| Présentation de l'équipe-projet ABSTRACTION <br> Patrick Cousot <br> Rocquencourt, I3 Janvier 2011 $\qquad$ | Abstract Domains | New numerical abstract domains <br> - Expressivity (non-convex relational domains) <br> - Efficiency (sound implementations with floats not rationals) | Array content analysis <br> - Existing abstract domains are imprecise or do not scale up <br> - We have developed efficient, expressive \& scalable abstract domain functors based on the idea of symbolic segmentation with dynamic bounds combined with decision trees $\qquad$ $\qquad$ |
| :---: | :---: | :---: | :---: |
|  | Abstract domain <br> - Algebraic structure (implemented as an analyzer module) $\langle A, \check{,}, \perp, \mathrm{~T}, \sqcup, \sqcap, \nabla, \Delta, \overline{\bar{f}}, \overline{\mathrm{~b}}, \overline{\mathrm{p}}, \ldots\rangle$ <br> such that <br> $\bar{P}, \bar{Q}, \ldots \in$ $\begin{aligned} & \because \in A \\ & \sqsubseteq \in A \times A \rightarrow \mathscr{E} \end{aligned}$ $\begin{aligned} & \quad \perp, \mathrm{T} \in A \\ & \sqcup, \Pi, \nabla, \Delta \in A \times A \rightarrow A \end{aligned}$ <br> infimum, supremum <br> $\in(x \times \mathbb{E}(x, f, p)) \rightarrow A \rightarrow A \quad$ abstract forward assignment transformer $\begin{array}{ll}\mathrm{b} \in(\mathrm{x} \times \mathbb{E}(\mathrm{x}, \mathrm{f}, \mathbb{p})) \rightarrow A \rightarrow A & \text { abstract backward assignment transformer } \\ \overline{\mathrm{p}} \in \mathbb{C}(\mathrm{x}, \mathbb{f}, \mathbb{p}) \rightarrow A \rightarrow A & \text { abstract condition transformer }\end{array}$ $\overline{\mathrm{p}} \in \mathbb{C}(\mathbb{x}, \mathbb{f}, \mathbb{p}) \rightarrow A \rightarrow A$ <br> abstract condition transformer | APRON Library <br> - Freely available (LGPL) library of numerical domains developed with INRIA-RA, including a web-based sample analyzer for demonstration, teaching and prototyping (Interproc) <br> - Common, Al-based API for numerical abstract domains with reference implementation of classic domains (intervals, polyhedra, octagons, ...) <br> - Easy to prototype new analyses or domains (e.g., interval polyhedra) with support for <br> - integer, rational, machine-integer and float data-types <br> - linear, non-linear and float expressions <br> - C, C++, Java, OCaml | Abstract Domains <br> (II) Symbolic domains <br> (b) Shape analysis |
| Abstract interpretation: from theory to practice <br> - Develop formal semantics of systems Theory (description of the possible evolution of discrete/continuous/hybrid systems over time) <br> - Formalize the specification and inference of properties of such semantics <br> - Develop precise, efficient, and scalable abstractions of system behavior properties <br> - Apply (maybe with approximation) to the inference of execution properties of systems <br> - Develop automatic static analysers <br> - Industrialize these static analysers | Abstract Domains <br> (I) Numerical Domains | Abstract Domains <br> (II) Symbolic domains | Shape abstract domains <br> - Challenges: <br> - dynamic structures with destructive updates <br> - complex structural invariants (red-black trees, callstack, ...) <br> - We develop relational and expressive parametric abstract domains based on inductive properties and disjointness of regions <br> Ongoing work: standardization of abstract domain interfaces to build composite abstractions (shapeshape and shape-numerical) $\qquad$ |
| Foundations | Problems in numerical abstract domains <br> - Traditional linear/polyedral abstract domains are implemented with the double description method, using rationals for soundness <br> - Inequalities: $P=\{x \mid A x=b, C x \geq d\}$ <br> - Generators: $P=\left\{x \mid x=L \lambda+R \mu+V \nu, \mu, \nu \geq 0, \sum \nu=1\right\}$ <br> - Results: sound algorithms using floats (with the inequalities representation only) and generalizations <br> - Current research: <br> - For efficiency, use the double description with floats while remaining sound <br> - For precision, non-linear and sound abstract domains | Abstract Domains <br> (II) Symbolic domains <br> (a) Array content analysis | Abstract Domains <br> (II) Symbolic domains <br> (c) Combining algebraic and logical abstract domains |

Logical \& algebraic abstraction

- Algebraic abstractions: in inference tools, based on the
iterated reduced product of a combination of abstract domains
- Logical abstraction: in verification tools, based on the Nelson-Oppen satisfiability procedure for combination $f$ theories (SMT solvers)
- The approaches can be combined by
- Generalizing logical abstractions (widenings, ...)
- Understanding Nelson-Oppen satisfiability procedure as an iterated reduced product

Fiscomen oneme

Abstract Domains
(III) Temporal domains


## Application:

Static Analysis of
Dynamic Systems
Static Analysis and
Verification
(I) Quasi-synchronous systems
Example
voter:


- Result of the static analysis:
Counter-example ?"? Specification proved
1

Static Analysis and Verification
(II) Parallel Programs

Static analysis of parallel software

- Critical embedded software are now parallel (e.g., IMA [Integrated Modular Avionics])
- Semantic challenges:
- implicit communications through shared memory
- weak memory consistency
- strict priorities of real-time schedulers
- Abstraction challenges, to scale up while being precise:
- history-sensitive abstractions of interleaved control
flows
- fixpoint strategies to compute interferences
- Ongoing work on the Thésée prototype

Static Analysis
(III) Biological systems
.

- Focus on models of signaling pathways described by
collating biochemical interactions.
- These systems usually suffer from a huge combinatorial complexity in the number of chemical species (i.e.
chemical complexes) which may be formed at run-time.
- We design scalable and precise analyses of the reachable
- Applications:
- Applications:
- pracompute eroperties so as to
fast stochastic simulation;
- automatic simplification of the model;
- compute an idiomatic description of
the systems.

Model reduction

- Quantitative (ODEs and stochastic) semantics are stochastic) semantics are
hard to compute, due to the combinatorial complexity ( $\sim 10^{20}$ variables).
- We use an approximation of the control flow between regions of chemical species, regions of chemical species,
so as to compute exact projections of these semantics.


Static analyzers

## ASTRÉE (commercialized)



- Routinely used by Airbus France (A380, A400M, A350)
- Now in the industrialization and commercialization phase (Abslnt)

THÉSÉE (in development)

- Challenging application to the FWS (A380, A400M, A350)


Dissemination

## Dissemination

- In the academic world (tutorials, summer schools, invited talks, conference organization, ...)
- In the industrial world (conferences, training, ...)

 $=5$ $\pm=$


Conferences recently organized

- SAS 2010 (Perpignan)

SASB in 2011 (Venice)


- NSAD 2010
- SASB 2010
- TAPAS 2010

Program committees
ESOP 2010/1I, HSCC 2009/II, PLDI 2010, POPL 2011, SAS 2010/1I,VMCAI 2009/10/11,...
$\cdots$
${ }_{3} 4$

## Research Contracts

Ongoing contracts

- SURVOL: FNRAE, Logiciel de Commande Embarqués: robustesse et sécurité, avec MIP-Univ. P. Sabatier, Toulouse, 2008-2011
- ASCERT: FNRAE, Analyses Statiques Certifiées, avec INRIA
Rennes, Rocquencourt, Grenoble, $2009-2012$
- SARDANES: Sémantique, analyse et transformation des applications
numériques embarruées synchrones. FNRAE avec Université de Perpignan \& Université de Brest, 2009-2012
- AbstractCell, long-term Junior ANR Chair of Excellence, J. Feret, Dec 2009 - Nov 2013.
- Abstract Interpretation and Code Obfuscation, Royal Society

UK, avec Imperial College, 2010-2012

- ANASTASY: Airbus France, Analyse statique et dynamique, 2010 $-2014$ $\qquad$


## Starting contracts

- MBAT: Model-based Analysis and Testing of Embedded Systems, European project, Artemis programme, with Airbus, Daimler Volvo, Eads, Thales, Rockwell Collins, ... \& Univ. Munich, Univ Aalborg ..., 2011 - 2014

Submitted contract applications

- ANR (appel à projets bioinformatiques)
- DYALOG: Hierarchical dynamical modelling and analysis of large regulatory networks controlling cell fate decisions, avec IBEns/
Ens, Institut Curie, Université Montpellier 2, projet INRIA Contraintes, Pasteur, IBisc (Erry), Technologies avancées pour le Génome et la clinique (Marseille), 13 December
2010
- 

Forthcoming contract applications

- ANR
- THÉSÉE: Static analysis of embedded parallel programs, avec Airbus France, March 201 I
- Europe (FP7):
- EvolBrid: A correct-by-evolution design and verification framework for hybrid-modeled complex embedded systems, FP7 (7 partners), January 201
- Ratis: Secure programming of embedded systems, FP7 (8 partners), December 201

> Main industrial collaborations

AbsInt Angewandte Informatik

- Industrialization of ASTRÉE www.absint.com/astree/

- 3 full-time engineers, everyday collaboration
$\qquad$
- Verification of space software


Sagem

- Analysis of inertial unit software

$x=2$

Airbus France

- Verification of parallel programs, mainly the FWS


Microsoft Research Redmond
- Verification of code contracts in Clousot (now distributed with Visual Studio Ultimate)
- Design of scalable abstract domains for
- Array content analysis
- Contract precondition inference from assertions



## Ongoing work ..

- Analysis of complex discrete/continuous systems:
- Parallel programs: complex data structures, interference, real-time, scheduling with dynamic priorities (e.g. priority ceiling protocols), ...
- Biological systems
- Closed loop analysis of control/command system
- Verification of complex properties:
- Security properties
- Probabilistic properties
- Eventuality properties

Conclusion

- Small team, smart project, great scientific ambitions :-)

