangle \mid X \in \gamma^{\mathfrak{S}}(\mathcal{S}[\![s_1]\!]) \wedge Y \in \gamma^{\mathfrak{S}}(\mathcal{S}[\![s_2]\!])\} \triangleq \gamma^{\mathfrak{S}}(\mathcal{S}[\![s_1]\!] * \mathcal{S}[\![s_2]\!])$ and so $\wp(\mathcal{S}[\![s_1]\!] \times \mathcal{S}[\![s_2]\!]) \in \gamma^{\mathfrak{S}}(\mathtt{P}\ (\mathcal{S}[\![s_1]\!] * \mathcal{S}[\![s_2]\!]))$. Similarly, $\wp(\mathcal{S}[\![s_1]\!]) \in \gamma^{\mathfrak{S}}(\mathtt{P}\ \mathcal{S}[\![s_1]\!])$ and $\wp(\mathcal{S}[\![s_2]\!]) \in \gamma^{\mathfrak{S}}(\mathtt{P}\ \mathcal{S}[\![s_2]\!])$ so $\wp(\mathcal{S}[\![s_1]\!]) \xrightarrow{\cup} \wp(\mathcal{S}[\![s_2]\!]) \in \{X \mapsto Y \mid X \in \gamma^{\mathfrak{S}}(\mathtt{P}\ \mathcal{S}[\![s_1]\!]) \wedge Y \in \gamma^{\mathfrak{S}}(\mathtt{P}\ \mathcal{S}[\![s_2]\!])\} \triangleq \gamma^{\mathfrak{S}}(\mathtt{P}\ \mathcal{S}[\![s_1]\!] \twoheadrightarrow \mathtt{P}\ \mathcal{S}[\![s_2]\!])$. It follows that
  $\mathcal{S}[\![\rightsquigarrow[s_1, s_2]]\!]$
  $\triangleq \{\mathcal{S}[\![s_1]\!] \notin \{\omega, \Omega\} \wedge \mathcal{S}[\![s_1]\!] \notin \{\omega, \Omega\} \mathbin{?} \langle \wp(\mathcal{S}[\![s_1]\!] \times \mathcal{S}[\![s_2]\!]), \subseteq\rangle \overset{\gamma^{\rightsquigarrow}}{\underset{\alpha^{\rightsquigarrow}}{\leftrightarrows}} \langle \wp(\mathcal{S}[\![s_1]\!]) \xrightarrow{\cup} \wp(\mathcal{S}[\![s_2]\!]), \dot{\subseteq}\rangle \mathbin{:} \mathtt{error}(\mathcal{S}[\![s_1]\!], \mathcal{S}[\![s_2]\!])\}$
  $= \langle \wp(\mathcal{S}[\![s_1]\!] \times \mathcal{S}[\![s_2]\!]), \subseteq\rangle \overset{\gamma^{\rightsquigarrow}}{\underset{\alpha^{\rightsquigarrow}}{\leftrightarrows}} \langle \wp(\mathcal{S}[\![s_1]\!]) \xrightarrow{\cup} \wp(\mathcal{S}[\![s_2]\!]), \dot{\subseteq}\rangle$
  $\in \{P \overset{\gamma}{\underset{\alpha}{\leftrightarrows}} P' \mid P \in \gamma^{\mathfrak{P}}(\mathtt{P}(\mathcal{S}[\![s_1]\!] * \mathcal{S}[\![s_2]\!]) \circledast \subseteq) \wedge P' \in \gamma^{\mathfrak{P}}(\mathtt{P}\ \mathcal{S}[\![s_1]\!] \twoheadrightarrow \mathtt{P}\ \mathcal{S}[\![s_2]\!])\}$
  $= \gamma^{\mathfrak{T}}(\mathtt{P}(\mathcal{S}[\![s_1]\!] * \mathcal{S}[\![s_2]\!]) \circledast \subseteq \rightleftharpoons \mathtt{P}\ \mathcal{S}[\![s_1]\!] \twoheadrightarrow \mathtt{P}\ \mathcal{S}[\![s_2]\!] \circledast \dot{\subseteq})$
  $= \gamma^{\mathfrak{T}}(\{\mathcal{S}[\![s_1]\!] \neq \mathtt{err} \wedge \mathcal{S}[\![s_2]\!] \neq \mathtt{err} \mathbin{?} \mathtt{P}(\mathcal{S}[\![s_1]\!] * \mathcal{S}[\![s_2]\!]) \circledast \subseteq \rightleftharpoons \mathtt{P}\ \mathcal{S}[\![s_1]\!] \twoheadrightarrow \mathtt{P}\ \mathcal{S}[\![s_2]\!] \circledast \dot{\subseteq} \mathbin{:} \mathtt{err}\})$
  $\subseteq \gamma^{\mathfrak{T}}(\mathcal{T}[\![\rightsquigarrow[s_1, s_2]]\!]) \cup \{\omega\}$ *QED.*

- $\mathcal{T}[\![\mapsto[s_1, s_2]]\!] \triangleq \{\mathcal{S}[\![s_1]\!] \neq \mathtt{err} \wedge \mathcal{S}[\![s_2]\!] \neq \mathtt{err} \mathbin{?} \mathtt{P}(\mathcal{S}[\![s_1]\!] \twoheadrightarrow \mathcal{S}[\![s_2]\!]) \circledast \subseteq \rightleftharpoons \mathtt{P}\ \mathcal{S}[\![s_1]\!] \twoheadrightarrow \mathtt{P}\ \mathcal{S}[\![s_2]\!] \circledast \dot{\subseteq} \mathbin{:} \mathtt{err}\}$

  In the non-trivial case $\mathcal{T}[\![\mapsto[s_1, s_2]]\!] \neq \mathtt{err}$, $\mathcal{S}[\![s_1]\!] \in \gamma^{\mathfrak{S}}(\mathcal{S}[\![s_1]\!])$, and $\mathcal{S}[\![s_2]\!] \in \gamma^{\mathfrak{S}}(\mathcal{S}[\![s_2]\!])$, we have $\mathcal{S}[\![s_1]\!] \mapsto \mathcal{S}[\![s_2]\!] \in \{X \mapsto Y \mid X \in \gamma^{\mathfrak{S}}(\mathcal{S}[\![s_1]\!]) \wedge Y \in \gamma^{\mathfrak{S}}(\mathcal{S}[\![s_2]\!])\} \triangleq \gamma^{\mathfrak{S}}(\mathcal{S}[\![s_1]\!] \twoheadrightarrow \mathcal{S}[\![s_2]\!])$ so $\wp(\mathcal{S}[\![s_1]\!] \mapsto \mathcal{S}[\![s_2]\!]) \in \gamma^{\mathfrak{S}}(\mathtt{P}\ (\mathcal{S}[\![s_1]\!] \twoheadrightarrow \mathcal{S}[\![s_2]\!]))$. Moreover $\wp(\mathcal{S}[\![s_1]\!]) \in \gamma^{\mathfrak{S}}(\mathtt{P}\ \mathcal{S}[\![s_1]\!])$ and $\wp(\mathcal{S}[\![s_2]\!]) \in \gamma^{\mathfrak{S}}(\mathtt{P}\ \mathcal{S}[\![s_1]\!])$ imply that $\wp(\mathcal{S}[\![s_1]\!]) \mapsto \wp(\mathcal{S}[\![s_2]\!]) \in \{X \mapsto Y \mid X \in \gamma^{\mathfrak{S}}(\mathtt{P}\ \mathcal{S}[\![s_1]\!]) \wedge Y \in \gamma^{\mathfrak{S}}(\mathtt{P}\ \mathcal{S}[\![s_1]\!])\} \triangleq \gamma^{\mathfrak{S}}((\mathtt{P}\ \mathcal{S}[\![s_1]\!]) \twoheadrightarrow (\mathtt{P}\ \mathcal{S}[\![s_1]\!]))$. It follows that
  $\mathcal{S}[\![\mapsto[s_1, s_2]]\!]$
  $\triangleq \{\mathcal{S}[\![s_1]\!] \notin \{\omega, \Omega\} \wedge \mathcal{S}[\![s_2]\!] \notin \{\omega, \Omega\} \mathbin{?} \langle \wp(\mathcal{S}[\![s_1]\!] \mapsto \mathcal{S}[\![s_2]\!]), \subseteq\rangle \overset{\gamma^{\mapsto}}{\underset{\alpha^{\mapsto}}{\leftrightarrows}} \langle \wp(\mathcal{S}[\![s_1]\!]) \mapsto \wp(\mathcal{S}[\![s_2]\!]), \dot{\subseteq}\rangle \mathbin{:} \mathtt{error}(\mathcal{S}[\![s_1]\!], \mathcal{S}[\![s_2]\!])\}$
  $= \langle \wp(\mathcal{S}[\![s_1]\!] \mapsto \mathcal{S}[\![s_2]\!]), \subseteq\rangle \overset{\gamma^{\mapsto}}{\underset{\alpha^{\mapsto}}{\leftrightarrows}} \langle \wp(\mathcal{S}[\![s_1]\!]) \mapsto \wp(\mathcal{S}[\![s_2]\!]), \dot{\subseteq}\rangle$
  $\in \{P \overset{\gamma}{\underset{\alpha}{\leftrightarrows}} P' \mid P \in \gamma^{\mathfrak{P}}(\mathtt{P}(\mathcal{S}[\![s_1]\!] \twoheadrightarrow \mathcal{S}[\![s_2]\!]) \circledast \subseteq) \wedge P' \in \gamma^{\mathfrak{P}}(\mathtt{P}\ \mathcal{S}[\![s_1]\!] \twoheadrightarrow \mathtt{P}\ \mathcal{S}[\![s_2]\!] \circledast \dot{\subseteq})\}$
  $= \gamma^{\mathfrak{T}}(\mathtt{P}(\mathcal{S}[\![s_1]\!] \twoheadrightarrow \mathcal{S}[\![s_2]\!]) \circledast \subseteq \rightleftharpoons \mathtt{P}\ \mathcal{S}[\![s_1]\!] \twoheadrightarrow \mathtt{P}\ \mathcal{S}[\![s_2]\!] \circledast \dot{\subseteq})$
  $= \gamma^{\mathfrak{T}}(\{\mathcal{S}[\![s_1]\!] \neq \mathtt{err} \wedge \mathcal{S}[\![s_2]\!] \neq \mathtt{err} \mathbin{?} \mathtt{P}(\mathcal{S}[\![s_1]\!] \twoheadrightarrow \mathcal{S}[\![s_2]\!]) \circledast \subseteq \rightleftharpoons \mathtt{P}\ \mathcal{S}[\![s_1]\!] \twoheadrightarrow \mathtt{P}\ \mathcal{S}[\![s_2]\!] \circledast \dot{\subseteq} \mathbin{:} \mathtt{err}\})$
  $\triangleq \gamma^{\mathfrak{T}}(\mathcal{T}[\![\mapsto[s_1, s_2]]\!])$ *QED.*

- $\mathcal{T}[\![\times[s_1, s_2]]\!] \triangleq \{\mathcal{S}[\![s_1]\!] \neq \mathtt{err} \wedge \mathcal{S}[\![s_2]\!] \neq \mathtt{err} \mathbin{?} \mathtt{P}(\mathcal{S}[\![s_1]\!] \twoheadrightarrow \mathcal{S}[\![s_2]\!]) \circledast \subseteq \rightleftharpoons \mathcal{S}[\![s_1]\!] \twoheadrightarrow \mathtt{P}\ \mathcal{S}[\![s_2]\!] \circledast \dot{\subseteq} \mathbin{:} \mathtt{err}\}$

  In the only non-trivial case $\mathcal{T}[\![\times[s_1, s_2]]\!] \neq \mathtt{err}$, $\mathcal{S}[\![s_1]\!], \mathcal{S}[\![s_2]\!] \notin \{\omega, \Omega\}$, we have $\mathcal{S}[\![s_i]\!] \in \gamma^{\mathfrak{S}}(\mathcal{S}[\![s_i]\!])$, $i = 1, 2$ by structural induction hypothesis. It follows that $\mathcal{S}[\![s_1]\!] \mapsto \mathcal{S}[\![s_2]\!] \in \{X \mapsto Y \mid X \in \gamma^{\mathfrak{S}}(\mathcal{S}[\![s_1]\!]) \wedge Y \in \gamma^{\mathfrak{S}}(\mathcal{S}[\![s_2]\!])\} \triangleq \gamma^{\mathfrak{S}}(\mathcal{S}[\![s_1]\!] \twoheadrightarrow \mathcal{S}[\![s_2]\!])$ so $\wp(\mathcal{S}[\![s_1]\!] \mapsto \mathcal{S}[\![s_2]\!]) \in \gamma^{\mathfrak{S}}(\mathtt{P}\ (\mathcal{S}[\![s_1]\!] \twoheadrightarrow \mathcal{S}[\![s_2]\!]))$ hence $\langle \wp(\mathcal{S}[\![s_1]\!] \mapsto \mathcal{S}[\![s_2]\!]), \subseteq\rangle \in \gamma^{\mathfrak{P}}(\mathtt{P}(\mathcal{S}[\![s_1]\!] \twoheadrightarrow \mathcal{S}[\![s_2]\!]) \circledast \subseteq)$. Moreover $\mathcal{S}[\![s_2]\!] \in \gamma^{\mathfrak{S}}(\mathcal{S}[\![s_2]\!])$ implies $\wp(\mathcal{S}[\![s_2]\!]) \in \gamma^{\mathfrak{S}}(\mathtt{P}\ \mathcal{S}[\![s_2]\!])$ so $\mathcal{S}[\![s_1]\!] \mapsto \wp(\mathcal{S}[\![s_2]\!]) \in \{X \mapsto Y \mid X \in \gamma^{\mathfrak{S}}(\mathcal{S}[\![s_1]\!]) \wedge Y \in \gamma^{\mathfrak{S}}(\mathtt{P}\ \mathcal{S}[\![s_2]\!])\} \triangleq \gamma^{\mathfrak{S}}(\mathcal{S}[\![s_1]\!] \twoheadrightarrow \mathtt{P}\ \mathcal{S}[\![s_2]\!])$ hence $\langle \mathcal{S}[\![s_1]\!] \mapsto \wp(\mathcal{S}[\![s_2]\!]), \dot{\subseteq}\rangle \in \gamma^{\mathfrak{P}}(\mathcal{S}[\![s_1]\!] \twoheadrightarrow \mathtt{P}\ \mathcal{S}[\![s_2]\!] \circledast \dot{\subseteq})$. We conclude that
  $\mathcal{S}[\![\times[s_1, s_2]]\!]$
  $\triangleq \{\mathcal{S}[\![s_1]\!] \notin \{\omega, \Omega\} \wedge \mathcal{S}[\![s_2]\!] \notin \{\omega, \Omega\} \mathbin{?} \langle \wp(\mathcal{S}[\![s_1]\!] \mapsto \mathcal{S}[\![s_2]\!]), \subseteq\rangle \overset{\gamma^{\times}}{\underset{\alpha^{\times}}{\leftrightarrows}} \langle \mathcal{S}[\![s_1]\!] \mapsto \wp(\mathcal{S}[\![s_2]\!]), \dot{\subseteq}\rangle \mathbin{:} \mathtt{error}(\mathcal{S}[\![s_1]\!], \mathcal{S}[\![s_2]\!])\}$
  $= \langle \wp(\mathcal{S}[\![s_1]\!] \mapsto \mathcal{S}[\![s_2]\!]), \subseteq\rangle \overset{\gamma^{\times}}{\underset{\alpha^{\times}}{\leftrightarrows}} \langle \mathcal{S}[\![s_1]\!] \mapsto \wp(\mathcal{S}[\![s_2]\!]), \dot{\subseteq}\rangle$
  $\in \{P \overset{\gamma}{\underset{\alpha}{\leftrightarrows}} P' \mid P \in \gamma^{\mathfrak{P}}(\mathtt{P}(\mathcal{S}[\![s_1]\!] \twoheadrightarrow \mathcal{S}[\![s_2]\!]) \circledast \subseteq) \wedge P' \in \gamma^{\mathfrak{P}}(\mathcal{S}[\![s_1]\!] \twoheadrightarrow \mathtt{P}\ \mathcal{S}[\![s_2]\!] \circledast \dot{\subseteq})\}$
  $\triangleq \gamma^{\mathfrak{T}}(\mathtt{P}(\mathcal{S}[\![s_1]\!] \twoheadrightarrow \mathcal{S}[\![s_2]\!]) \circledast \subseteq \rightleftharpoons \mathcal{S}[\![s_1]\!] \twoheadrightarrow \mathtt{P}\ \mathcal{S}[\![s_2]\!] \circledast \dot{\subseteq})$
  $= \gamma^{\mathfrak{T}}(\{\mathcal{S}[\![s_1]\!] \neq \mathtt{err} \wedge \mathcal{S}[\![s_2]\!] \neq \mathtt{err} \mathbin{?} \mathtt{P}(\mathcal{S}[\![s_1]\!] \twoheadrightarrow \mathcal{S}[\![s_2]\!]) \circledast \subseteq \rightleftharpoons \mathcal{S}[\![s_1]\!] \twoheadrightarrow \mathtt{P}\ \mathcal{S}[\![s_2]\!] \circledast \dot{\subseteq} \mathbin{:} \mathtt{err}\})$
  $\triangleq \gamma^{\mathfrak{T}}(\mathcal{T}[\![\times[s_1, s_2]]\!])$ *QED.*

- $\mathcal{T}[\![\mathsf{R}[g]]\!] \triangleq \mathcal{T}[\![g]\!]$

  In case $\mathcal{T}[\![\mathsf{R}[g]]\!] \neq \mathtt{err}$ and $\mathcal{S}[\![\mathsf{R}[g]]\!] \notin \{\omega, \Omega\}$, we have
  $\mathcal{S}[\![\mathsf{R}[g]]\!]$
  $\triangleq \{\mathcal{S}[\![g]\!] = \langle \mathcal{C}, \sqsubseteq\rangle \overset{\gamma}{\underset{\alpha}{\leftrightarrows}} \langle \mathcal{A}, \leqslant\rangle \mathbin{?} \langle \mathcal{C}, \sqsubseteq\rangle \overset{\gamma}{\underset{\alpha}{\leftrightarrows}} \langle \{\alpha(P) \mid P \in \mathcal{C}\}, \leqslant\rangle \mathbin{:} (\mathcal{S}[\![g]\!] = \omega \mathbin{?} \omega \mathbin{:} \Omega)\}$
  $= \langle \mathcal{C}, \sqsubseteq\rangle \overset{\gamma}{\underset{\alpha}{\leftrightarrows}} \langle \{\alpha(P) \mid P \in \mathcal{C}\}, \leqslant\rangle$
  By structural induction hypothesis, $\mathcal{S}[\![g]\!] = \langle \mathcal{C}, \sqsubseteq\rangle \overset{\gamma}{\underset{\alpha}{\leftrightarrows}} \langle \mathcal{A}, \leqslant\rangle \in \gamma^{\mathfrak{T}}(\mathcal{T}[\![g]\!])$ so $\mathcal{T}[\![g]\!] = \mathsf{P} \rightleftharpoons \mathsf{P}'$ with $\langle \mathcal{C}, \sqsubseteq\rangle \overset{\gamma}{\underset{\alpha}{\leftrightarrows}} \langle \mathcal{A}, \leqslant\rangle \in \gamma^{\mathfrak{T}}(\mathsf{P} \rightleftharpoons \mathsf{P}') = \{P \overset{\gamma}{\underset{\alpha}{\leftrightarrows}} P' \mid P \in \gamma^{\mathfrak{P}}(\mathsf{P}) \wedge P' \in \gamma^{\mathfrak{P}}(\mathsf{P}')\}$. Necessarily, $P' = \mathsf{S}' \circledast \mathsf{O}'$ with $\langle \mathcal{A}, \leqslant\rangle \in \gamma^{\mathfrak{P}}(\mathsf{P}') = \gamma^{\mathfrak{S}}(\mathsf{S}') \times$

$\gamma^{\mathfrak{O}}(\mathsf{O}')$ so $\mathcal{A} \in \gamma^{\mathfrak{S}}(\mathsf{S}')$. But $\{\alpha(P) \mid P \in \mathcal{C}\} \subseteq \mathcal{A}$ and $\gamma^{\mathfrak{S}}(\mathsf{S}')$ is $\subseteq$-downward closed so $\{\alpha(P) \mid P \in \mathcal{C}\} \in \gamma^{\mathfrak{S}}(\mathsf{S}')$ and $\langle \{\alpha(P) \mid P \in \mathcal{C}\}, \leqslant \rangle \in \gamma^{\mathfrak{S}}(\mathsf{S}') \times \gamma^{\mathfrak{O}}(\mathsf{O}') = \gamma^{\mathfrak{P}}(\mathsf{P}')$. It follows that

$\langle \mathcal{C}, \sqsubseteq \rangle \underset{\alpha}{\overset{\gamma}{\leftrightarrows}} \langle \{\alpha(P) \mid P \in \mathcal{C}\}, \leqslant \rangle$
$\in \{P \underset{\alpha}{\overset{\gamma}{\leftrightarrows}} P' \mid P \in \gamma^{\mathfrak{P}}(\mathsf{P}) \wedge P' \in \gamma^{\mathfrak{P}}(\mathsf{P}')\}$
$\triangleq \gamma^{\mathfrak{T}}(\mathsf{P} \leftrightharpoons \mathsf{P}')$
$\triangleq \gamma^{\mathfrak{T}}(\mathcal{T}[\![\mathsf{R}[g]]\!])$ *QED.*

- $\mathcal{T}[\![s \to g]\!] \triangleq \big(\mathcal{S}[\![s]\!] \neq \mathtt{err} \wedge \mathcal{T}[\![g]\!] \neq \mathtt{err} \;?\; \mathcal{S}[\![s]\!] \mathbin{\ast\!\!\to} \mathcal{T}[\![g]\!] \,;\, \mathtt{err}\big)$

In the non trivial case $\mathcal{T}[\![s \to g]\!] \neq \mathtt{err}$, $\mathcal{S}[\![s]\!] \notin \{\omega, \Omega\}$, and $\mathcal{S}[\![g]\!] = \langle \mathcal{C}, \sqsubseteq \rangle \underset{\alpha}{\overset{\gamma}{\leftrightarrows}} \langle \mathcal{A}, \leqslant \rangle \notin \{\omega, \Omega\}$, we have $\mathcal{S}[\![s]\!] \in \gamma^{\mathfrak{S}}(\mathcal{S}[\![s]\!])$ and $\langle \mathcal{C}, \sqsubseteq \rangle \underset{\alpha}{\overset{\gamma}{\leftrightarrows}} \langle \mathcal{A}, \leqslant \rangle \in \gamma^{\mathfrak{T}}(\mathcal{T}[\![g]\!])$ by structural induction hypothesis, so that we conclude that

$\mathcal{S}[\![s \to g]\!]$
$\triangleq \big(\mathcal{S}[\![s]\!] = X \notin \{\omega, \Omega\} \wedge \mathcal{S}[\![g]\!] = \langle \mathcal{C}, \sqsubseteq \rangle \underset{\alpha}{\overset{\gamma}{\leftrightarrows}} \langle \mathcal{A}, \leqslant \rangle \;?\; \langle X \mapsto \mathcal{C}, \dot{\sqsubseteq} \rangle \underset{\boldsymbol{\lambda}\,\rho \bullet \alpha \circ \rho}{\overset{\boldsymbol{\lambda}\,\overline{\rho} \bullet \gamma \circ \overline{\rho}}{\leftrightarrows}} \langle X \mapsto \mathcal{A}, \dot{\leqslant} \rangle \,;\, \mathtt{error}(\mathcal{S}[\![s]\!], \big(\mathcal{S}[\![g]\!] = \omega \;?\; \omega \,;\, \Omega\big))\big)$
$= \langle \mathcal{S}[\![s]\!] \mapsto \mathcal{C}, \dot{\sqsubseteq} \rangle \underset{\boldsymbol{\lambda}\,\rho \bullet \alpha \circ \rho}{\overset{\boldsymbol{\lambda}\,\overline{\rho} \bullet \gamma \circ \overline{\rho}}{\leftrightarrows}} \langle \mathcal{S}[\![s]\!] \mapsto \mathcal{A}, \dot{\leqslant} \rangle$
$\in \{\langle X \mapsto \mathcal{C}, \dot{\sqsubseteq} \rangle \underset{\alpha'}{\overset{\gamma'}{\leftrightarrows}} \langle X \mapsto \mathcal{A}, \dot{\leqslant} \rangle \mid X \in \gamma^{\mathfrak{S}}(\mathcal{S}[\![s]\!]) \wedge \langle \mathcal{C}, \sqsubseteq \rangle \underset{\alpha}{\overset{\gamma}{\leftrightarrows}} \langle \mathcal{A}, \leqslant \rangle \in \gamma^{\mathfrak{T}}(\mathcal{T}[\![g]\!])\}$
$\triangleq \gamma^{\mathfrak{T}}(\mathcal{S}[\![s]\!] \mathbin{\ast\!\!\to} \mathcal{T}[\![g]\!])$
$\triangleq \gamma^{\mathfrak{T}}(\mathcal{T}[\![s \to g]\!])$ *QED.*

- $\mathcal{T}[\![g_1 \,;\, g_2]\!] \triangleq \big(\mathcal{T}[\![g_1]\!] = \mathsf{P}_1 \leftrightharpoons \mathsf{P}_2 \wedge \mathcal{T}[\![g_2]\!] = \mathsf{P}_3 \leftrightharpoons \mathsf{P}_4 \wedge \mathsf{P}_2 \cong \mathsf{P}_3 \;?\; \mathsf{P}_1 \leftrightharpoons \mathsf{P}_4 \,;\, \mathtt{err}\big)$

In case $\mathcal{T}[\![g_1 \,;\, g_2]\!] \neq \mathtt{err}$ and $\mathcal{S}[\![g_1]\!], \mathcal{S}[\![g_2]\!] \notin \{\Omega, \omega\}$, we have, by structural induction hypothesis, that $\mathcal{S}[\![g_1]\!] = p_1 \underset{\alpha_1}{\overset{\gamma_1}{\leftrightarrows}} p_2 \in \gamma^{\mathfrak{T}}(\mathcal{T}[\![g_1]\!]) = \gamma^{\mathfrak{T}}(\mathsf{P}_1 \leftrightharpoons \mathsf{P}_2) = \{P_1 \underset{\alpha_1}{\overset{\gamma_1}{\leftrightarrows}} P_2 \mid P_1 \in \gamma^{\mathfrak{P}}(\mathsf{P}_1) \wedge P_2 \in \gamma^{\mathfrak{P}}(\mathsf{P}_2)\}$ so $p_1 \in \gamma^{\mathfrak{P}}(\mathsf{P}_1)$ and $\mathcal{S}[\![g_2]\!] = p_3 \underset{\alpha_2}{\overset{\gamma_2}{\leftrightarrows}} p_4 \in \gamma^{\mathfrak{T}}(\mathcal{T}[\![g_2]\!]) = \gamma^{\mathfrak{T}}(\mathsf{P}_3 \leftrightharpoons \mathsf{P}_4) = \{P_3 \underset{\alpha_2}{\overset{\gamma_2}{\leftrightarrows}} P_4 \mid P_3 \in \gamma^{\mathfrak{P}}(\mathsf{P}_3) \wedge P_4 \in \gamma^{\mathfrak{P}}(\mathsf{P}_4)\}$ so $p_4 \in \gamma^{\mathfrak{P}}(\mathsf{P}_4)$. Moreover $\mathsf{P}_4 \cong \mathsf{P}_4$ so $p_2 \in \gamma^{\mathfrak{P}}(\mathsf{P}_2) = \gamma^{\mathfrak{P}}(\mathsf{P}_3) \ni p_3$. Nevertheless this type requirement is too weak to ensure that $p_2 = p_3$. When $p_2 \neq p_3$, the semantics returns a dynamic error ω allowed by the type system. Otherwise $p_2 = p_3$, in which case

$\mathcal{S}[\![g_1 \,;\, g_2]\!]$
$\triangleq \big(\mathcal{S}[\![g_1]\!] = p_1 \underset{\alpha_1}{\overset{\gamma_1}{\leftrightarrows}} p_2 \wedge \mathcal{S}[\![g_2]\!] = p_3 \underset{\alpha_2}{\overset{\gamma_2}{\leftrightarrows}} p_4 \;?\; \big(p_2 = p_3 \;?\; p_1 \underset{\alpha_2 \circ \alpha_1}{\overset{\gamma_1 \circ \gamma_2}{\leftrightarrows}} p_4 \,;\, \omega\big) \,;\, \mathtt{error}(\big(\mathcal{S}[\![g_1]\!] = \omega \;?\; \omega \,;\, \Omega\big)), \big(\mathcal{S}[\![g_2]\!] = \omega \;?\; \omega \,;\, \Omega\big)\big)\big)$
$= p_1 \underset{\alpha_2 \circ \alpha_1}{\overset{\gamma_1 \circ \gamma_2}{\leftrightarrows}} p_4$ ⟅in the non-trivial case⟆
$\in \{P_1 \underset{\alpha}{\overset{\gamma}{\leftrightarrows}} P_4 \mid P_1 \in \gamma^{\mathfrak{P}}(\mathsf{P}_1) \wedge P_4 \in \gamma^{\mathfrak{P}}(\mathsf{P}_4)\}$
⟅since $p_1 \in \gamma^{\mathfrak{P}}(\mathsf{P}_1)$ and $p_4 \in \gamma^{\mathfrak{P}}(\mathsf{P}_4)$⟆
$\triangleq \gamma^{\mathfrak{T}}(\mathsf{P}_1 \leftrightharpoons \mathsf{P}_4)$
$\triangleq \gamma^{\mathfrak{T}}(\mathcal{T}[\![g_1 \,;\, g_2]\!])$ *QED.*

- $\mathcal{T}[\![g_1 \Cap g_2]\!] \triangleq \big(\mathcal{T}[\![g_1]\!] = \mathsf{S}_1 \circledast \mathsf{O}_1 \leftrightharpoons \mathsf{S}_2 \circledast \mathsf{O}_2 \wedge \mathcal{T}[\![g_2]\!] = \mathsf{S}_3 \circledast \mathsf{O}_3 \leftrightharpoons \mathsf{S}_4 \circledast \mathsf{O}_4 \;?\; \mathsf{S}_1 \ast \mathsf{S}_2 \circledast \mathsf{O}_1 \star \mathsf{O}_2 \leftrightharpoons \mathsf{S}_3 \ast \mathsf{S}_4 \circledast \mathsf{O}_3 \star \mathsf{O}_4 \,;\, \mathtt{err}\big)$

In case $\mathcal{T}[\![g_1 \Cap g_2]\!] \neq \mathtt{err}$, $\mathcal{S}[\![g_1]\!] = \langle \mathcal{C}_1, \subseteq \rangle \underset{\alpha_1}{\overset{\gamma_1}{\leftrightarrows}} \langle \mathcal{A}_1, \sqsubseteq \rangle \neq \omega$, and $\mathcal{S}[\![g_2]\!] = \langle \mathcal{C}_2, \Subset \rangle \underset{\alpha_2}{\overset{\gamma_2}{\leftrightarrows}} \langle \mathcal{A}_2, \preceq \rangle \neq \omega$, we have $\langle \mathcal{C}_1, \subseteq \rangle \in \gamma^{\mathfrak{P}}(\mathsf{S}_1 \circledast \mathsf{O}_1)$, $\langle \mathcal{A}_1, \sqsubseteq \rangle \in \gamma^{\mathfrak{P}}(\mathsf{S}_2 \circledast \mathsf{O}_2)$, $\langle \mathcal{C}_2, \Subset \rangle \in \gamma^{\mathfrak{P}}(\mathsf{S}_3 \circledast \mathsf{O}_3)$, and $\langle \mathcal{A}_2, \preceq \rangle \in \gamma^{\mathfrak{P}}(\mathsf{S}_4 \circledast \mathsf{O}_4)$, by structural induction hypothesis.

It follows that $\mathcal{C}_1 \in \gamma^{\mathfrak{S}}(\mathsf{S}_1)$, $\mathcal{C}_2 \in \gamma^{\mathfrak{S}}(\mathsf{S}_2)$, $\subseteq \in \gamma^{\mathfrak{O}}(\mathsf{O}_1)$, $\Subset \in \gamma^{\mathfrak{O}}(\mathsf{O}_2)$, and therefore $\langle \mathcal{C}_1 \times \mathcal{C}_2, \subseteq \times \Subset \rangle \in (\{X \times Y \mid X \in \gamma^{\mathfrak{S}}(\mathsf{S}_1) \wedge Y \in \gamma^{\mathfrak{S}}(\mathsf{S}_2)\}) \times (\gamma^{\mathfrak{O}}(\mathsf{O}_1) \times \gamma^{\mathfrak{O}}(\mathsf{O}_2)) = \gamma^{\mathfrak{S}}(\mathsf{S}_1 \ast \mathsf{S}_2) \times \gamma^{\mathfrak{O}}(\mathsf{O}_1 \star \mathsf{O}_2) = \gamma^{\mathfrak{P}}(\mathsf{S}_1 \ast \mathsf{S}_2 \circledast \mathsf{O}_1 \star \mathsf{O}_2)$ and similarly $\langle \mathcal{A}_1 \times \mathcal{A}_2, \sqsubseteq \times \preceq \rangle \in \gamma^{\mathfrak{P}}(\mathsf{S}_3 \ast \mathsf{S}_4 \circledast \mathsf{O}_3 \star \mathsf{O}_4)$

We conclude that

$\mathcal{S}[\![g_1 \Cap g_2]\!]$
$\triangleq \langle \mathcal{C}_1 \times \mathcal{C}_2, \subseteq \times \Subset \rangle \underset{\alpha_1 \times \alpha_2}{\overset{\gamma_1 \times \gamma_2}{\leftrightarrows}} \langle \mathcal{A}_1 \times \mathcal{A}_2, \sqsubseteq \times \preceq \rangle$
$\in \{P \underset{\alpha}{\overset{\gamma}{\leftrightarrows}} P' \mid P \in \gamma^{\mathfrak{P}}(\mathsf{S}_1 \ast \mathsf{S}_2 \circledast \mathsf{O}_1 \star \mathsf{O}_2) \wedge P' \in \gamma^{\mathfrak{P}}(\mathsf{S}_3 \ast \mathsf{S}_4 \circledast \mathsf{O}_3 \star \mathsf{O}_4)\}$
$\triangleq \gamma^{\mathfrak{T}}(\mathcal{T}[\![\mathsf{S}_1 \ast \mathsf{S}_2 \circledast \mathsf{O}_1 \star \mathsf{O}_2 \leftrightharpoons \mathsf{S}_3 \ast \mathsf{S}_4 \circledast \mathsf{O}_3 \star \mathsf{O}_4]\!])$
$\triangleq \gamma^{\mathfrak{T}}(\mathcal{T}[\![g_1 \Cap g_2]\!])$ *QED.*

- $\mathcal{T}[\![g_1 \Longrightarrow g_2]\!] \triangleq \big(\mathcal{T}[\![g_1]\!] = \mathsf{S}_1 \circledast \mathsf{O}_1 \leftrightharpoons \mathsf{S}_2 \circledast \mathsf{O}_2 \wedge \mathcal{T}[\![g_2]\!] = \mathsf{S}_3 \circledast \mathsf{O}_3 \leftrightharpoons \mathsf{S}_4 \circledast \mathsf{O}_4 \;?\; \mathsf{S}_1 \mathbin{\ast\!\!\to} \mathsf{S}_3 \circledast \dot{\mathsf{O}}_3 \leftrightharpoons \mathsf{S}_2 \mathbin{\ast\!\!\to} \mathsf{S}_4 \circledast \dot{\mathsf{O}}_4 \,;\, \mathtt{err}\big)$

$\mathcal{T}[\![g_1 \Longrightarrow g_2]\!] \neq \mathtt{err}$, $\mathcal{S}[\![g_1]\!] = \langle \mathcal{C}_1, \subseteq \rangle \underset{\alpha_1}{\overset{\gamma_1}{\leftrightarrows}} \langle \mathcal{A}_1, \sqsubseteq \rangle \neq \omega$ so $\mathcal{C}_1 \in \gamma^{\mathfrak{S}}(\mathsf{S}_1)$, $\subseteq \in \gamma^{\mathfrak{O}}(\mathsf{O}_1)$, $\mathcal{A}_1 \in \gamma^{\mathfrak{S}}(\mathsf{S}_2)$, $\sqsubseteq \in \gamma^{\mathfrak{O}}(\mathsf{O}_2)$ by structural induction hypothesis, and similarly for $\mathcal{S}[\![g_2]\!] = \langle \mathcal{C}_2, \preceq \rangle \underset{\alpha_2}{\overset{\gamma_2}{\leftrightarrows}} \langle \mathcal{A}_2, \leqslant \rangle \neq \omega$ where $\preceq \in \gamma^{\mathfrak{O}}(\mathsf{O}_3)$ so $\dot{\preceq} \in \gamma^{\mathfrak{O}}(\dot{\mathsf{O}}_3)$ and $\leqslant \in \gamma^{\mathfrak{O}}(\mathsf{O}_4)$ so $\dot{\leqslant} \in \gamma^{\mathfrak{O}}(\dot{\mathsf{O}}_4)$. It follows that $\mathcal{C}_1 \xrightarrow{\nearrow} \mathcal{C}_2 \in \{X \mapsto Y \mid X \in \gamma^{\mathfrak{S}}(\mathsf{S}_1) \wedge Y \in \gamma^{\mathfrak{S}}(\mathsf{S}_3)\} \triangleq \gamma^{\mathfrak{S}}(\mathsf{S}_1 \mathbin{\ast\!\!\to} \mathsf{S}_3)$ and similarly $\mathcal{A}_1 \xrightarrow{\nearrow} \mathcal{A}_2 \in \gamma^{\mathfrak{S}}(\mathsf{S}_2 \mathbin{\ast\!\!\to} \mathsf{S}_4)$. We conclude that

$\mathcal{S}[\![g_1 \Longrightarrow g_2]\!]$
$\triangleq \langle \mathcal{C}_1 \xrightarrow{\nearrow} \mathcal{C}_2, \dot{\preceq} \rangle \underset{\boldsymbol{\lambda}\,f \bullet \alpha_2 \circ f \circ \gamma_1}{\overset{\boldsymbol{\lambda}\,g \bullet \gamma_2 \circ g \circ \alpha_1}{\leftrightarrows}} \langle \mathcal{A}_1 \xrightarrow{\nearrow} \mathcal{A}_2, \dot{\leqslant} \rangle$
$= \{\langle S, O \rangle \underset{\alpha}{\overset{\gamma}{\leftrightarrows}} \langle S', O' \rangle \mid S \in \gamma^{\mathfrak{S}}(\mathsf{S}_1 \mathbin{\ast\!\!\to} \mathsf{S}_3) \wedge O \in \gamma^{\mathfrak{O}}(\dot{\mathsf{O}}_3) \wedge S' \in \gamma^{\mathfrak{S}}(\mathsf{S}_2 \mathbin{\ast\!\!\to} \mathsf{S}_4) \wedge O' \in \gamma^{\mathfrak{O}}(\dot{\mathsf{O}}_4)\}$
$= \{P \underset{\alpha}{\overset{\gamma}{\leftrightarrows}} P' \mid P \in \gamma^{\mathfrak{P}}(\mathsf{S}_1 \mathbin{\ast\!\!\to} \mathsf{S}_3 \circledast \dot{\mathsf{O}}_3) \wedge P' \in \gamma^{\mathfrak{P}}(\mathsf{S}_2 \mathbin{\ast\!\!\to} \mathsf{S}_4 \circledast \dot{\mathsf{O}}_4)\}$
$\triangleq \gamma^{\mathfrak{T}}(\mathsf{S}_1 \mathbin{\ast\!\!\to} \mathsf{S}_3 \circledast \dot{\mathsf{O}}_3 \leftrightharpoons \mathsf{S}_2 \mathbin{\ast\!\!\to} \mathsf{S}_4 \circledast \dot{\mathsf{O}}_4)$
$\triangleq \gamma^{\mathfrak{T}}(\mathcal{T}[\![g_1 \Longrightarrow g_2]\!])$ *QED.*

11. Rule-based typing deduction system Abstract interpretations can always be presented by a suitable generalization of rule-based deduction systems [4].

We write $x \vdash \mathsf{T}$ to mean that $\mathscr{E}[\![x]\!]/\mathcal{S}[\![x]\!]/\mathcal{O}[\![x]\!]/\mathcal{P}[\![x]\!]/\mathcal{T}[\![x]\!] = \mathsf{T}$ (with discrimination on the syntax of x). The rule-based typing deduction system can only derive $x \vdash \mathsf{T}$ when $\mathscr{E}[\![x]\!]/\mathcal{S}[\![x]\!]/\mathcal{O}[\![x]\!]/\mathcal{P}[\![x]\!]/\mathcal{T}[\![x]\!] \neq \mathtt{err}$. So an expression which is not typable by the type system is understood to have type $\mathtt{err}$. Otherwise stated, $\mathtt{err}$ encodes untypable.

[4] Patrick Cousot, Radhia Cousot: Compositional and Inductive Semantic Definitions in Fixpoint, Equational, Constraint, Closure-condition, Rule-based and Game-Theoretic Form. CAV 1995: 293–308.

11.1 Elements

$\mathsf{true} \vdash \mathtt{bool}$ $\quad \mathsf{false} \vdash \mathtt{bool}$ $\quad 0 \vdash \mathtt{int}$ $\quad 1 \vdash \mathtt{int}$ $\quad \infty \vdash \mathtt{int}$

$\mathtt{x} \vdash \mathtt{var}$ $\quad \mathtt{y} \vdash \mathtt{var}$ $\quad \ell \vdash \mathtt{lab}$ $\quad \dfrac{e \vdash \mathtt{bool}}{-e \vdash \mathtt{bool}}$ $\quad \dfrac{e \vdash \mathtt{int}}{-e \vdash \mathtt{int}}$

11.2 Sets

$\mathbb{B} \vdash \mathtt{P\ bool}$ $\quad \mathbb{Z} \vdash \mathtt{P\ int}$ $\quad \mathbb{X} \vdash \mathtt{P\ var}$

$\mathbb{L} \vdash \mathtt{P\ lab}$

$$\frac{s \vdash \mathsf{S}}{\mathbb{I}(s,o) \vdash \mathtt{P}\ \mathsf{S}} \qquad \frac{e \vdash \mathsf{E}}{\{e\} \vdash \mathtt{P}\ \mathsf{E}} \qquad \frac{e_1 \vdash \mathsf{E}_1, \quad e_2 \vdash \mathsf{E}_2, \quad \mathsf{E}_1 \cong \mathsf{E}_2}{[e_1,e_2] \vdash \mathtt{P}\ \mathsf{E}_1}$$

$$\frac{s \vdash \mathsf{S}}{s^\infty \vdash \mathtt{seq}\ \mathsf{S}} \qquad \frac{s_1 \vdash \mathsf{S}_1, \quad s_2 \vdash \mathsf{S}_2, \quad \mathsf{S}_1 \cong \mathsf{S}_2}{s_1 \cup s_2 \vdash \mathsf{S}_1}$$

$$\frac{s_1 \vdash \mathsf{S}_1, \quad s_2 \vdash \mathsf{S}_2}{s_1 \mapsto s_2 \vdash \mathsf{S}_1 \nrightarrow \mathsf{S}_2} \qquad \frac{s_1 \vdash \mathsf{S}_1, \quad s_2 \vdash \mathsf{S}_2}{s_1 \times s_2 \vdash \mathsf{S}_1 * \mathsf{S}_2} \qquad \frac{s \vdash \mathsf{S}}{\wp(s) \vdash \mathtt{P}\ \mathsf{S}}$$

11.3 Partial orders and posets

$o \vdash o, \quad o \in \{\Rightarrow, \Leftrightarrow, \leq, \subseteq, =\}$ $\quad \dot{\subseteq} \vdash \subseteq$ $\quad \dfrac{o \vdash \mathsf{O}}{o^{-1} \vdash (\mathsf{O})^{-1}}$

$$\frac{o_1 \vdash \mathsf{O}_1, \quad o_2 \vdash \mathsf{O}_2}{o_1 \times o_2 \vdash \mathsf{O}_1 * \mathsf{O}_2} \qquad \frac{o \vdash \mathsf{O}}{\dot{o} \vdash (\dot{\mathsf{O}})} \qquad \frac{o \vdash \mathsf{O}}{\ddot{o} \vdash (\ddot{(\mathsf{O})})} \qquad \frac{s \vdash \mathsf{S}, \quad o \vdash \mathsf{o}'}{\langle s, o \rangle \vdash \mathsf{S} \circledast o'}$$

11.4 Galois connections

$$\frac{p \vdash \mathsf{P},}{\mathbb{I}[p] \vdash \mathsf{P} \rightleftharpoons \mathsf{P}} \qquad \frac{e \vdash \mathsf{E}, \quad p \vdash \mathsf{S} \circledast \mathsf{O}, \quad \mathsf{E} \Subset \mathsf{S}}{\top[p,e] \vdash \mathsf{S} \circledast \mathsf{O} \rightleftharpoons \mathsf{S} \circledast \mathsf{O}}$$

$$\frac{s \vdash \mathsf{S}, \quad b \vdash \mathsf{E}_b, \quad t \vdash \mathsf{E}_t, \quad \mathsf{E}_b \Subset \mathsf{S}, \quad \mathsf{E}_t \Subset \mathsf{S}}{\mathbb{I}[\langle s, o \rangle, b, t] \vdash (\mathtt{P}\ \mathsf{S} \circledast \subseteq) \rightleftharpoons (\mathtt{P}\ \mathsf{S} \circledast \subseteq)}$$

$$\frac{s_{\mathbb{L}} \vdash \mathsf{S}_{\mathbb{L}}, \quad s_{\mathcal{M}} \vdash \mathsf{S}_{\mathcal{M}}}{\curvearrowright[s_{\mathbb{L}}, s_{\mathcal{M}}] \vdash \mathtt{P}\ (\mathsf{S}_{\mathbb{L}} * \mathsf{S}_{\mathcal{M}}) \circledast \subseteq \rightleftharpoons \mathsf{S}_{\mathbb{L}} \nrightarrow \mathtt{P}\ \mathsf{S}_{\mathcal{M}} \circledast \dot{\subseteq}}$$

$$\frac{s \vdash \mathsf{S}}{\cup[s] \vdash \mathtt{P}\ (\mathtt{P}\ \mathsf{S}) \circledast \subseteq \rightleftharpoons \mathtt{P}\ \mathsf{S} \circledast \subseteq} \qquad \frac{s \vdash \mathsf{S}}{\neg[s] \vdash \mathtt{P}\ \mathsf{S} \circledast \subseteq \rightleftharpoons \mathtt{P}\ \mathsf{S} \circledast \subseteq^{-1}}$$

$$\frac{s \vdash \mathsf{S}}{\infty[s] \vdash \mathtt{P}\ (\mathtt{seq}\ \mathsf{S}) \circledast \subseteq \rightleftharpoons \mathtt{P}\ \mathsf{S} \circledast \subseteq}$$

$$\frac{s_1 \vdash \mathsf{S}_1, \quad s_2 \vdash \mathsf{S}_2}{\rightsquigarrow[s_1,s_2] \vdash \mathtt{P}(\mathsf{S}_1 * \mathsf{S}_2) \circledast \subseteq \rightleftharpoons \mathtt{P}\ \mathsf{S}_1 \nrightarrow \mathtt{P}\ \mathsf{S}_2 \circledast \dot{\subseteq}}$$

$$\frac{s_1 \vdash \mathsf{S}_1, \quad s_2 \vdash \mathsf{S}_2}{\mapsto[s_1,s_2] \vdash \mathtt{P}(\mathsf{S}_1 \nrightarrow \mathsf{S}_2) \circledast \subseteq \rightleftharpoons \mathtt{P}\ \mathsf{S}_1 \nrightarrow \mathtt{P}\ \mathsf{S}_2 \circledast \dot{\subseteq}}$$

$$\frac{s_1 \vdash \mathsf{S}_1, \quad s_2 \vdash \mathsf{S}_2}{\times[s_1,s_2] \vdash \mathtt{P}(\mathsf{S}_1 \nrightarrow \mathsf{S}_2) \circledast \subseteq \rightleftharpoons \mathsf{S}_1 \nrightarrow \mathtt{P}\ \mathsf{S}_2 \circledast \dot{\subseteq}}$$

$$\frac{g \vdash \mathsf{T}}{\mathsf{R}[g] \vdash \mathsf{T}} \qquad \frac{s \vdash \mathsf{S}, \quad g \vdash \mathsf{T}}{s \twoheadrightarrow g \vdash \mathsf{S} \nrightarrow \mathsf{T}}$$

$$\frac{g_1 \vdash \mathsf{P}_1 \rightleftharpoons \mathsf{P}_2, \quad g_2 \vdash \mathsf{P}_3 \rightleftharpoons \mathsf{P}_4, \quad \mathsf{P}_2 \cong \mathsf{P}_3}{g_1 \fatsemi g_2 \vdash \mathsf{P}_1 \rightleftharpoons \mathsf{P}_4}$$

$$\frac{g_1 \vdash \mathsf{S}_1 \circledast \mathsf{O}_1 \rightleftharpoons \mathsf{S}_2 \circledast \mathsf{O}_2, \quad g_2 \vdash \mathsf{S}_3 \circledast \mathsf{O}_3 \rightleftharpoons \mathsf{S}_4 \circledast \mathsf{O}_4}{g_1 \circledast g_2 \vdash \mathsf{S}_1 * \mathsf{S}_2 \circledast \mathsf{O}_1 \star \mathsf{O}_2 \rightleftharpoons \mathsf{S}_3 * \mathsf{S}_4 \circledast \mathsf{O}_3 \star \mathsf{O}_4}$$

$$\frac{g_1 \vdash \mathsf{S}_1 \circledast \mathsf{O}_1 \rightleftharpoons \mathsf{S}_2 \circledast \mathsf{O}_2, \quad g_2 \vdash \mathsf{S}_3 \circledast \mathsf{O}_3 \rightleftharpoons \mathsf{S}_4 \circledast \mathsf{O}_4}{g_1 \Longmapsto g_2 \vdash \mathsf{S}_1 \nrightarrow \mathsf{S}_3 \circledast \dot{\mathsf{O}}_3 \rightleftharpoons \mathsf{S}_2 \nrightarrow \mathsf{S}_4 \circledast \dot{\mathsf{O}}_4}$$

12. Principal type As usual for syntactic abstractions like types, there is in general no best abstraction of an arbitrary set[5]. For example there is no best regular expression to describe as precisely as possible a non-regular language. There is no best set type for an empty set. As shown in [2], the problem is solved for typing both by restricting types (*e.g.* no disjunctive type in ML which requires to tag the alternatives) and the programming language (*e.g.* alternatives of conditionals in ML cannot return objects of different types). The same is true for the *GC* calculus, in that there is a best abstraction for the properties of semantic objects generated by the language, although none exists in general.

[5] Patrick Cousot, Radhia Cousot: Formal Language, Grammar and Set-Constraint-Based Program Analysis by Abstract Interpretation. FPCA 1995: 170–181

12.1 Best abstraction of ur-element properties

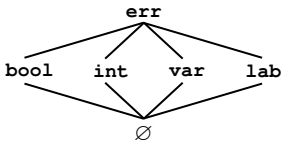
The ur-element properties are $\mathcal{P}^{\mathfrak{E}} \triangleq \wp(\{\mathcal{S}[\![e]\!] \mid e \in \mathbb{E}\})$. Define

- $\alpha^{\mathfrak{E}}(\emptyset) \triangleq \varnothing$[6], and otherwise
- $\alpha^{\mathfrak{E}}(P) \triangleq \mathtt{bool}$, when $\emptyset \neq P \subseteq \mathbb{B}$;
- $\alpha^{\mathfrak{E}}(P) \triangleq \mathtt{int}$, when $\emptyset \neq P \subseteq \mathbb{Z} \cup \{-\infty, \infty\}$;
- $\alpha^{\mathfrak{E}}(P) \triangleq \mathtt{var}$, when $\emptyset \neq P \subseteq \mathbb{X}$;
- $\alpha^{\mathfrak{E}}(P) \triangleq \mathtt{lab}$, when $\emptyset \neq P \subseteq \mathbb{L}$;
- $\alpha^{\mathfrak{E}}(P) \triangleq \mathtt{err}$, otherwise.

We have defined

- $\gamma^{\mathfrak{E}}(\varnothing) \triangleq \emptyset$;
- $\gamma^{\mathfrak{E}}(\mathtt{bool}) \triangleq \mathbb{B}$;
- $\gamma^{\mathfrak{E}}(\mathtt{int}) \triangleq \mathbb{Z} \cup \{-\infty, \infty\}$;
- $\gamma^{\mathfrak{E}}(\mathtt{var}) \triangleq \mathbb{X}$;
- $\gamma^{\mathfrak{E}}(\mathtt{lab}) \triangleq \mathbb{L}$;
- $\gamma^{\mathfrak{E}}(\mathtt{err}) \triangleq \{\mathcal{S}[\![e]\!] \mid e \in \mathbb{E}\}$.

We extend $\trianglelefteqslant$ to $\varnothing \trianglelefteqslant \varnothing \trianglelefteqslant \mathsf{E}$ for all $\mathsf{E} \in \mathfrak{E}$ so that we have the following lattice of element types.



It follows that

$$\begin{array}{lll} & \alpha^{\mathfrak{E}}(P) \trianglelefteqslant \mathsf{E} & \\ \Leftrightarrow & \alpha^{\mathfrak{E}}(P) = \mathsf{E} & \text{〈def. } \trianglelefteqslant \text{ on } \mathfrak{E}\text{〉} \\ \Leftrightarrow & P \subseteq \gamma^{\mathfrak{E}}(\mathsf{E}) & \text{〈def. } \alpha^{\mathfrak{E}} \text{ and } \gamma^{\mathfrak{E}}\text{〉} \end{array}$$

and so $\langle \mathcal{P}^{\mathfrak{E}}, \subseteq \rangle \underset{\alpha^{\mathfrak{E}}}{\overset{\gamma^{\mathfrak{E}}}{\leftrightarrows}} \langle \mathfrak{E} \cup \{\varnothing\}, \trianglelefteqslant \rangle$. Moreover $\gamma^{\mathfrak{E}}$ is injective so $\langle \mathcal{P}^{\mathfrak{E}}, \subseteq \rangle \underset{\alpha^{\mathfrak{E}}}{\overset{\gamma^{\mathfrak{E}}}{\leftrightarrows}} \langle \mathfrak{E} \cup \{\varnothing\}, \trianglelefteqslant \rangle$. We conclude that any ur-element semantic property $P \in \mathcal{P}^{\mathfrak{E}}$ now has a $\trianglelefteqslant$-best abstraction in $\mathfrak{E} \cup \{\varnothing\}$ which is their principal type $\alpha^{\mathfrak{E}}(P)$.

12.2 Best abstraction of set properties

The set semantic properties are $\mathcal{P}^{\mathfrak{S}} \triangleq \wp(\{\mathcal{S}[\![s]\!] \mid s \in \mathbb{S}\})$ *i.e.* the properties of those mathematical sets describable by the formal language $\mathbb{S}$ such as $\wp(\emptyset)$ which best abstraction is $\mathtt{P}\ \varnothing$, $\{\{\mathsf{true}\}, \mathbb{B}\}$ which best abstraction is $\mathtt{P}\ \mathbb{B}$ and $\{\{1\}, \mathbb{B}\}$ which best abstraction is $\mathtt{err}$.

If $P \in \mathcal{P}^{\mathfrak{S}}$ then either $P = \emptyset$ and $\alpha^{\mathfrak{S}}(\emptyset) \triangleq \varnothing$ else $P = \{\mathcal{S}[\![s_i]\!] \mid i \in \Delta \wedge \forall i \in \Delta : s_i \in \mathbb{S}\}$, $\Delta \neq \emptyset$. Then consider $\mathsf{S}_i \triangleq \mathcal{S}[\![s_i]\!]$, $i \in \Delta$. If $\forall i, j \in \Delta : \mathsf{S}_i \cong \mathsf{S}_j$ then $\alpha^{\mathfrak{S}}(P) = [\mathsf{S}_i]_{\cong}$ for some $i \in \Delta$ which choice is irrelevant. Else $\exists i, j \in \Delta : \mathsf{S}_i \ncong \mathsf{S}_j$, in which case $\alpha^{\mathfrak{S}}(P) = \mathtt{err}$.

Observe that $\alpha^{\mathfrak{S}}$ is increasing, sound since $\alpha^{\mathfrak{S}}(P) \trianglelefteqslant \mathsf{S} \Rightarrow P \subseteq \gamma^{\mathfrak{S}}(\mathsf{S})$, and the best abstraction since $P \subseteq \gamma^{\mathfrak{S}}(\mathsf{S}) \Rightarrow \alpha^{\mathfrak{S}}(P) \trianglelefteqslant \mathsf{S}$.

It follows that $\langle \mathcal{P}^{\mathfrak{S}}, \subseteq \rangle \underset{\alpha^{\mathfrak{S}}}{\overset{\gamma^{\mathfrak{S}}}{\leftrightarrows}} \langle (\mathfrak{S} \cup \{\varnothing\})_{/\cong}, \trianglelefteqslant \rangle$ where $\trianglelefteqslant$ is naturally expended to equivalence classes since $\alpha^{\mathfrak{S}}$ is onto/$\gamma^{\mathfrak{S}}$ injective.

[6] By analogy with a programming language, $\emptyset$ would be the false invariant so that $\varnothing$ would be the type of unreachable code.

12.3 Best abstraction of partial order and poset properties

The partial order properties are $\mathcal{P}^{\mathfrak{O}} \triangleq \wp(\{\mathcal{S}[\![o]\!] \mid o \in \mathbb{O}\})$. Define $\alpha^{\mathfrak{O}}(P) \triangleq \big(P = \emptyset \mathbin{?} \varnothing \parallel \exists \mathsf{O} \in \mathfrak{O} : \forall o \in \mathbb{O} : \mathcal{S}[\![o]\!] \in P \Rightarrow \mathcal{O}[\![o]\!] \cong \mathsf{O} \mathbin{?} [\mathsf{O}]_{\cong} \mathbin{:} \mathtt{err}\big)$. Then $\langle \mathcal{P}^{\mathfrak{O}}, \subseteq\rangle \underset{\alpha^{\mathfrak{O}}}{\overset{\gamma^{\mathfrak{O}}}{\leftrightarrows}} \langle (\mathfrak{O} \cup \{\varnothing\})_{/\cong}, \trianglelefteq\rangle$.

Similarly, poset properties are $\mathcal{P}^{\mathfrak{P}} \triangleq \wp(\{\mathcal{S}[\![p]\!] \mid p \in \mathbb{P}\})$. Define $\alpha^{\mathfrak{P}}(P) \triangleq \big(P = \emptyset \mathbin{?} \varnothing \parallel \exists \mathsf{P} : \forall \mathcal{S}[\![o]\!] \in P : \mathcal{P}[\![p]\!] \cong \mathsf{P} \mathbin{?} [\mathsf{P}]_{\cong} \mathbin{:} \mathtt{err}\big)$. Then $\langle \mathcal{P}^{\mathfrak{P}}, \subseteq\rangle \underset{\alpha^{\mathfrak{P}}}{\overset{\gamma^{\mathfrak{P}}}{\leftrightarrows}} \langle (\mathfrak{P} \cup \{\varnothing\})_{/\cong}, \trianglelefteq\rangle$.

12.4 Best abstraction of *GC* properties

Finally, the semantic properties $P = \{\mathcal{S}[\![g_i]\!] \mid i \in \Delta\}$ of *GC* expressions, the principal type is $\alpha^{\mathfrak{T}}(\emptyset) \triangleq \varnothing$, $\alpha^{\mathfrak{T}}(P) \triangleq [\mathsf{T}]_{\cong}$ when $\forall i \in \Delta \neq \emptyset : \mathcal{T}[\![g_i]\!] \cong \mathsf{T}$ else $\alpha^{\mathfrak{T}}(P) \triangleq \mathtt{err}$ so $\langle \wp(\{\mathcal{S}[\![g]\!] \mid g \in \mathbb{G}\}), \subseteq\rangle \underset{\alpha^{\mathfrak{T}}}{\overset{\gamma^{\mathfrak{T}}}{\leftrightarrows}} \langle (\mathfrak{T} \cup \{\varnothing\})_{/\cong}, \trianglelefteq\rangle$ since $\alpha^{\mathfrak{T}}$ is onto.

13. Types of types To add one more level of abstraction, we can consider types of types.

Types $\mathcal{T} \triangleq \{\mathfrak{E}, \mathfrak{S}, \mathfrak{O}, \mathfrak{P}, \mathfrak{T}\}$ have properties which are sets of types *i.e.* $\mathcal{P} \triangleq \wp(\mathfrak{E} \cup \mathfrak{S} \cup \mathfrak{O} \cup \mathfrak{P} \cup \mathfrak{T}) = \wp(\bigcup \mathcal{T})$ abstracted by types of types $\overline{\mathfrak{T}} ::= \overline{\varnothing} \mid \overline{\mathfrak{E}} \mid \overline{\mathfrak{S}} \mid \overline{\mathfrak{O}} \mid \overline{\mathfrak{P}} \mid \overline{\mathfrak{T}} \mid \overline{\mathtt{err}}$.

The abstraction is $\alpha^{\overline{\mathfrak{T}}}(P) \triangleq \big(P = \emptyset \mathbin{?} \overline{\varnothing} \parallel P \subseteq \mathsf{T}, \mathsf{T} \in \{\mathfrak{E}, \mathfrak{S}, \mathfrak{O}, \mathfrak{P}, \mathfrak{T}\} \mathbin{?} \overline{\mathsf{T}} \mathbin{:} \overline{\mathtt{err}}\big)$.

The concretization is $\gamma^{\overline{\mathfrak{T}}}(\overline{\mathsf{T}}) \triangleq \big(\overline{\mathsf{T}} = \overline{\varnothing} \mathbin{?} \emptyset \parallel \overline{\mathsf{T}} \in \{\overline{\mathfrak{E}}, \overline{\mathfrak{S}}, \overline{\mathfrak{O}}, \overline{\mathfrak{P}}, \overline{\mathfrak{T}}\} \mathbin{?} \mathsf{T} \mathbin{:} \mathfrak{E} \cup \mathfrak{S} \cup \mathfrak{O} \cup \mathfrak{P} \cup \mathfrak{T}\big)$ so $\gamma^{\overline{\mathfrak{T}}}(\overline{\mathtt{err}}) = \mathfrak{E} \cup \mathfrak{S} \cup \mathfrak{O} \cup \mathfrak{P} \cup \mathfrak{T}$.

Defining $\overline{\mathsf{T}} \overline{\trianglelefteq} \overline{\mathsf{T}}' \triangleq \gamma^{\overline{\mathfrak{T}}}(\overline{\mathsf{T}}) \subseteq \gamma^{\overline{\mathfrak{T}}}(\overline{\mathsf{T}}')$, we get the flat partial order $\overline{\varnothing} \overline{\trianglelefteq} \overline{\varnothing} \overline{\trianglelefteq} \overline{\mathsf{T}} \overline{\trianglelefteq} \overline{\mathsf{T}} \overline{\trianglelefteq} \overline{\mathtt{err}} \overline{\trianglelefteq} \overline{\mathtt{err}}$ for all $\mathsf{T} \in \{\mathfrak{E}, \mathfrak{S}, \mathfrak{O}, \mathfrak{P}, \mathfrak{T}\}$.

The abstraction is $\langle \mathcal{P}, \subseteq\rangle \underset{\alpha^{\overline{\mathfrak{T}}}}{\overset{\gamma^{\overline{\mathfrak{T}}}}{\leftrightarrows}} \langle \overline{\mathfrak{T}}, \overline{\trianglelefteq}\rangle$ where $\alpha^{\overline{\mathfrak{T}}}(\mathcal{P})$ is the principal type of type of the type property $\mathcal{P}$.

Define the type inference $\overline{\mathcal{T}}[\![\mathsf{T}]\!] \triangleq \alpha^{\overline{\mathfrak{T}}}(\{\mathsf{T}\})$ for all $\mathsf{T} \in \bigcup \mathcal{T}$. The typing rules for the meta-language that we have used to define $\mathcal{T}$ include traditional rules $\mathcal{T}^{\mathbb{B}}$ for Boolean expressions as well as

- $\overline{\mathcal{T}}[\![\mathtt{bool}]\!] = \overline{\mathcal{T}}[\![\mathtt{int}]\!] = \overline{\mathcal{T}}[\![\mathtt{var}]\!] = \overline{\mathcal{T}}[\![\mathtt{lab}]\!] \triangleq \overline{\mathfrak{E}}$
- $\overline{\mathcal{T}}[\![\mathtt{P}\ X]\!] \triangleq \big(\overline{\mathcal{T}}[\![X]\!] = \mathfrak{E} \vee \overline{\mathcal{T}}[\![X]\!] = \mathfrak{S} \mathbin{?} \mathfrak{S} \mathbin{:} \overline{\mathtt{err}}\big)$
- $\overline{\mathcal{T}}[\![\subseteq]\!] \triangleq \overline{\mathfrak{O}}$
- $\overline{\mathcal{T}}[\![(X)^{-1}]\!] \triangleq \big(\overline{\mathcal{T}}[\![X]\!] = \overline{\mathfrak{O}} \mathbin{?} \overline{\mathfrak{O}} \mathbin{:} \overline{\mathtt{err}}\big)$
- $\overline{\mathcal{T}}[\![X_1 \circledast X_2]\!] \triangleq \big(\overline{\mathcal{T}}[\![X_1]\!] = \overline{\mathfrak{S}} \wedge \overline{\mathcal{T}}[\![X_2]\!] = \overline{\mathfrak{O}} \mathbin{?} \overline{\mathfrak{P}} \mathbin{:} \overline{\mathtt{err}}\big)$
- $\overline{\mathcal{T}}[\![X_1 \leftrightharpoons X_2]\!] \triangleq \big(\overline{\mathcal{T}}[\![X_1]\!] = \overline{\mathcal{T}}[\![X_2]\!] = \overline{\mathfrak{P}} \mathbin{?} \overline{\mathfrak{T}} \mathbin{:} \overline{\mathtt{err}}\big)$
- $\overline{\mathcal{T}}[\![X_1 \mathbin{;} X_2]\!] \triangleq \big(\overline{\mathcal{T}}[\![X_1]\!] = \overline{\mathcal{T}}[\![X_2]\!] = \overline{\mathfrak{T}} \mathbin{?} \overline{\mathfrak{T}} \mathbin{:} \overline{\mathtt{err}}\big)$
- $\overline{\mathcal{T}}[\![(X_1 \ ? \ X_2 \ : \ X_3)]\!] \triangleq \big(\mathcal{T}^{\mathbb{B}}[\![X_1]\!] = \mathtt{bool} \wedge (\overline{\mathcal{T}}[\![X_2]\!] = \overline{\mathcal{T}}[\![X_3]\!] \vee X_3 = \mathtt{err}) \mathbin{?} \overline{\mathcal{T}}[\![X_2]\!] \mathbin{:} \overline{\mathtt{err}}\big)$
- *etc.*

Well-typing of types requires that $\overline{\mathcal{T}}[\![\mathsf{T}]\!] \neq \overline{\mathtt{err}}$. In particular, for all $g \in \mathbb{G}$, $\overline{\mathcal{T}}[\![\mathcal{T}[\![g]\!]]\!] \neq \overline{\mathtt{err}}$ so that $\mathcal{T}[\![g]\!]$ is well-defined that is "well-types types cannot go wrong".

Notice that the typing of types rules correspond to the routine informal verification of mathematical texts by checking that operators take their arguments and return their results in their domains of definition.

14. Abstraction of induction by widening/narrowing In static analyzers [1, 12, 14] *GCs* specify abstract domains modules and Galois connectors their combinations by functors. For scalability and precision, rapid convergence acceleration of infinite fixpoint computations by widening/narrowing abstracting induction and/or their duals for co-induction [3–5] is more precise than finite abstractions [8].

This can be illustrated for recursive definitions of sets such as $\mathbf{lfp}^{\subseteq}\ \boldsymbol{\lambda}\, S \bullet \{\emptyset\} \cup \wp(S) = \{\emptyset\} \cup \wp(\{\emptyset\}) \cup \wp(\{\emptyset\} \cup \wp(\{\emptyset\})) \cup \ldots\} = \{\emptyset\} \cup \{\emptyset, \{\emptyset\}\} \cup \{\emptyset, \{\emptyset\}, \{\{\emptyset\}\}, \{\emptyset, \{\emptyset\}\}\} \cup \ldots\} = \{\emptyset, \{\emptyset\}, \{\{\emptyset\}\}, \{\emptyset, \{\emptyset\}\}, \ldots\}$ which can be typed by introducing recursive/fixpoint type definitions such as $\mathbf{lfp}^{\trianglelefteq}\ \boldsymbol{\lambda}\, \mathsf{S} \bullet \mathtt{P}\ \varnothing \cup \mathtt{P}\ \mathsf{S}$.

This requires a widening, such as the trivial one considered in ML typing, hidden in the type inference rules, and requiring that all actual arguments in recursive calls have the same type [2] .

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