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Challenges in control/ command software analysis

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Motivation
Ποτινατιστι
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Computer controlled systems







Australian Transport "The Safety Bureau (ATSB) found that the main probable cause of this incident was a *latent* software error which allowed the ADIRU to use data from a failed accelerometer"

ADIRU = Air Data Inertial Reference Unit (provides air speed, altitude & position) 4

- The initial effects of the fault were:
 - false stall and overspeed warnings
 - loss of attitude information on the Captain's primary flight display
 - several Electronic Centralised Aircraft Monitoring (ECAM) system warnings
- About two minutes later, ADIRU #1, which was providing data to the captain's primary flight display, provided very high (and false) indications for the aircraft's angle of attack, leading to:
 - the flight control computers commanding a nose-down aircraft movement, which resulted in the aircraft pitching down to a maximum of about 8.5 degrees,
 - the triggering of a Flight Control Primary Computer pitch fault.
- On 15 January 2009 the EASA issued an Emergency Airworthiness Directive to address the above A330 and A340 Northrop-Grumman ADIRU problem of incorrectly responding to a defective inertial reference.₅

• A memo leaked from Airbus on the Flight 447 (from Rio de Janeiro, Brazil, to Paris, that crashed into the Atlantic Ocean on I June 2009) suggests that there was no evidence that Honeywell manufactured ADIRU malfunction was similar to the failure of the Northrop Grumman manufactured ADIRU in Qantas flight incidents



An Air Data Inertial Reference Unit (ADIRU) is a key component of the integrated Air Data Inertial Reference System (ADIRS), that supplies air data (airspeed, angle of attack and altitude) and inertial reference (position and attitude) information to the pilots' Electronic Flight Instrument System displays as well as other systems on the aircraft such as the engines, autopilot, flight control and landing gear systems Brief overview of one approach to the design of an aircraft digital fly by wire control system



Mechanical flight-control systems

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Simulation of the controlled system e.g. in Simulink™





Design of a fault-tolerant control system (I)

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- Separate flight computers with different software and inputs
- The results of these command and monitoring units are checked by a comparator
- In case of desagreement control switches to another pair of computers
- In total 10 primary and secondary computers

Traverse P., Lacaze I., Souyris J., «Airbus Fly-by-Wire: A Total Approach to Dependability », Proceedings 18th IFIP World Computer Congress, Building the Information Society, Toulouse, France, 22-27 août 2004, pp. 191-212.

(I) Airbus



Design of a fault-tolerant control system (II)

LANE 1

POWER

PHOCESS AMD 750

ARINC 829 BITERFACE LANE 2

POWER

ARINC 629 LANE 3

POWER

ARINC 529

- Three different primary flight computers
- The computers have the same input signals and are all active
- The outputs are connected to a voter that compares signals
- The correct signal is chosen by a majority vote

Yeh Y.C., « Triple-Triple Redundant 777 Primary Flight Computers », Proceedings IEEE Aerospace Applications Conference, Aspen, CO, USA, 3-10 février 1996, pp. 293-307.

(II) Boeing

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Generation of C code

- Automatic C code generation
- The primitives are programmed by hand



Automatic compilation and linkage in machine code



Simulation and testing

- From classical software testing and inspection techniques (e.g. Rational Suite DevelopmentStudio⁽¹⁾)
- To full-scale similation benches for system integration:
- To the iron bird:
- To flight tests:

(1) www-01.ibm.com/software/awdtools/suite/dstudio/unix/
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Examples of model-checkers

 SCADE Design Verifier (<u>www.esterel-technologies.com/products/</u> <u>scade-suite/design-verifier</u>): check the logic of flight-control primitives

Examples of provers

• Caveat (<u>http://www-list.cea.fr/labos/gb/LSL/caveat/</u> <u>index.html</u>): Hoare logic prover with user-provided invariants, used to prove the Autotest primitive functions

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Examples of static analyzers

- AiT WCET Analyzers (<u>www.absint.de/ait/</u>): compute tight bounds for the worst-case execution time (WCET) of tasks in synchronous real-time systems.
- Astrée (www.astree.ens.fr, www.absint.de/astree/): proof of absence of runtime error in the primary flight control software of the Airbus A340/600 & A380 fly-by-wire system.
- Fluctuat (www-list.cea.fr/labos/gb/LSL/fluctuat/ flopsoft.html): contribution of individual program operations to the global error on float variables, used on DSPs (.
- StackAnalyzer (<u>www.absint.de/stackanalyzer/</u>): worstcase stack usage of the tasks in all applications.

Short term challenges for formal methods

From sequential/synchronous to parallel/ distributed



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Flight warning system



Long-term challenges for formal methods

Future challenges

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A software bug discovered during flight tests is extremely costly, but easy to correct during the design phase

- Present day: process validation + a posteriori software validation
- Future: early definition and validation of [software] systems + final product validation

Models and specifications

- Absence of formal specifications: model-based development in Simulink[™], Scade[™], etc. means that the model is the only available formal specification ⇒ no formal specification of the model itself
- Implicit specifications: can be embedded in verification tools (and documented e.g. in simple data files), examples:
- Integrity: no computer can send wrong information to actuators
- Availability: in a fault-tolerant system, no component failure can result in the unavailability of a service
- Electromagnetic radiations: memory glitches cannot make programs to fail
- etc

• Explicit formal specifications ultimately needed: in which form?



Environment specifications

• Precise verification:



Hybrid models of discrete control programs and continuous physical systems



- Recent improvements on bounded time analysis of continuous and hybrid systems⁽¹⁾
- Not so much in duration (``unbounded time")
- \Rightarrow should scale from seconds to minutes and more.

(I) Colas Le Guernic. Reachability Analysis of Hybrid Systems with Linear Continuous Dynamics. PhD thesis, University Joseph Fourier of Grenoble, October 29, 2009.

Conclusion

- Formal methods should be used:
- All along the whole design chain
- On functional properties (specification?)
- On a long range of time (s \rightarrow mn \rightarrow h: scalability)
- For parallel/distributed systems (semantics?)
- For end-users (non-intrusion)

Thank you for your attention

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Conclusion

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- On functional properties (specification?)
- On a long range of time (s \rightarrow mn \rightarrow h: scalability)
- For parallel/distributed systems (semantics?)
- For end-users (non-intrusive methods)
 - Robust software has simple abstractions
 - Long proofs imply bugs
 Paraphrasing John Doyle

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