

# Abstracting Induction by Extrapolation and Interpolation

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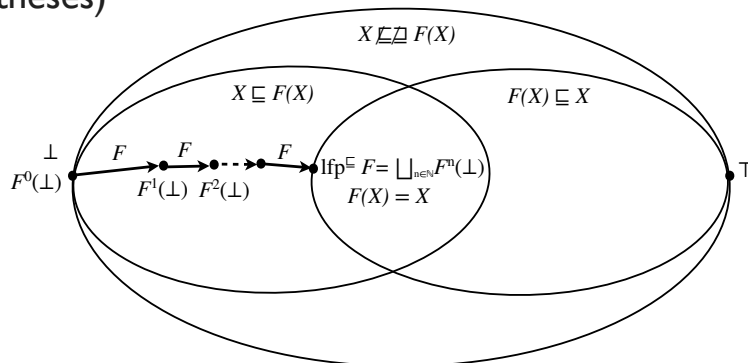
## Abstract Interpreters

- **Transitional abstract interpreters:** proceed by induction on program steps
- **Structural abstract interpreters:** proceed by induction on the program syntax
- **Common main problem:** over/under-approximate fixpoints in non-Noetherian<sup>(\*)</sup> abstract domains<sup>(\*\*)</sup>

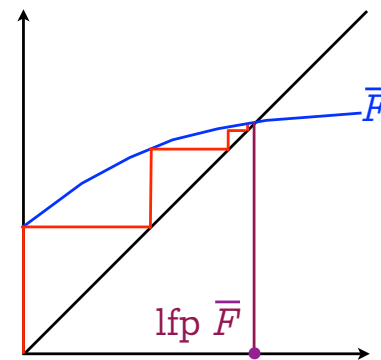
(\*) Iterative fixpoint computations may not converge in finitely many steps  
(\*\*) Or convergence may be guaranteed but to slow.

## Fixpoints

- Poset (or pre-order)  $\langle D, \sqsubseteq, \perp, \top \rangle$
- Transformer:  $F \in D \mapsto D$
- Least fixpoint:  $\text{lfp}^{\sqsubseteq} F = \bigsqcup_{n \in \mathbb{N}} F^n(\perp)$  (under appropriate hypotheses)

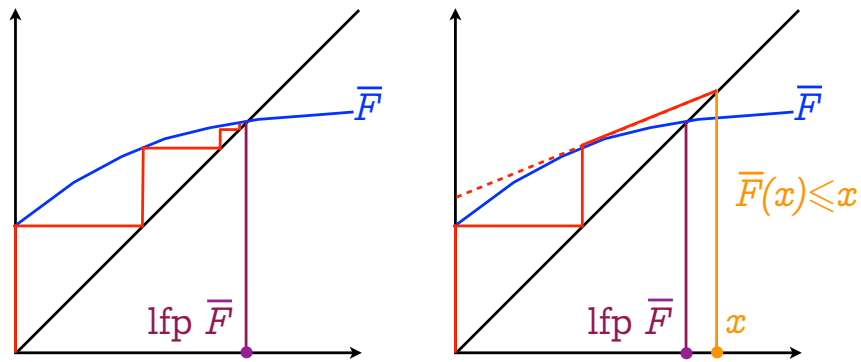


## Convergence acceleration with widening



Infinite iteration

# Convergence acceleration with widening



Infinite iteration

Accelerated iteration with widening  
(e.g. with a widening based on the derivative as in Newton-Raphson method<sup>(\*)</sup>)

<sup>(\*)</sup> Javier Esparza, Stefan Kiefer, Michael Luttenberger: Newtonian program analysis. J. ACM 57(6): 33 (2010)

# Extrapolation by Widening

- $X^0 = \perp$  (increasing iterates with widening)
- $X^{n+1} = X^n \nabla F(X^n)$  when  $F(X^n) \not\subseteq X^n$
- $X^{n+1} = X^n$  when  $F(X^n) \subseteq X^n$
- Widening  $\nabla$ :
  - $Y \subseteq X \nabla Y$  (extrapolation)
  - Enforces convergence of increasing iterates with widening (to a limit  $X^\ell$ )

# The oldest widenings

- Primitive widening [1,2]

```

(x  $\bar{\vee}$  y) = max x < y1, y < y2 dans
  [ 0, ? => y ;
    ? , 0 => x ;
    [a1, a2] , [b1, b2] =>
      [if a1 < a2 then -∞ else a1 fi ;
       if b1 < b2 then +∞ else b1 fi] ;
  fincas ;
    
```

```

[a1, b1]  $\bar{\vee}$  [a2, b2] =
  [if a2 < a1 then -∞ else a1 fi,
   if b2 > b1 then +∞ else b1 fi]
    
```

- Widening with thresholds [3]

$$\forall x \in \bar{L}_2, \perp \nabla_2(j) x = x \nabla_2(j) \perp = x$$

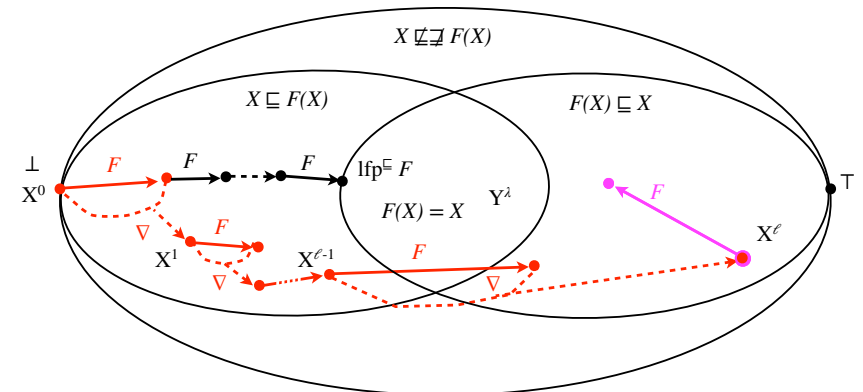
$$[l_1, u_1] \nabla_2(j) [l_2, u_2]$$

$$= [\text{if } 0 \leq l_2 < l_1 \text{ then } 0 \text{ elsif } l_2 < l_1 \text{ then } -b - 1 \text{ else } l_1 \text{ fi},$$

$$\text{if } u_1 < u_2 \leq 0 \text{ then } 0 \text{ elsif } u_1 < u_2 \text{ then } b \text{ else } u_1 \text{ fi}]$$

[1] Patrick Cousot, Radhia Cousot: Vérification statique de la cohérence dynamique des programmes. Rapport du contrat IRIA-SESORI No 75-032, 23 septembre 1975.  
[2] Patrick Cousot, Radhia Cousot: Abstract Interpretation: A Unified Lattice Model for Static Analysis of Programs by Construction or Approximation of Fixpoints. POPL 1977: 238-252  
[3] Patrick Cousot, Semantic foundations of program analysis, Ch. 10 of Program flow analysis: theory and practice, N. Jones & S. Muchnick (eds), Prentice Hall, 1981.

# Extrapolation with widening



# Interpolation with narrowing

- $Y^0 = X^\ell$  (decreasing iterates with narrowing)

$$Y^{n+1} = Y^n \Delta F(Y^n) \quad \text{when } F(Y^n) \subseteq Y^n$$

$$Y^{n+1} = Y^n \quad \text{when } F(Y^n) = Y^n$$

- **Narrowing  $\Delta$ :**

- $Y \subseteq X \implies Y \subseteq X \Delta Y \subseteq X$  (interpolation)

- Enforces **convergence** of decreasing iterates with narrowing (to a limit  $Y^\lambda$ )

# The oldest narrowing

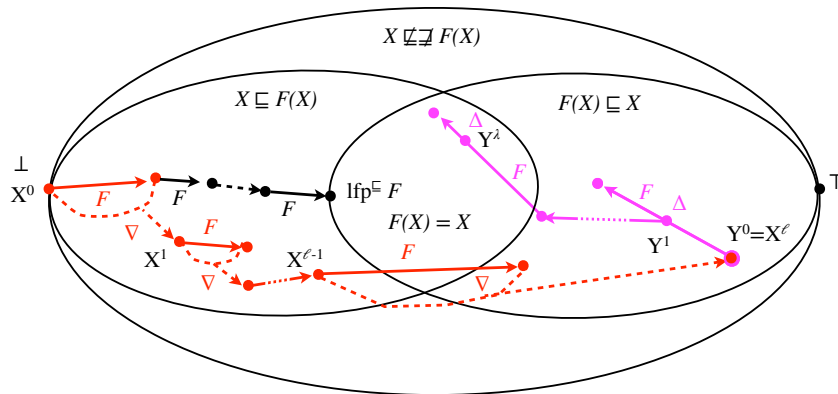
- [2]

$$[a_1, b_1] \bar{\Delta} [a_2, b_2] =$$

$$[\text{if } a_1 = -\infty \text{ then } a_2 \text{ else } \text{MIN}(a_1, a_2),$$

$$\text{if } b_1 = +\infty \text{ then } b_2 \text{ else } \text{MAX}(b_1, b_2)]$$

# Interpolation with narrowing



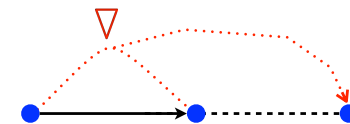
Could stop when  $F(X) \not\subseteq X \wedge F(F(X)) \subseteq F(X)$  but not the current practice.

# Duality

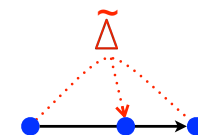
	Convergence above the limit	Convergence below the limit
Increasing iteration	Widening $\nabla$	Dual-narrowing $\bar{\Delta}$
Decreasing iteration	Narrowing $\Delta$	Dual widening $\bar{\nabla}$

Extrapolators ( $\nabla, \bar{\nabla}$ ) and interpolators ( $\Delta, \bar{\Delta}$ )

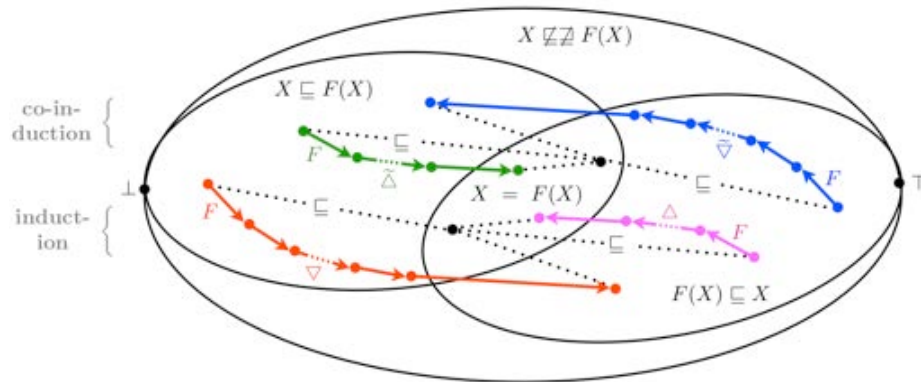
- **Extrapolators:**



- **Interpolators:**



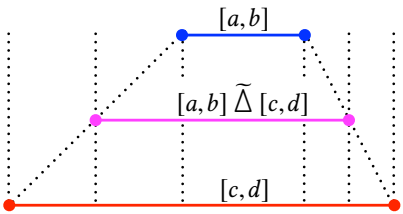
# Extrapolators, Interpolators, and Duals



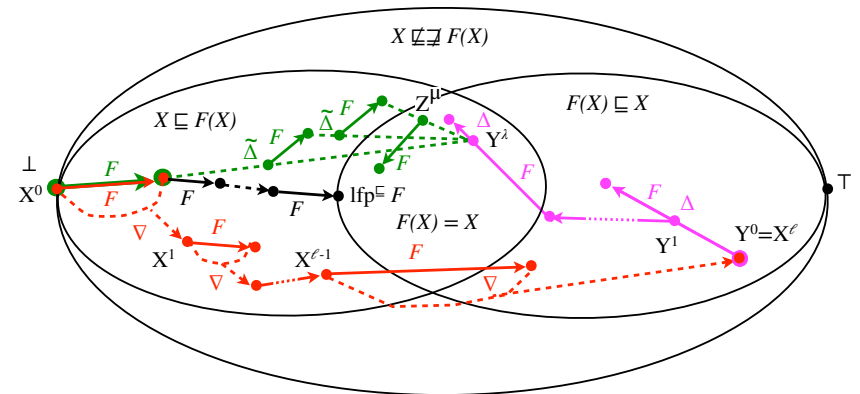
# Interpolation with dual narrowing

- $Z^0 = \perp$  (increasing iterates with dual-narrowing)
- $Z^{n+1} = F(Z^n) \tilde{\Delta} Y^\lambda$  when  $F(Z^n) \not\subseteq Z^n$
- $Z^{n+1} = Z^n$  when  $F(Z^n) \subseteq Z^n$
- Dual-narrowing  $\tilde{\Delta}$ :
  - $X \subseteq Y \Rightarrow X \subseteq X \tilde{\Delta} Y \subseteq Y$  (interpolation)
  - Enforces convergence of increasing iterates with dual-narrowing

# Example of dual-narrowing

- 
- $[a, b] \tilde{\Delta} [c, d] \triangleq [(c = -\infty ? a : [(a+c)/2]), (d = \infty ? b : [(b+d)/2])]$
- The first method we tried in the late 70's with Radhia
  - Slow
  - Does not easily generalize (e.g. to polyhedra)

# Interpolation with dual-narrowing

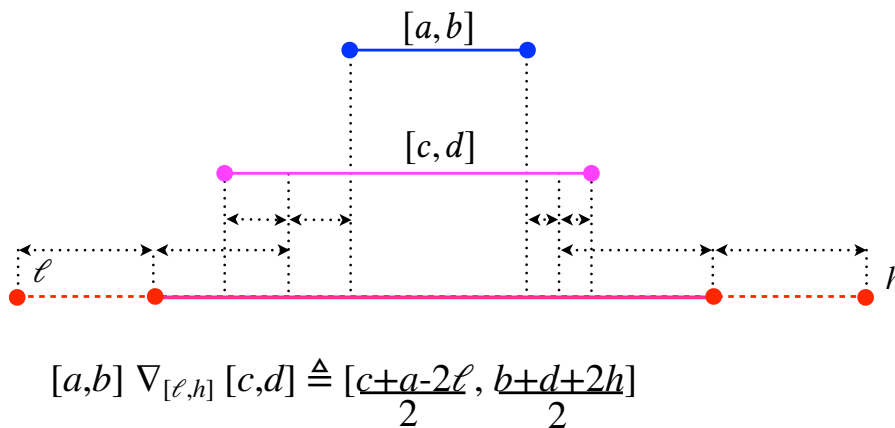


## Relationship between narrowing and dual-narrowing

- $\tilde{\Delta} = \Delta^{-1}$
- $Y \sqsubseteq X \Rightarrow Y \sqsubseteq X \Delta Y \sqsubseteq X$  (narrowing)
- $Y \sqsubseteq X \Rightarrow Y \sqsubseteq Y \tilde{\Delta} X \sqsubseteq X$  (dual-narrowing)
- Example: Craig interpolation
- Why not use a bounded widening (bounded by B)?
  - $F(X) \sqsubseteq B \Rightarrow F(X) \sqsubseteq F(X) \tilde{\Delta} B \sqsubseteq B$  (dual-narrowing)
  - $X \sqsubseteq F(X) \sqsubseteq B \Rightarrow F(X) \sqsubseteq X \nabla_B F(X) \sqsubseteq B$  (bounded widening)

## Example of widenings (cont'd)

- Bounded widening (in  $[\ell, h]$ ):



More in the paper...

Widenings

## Widenings are not increasing

- A **well-known** fact

$$[1,1] \subseteq [1,2] \text{ but } [1,1] \nabla [1,2] = [1, \phi] \subseteq [1,2] \nabla [1,2] = [1,2]$$

- A widening cannot both:
  - Be **increasing** in its first parameter
  - Enforce **termination** of the iterates
  - Avoid **useless over-approximations** as soon as a solution is found<sup>(\*)</sup>

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<sup>(\*)</sup> A counter-example is  $x \nabla y = \top$

## Soundness

## Soundness

- In the paper, the fixpoint approximation soundness theorems are expressed with **minimalist hypotheses**:
  - No need for complete lattices, complete partial orders (CPO's):
    - The concrete domain is a poset
    - The abstract domain is a pre-order
    - The concretization is defined for the abstract iterates only.

## Soundness (cont'd)

- No need for increasingness/monotony hypotheses for fixpoint theorems (Tarski, Kleene, etc)
  - The concrete transformer is increasing and the limit of the iterations does exist in the concrete domain
  - No hypotheses on the abstract transformer (no need for fixpoints in the abstract)
  - Soundness hypotheses on the extrapolators/interpolators with respect to the concrete
- In addition, termination hypotheses on the extrapolators/interpolators ensure convergence in finitely many steps

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- The concrete transformer is increasing and the limit of the iterations does exist in the concrete domain
- No hypotheses on the abstract transformer (no need for fixpoints in the abstract)
- Soundness hypotheses on the extrapolators/interpolators with respect to the concrete

## Examples of interpolators

## Craig interpolation

- Craig interpolation:

Given  $P \implies Q$  find  $I$  such that  $P \implies I \implies Q$  with  $\text{var}(I) \subseteq \text{var}(P) \cap \text{var}(Q)$

is a **dual narrowing** (already observed by Vijay D'Silva and Leopold Haller as an inversed narrowing)

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- This evidence looked very controversial to some reviewers
- The generalization of an idea does not diminish in any way the merits and originality of this idea

## Conclusion

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- Abstract interpretation in infinite domains is traditionally by **iteration with widening/narrowing**.
- We have shown how to use **iteration with dual-narrowing** (alone or after widening/narrowing).
- These ideas of the 70's **generalize Craig interpolation from logic to arbitrary abstract domains**.

# The End, Thank You